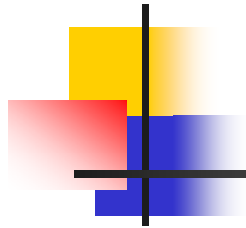




# **Chapter 10** Spiral-Multi-Path QoS Routing Protocol in Wireless Mobile Ad-Hoc Networks

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Information Engineering  
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## Publication

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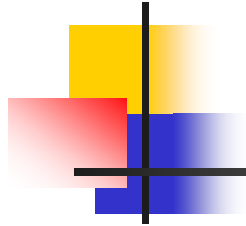
- **Yuh-Shyan Chen** and Yu-Ting Yu,  
“Spiral-Multi-Path QoS Routing Protocol  
in Wireless Mobile Ad-Hoc Networks”,  
***IEICE Transactions on  
Communications***, Vol.E87-B, No.1,  
pp.104-116, Jan. 2004.



# Outline

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- Introduction
- Basic Idea
- Our proposed SMPQ Protocol
- Experimental Results
- Conclusion



## Introduction

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- This paper proposes a QoS routing protocol in MANET
  - MAC layer adopts CDMA-over-TDMA scheme (multiple access scheme)
- Two important schemes are integrated
  - Spiral-path
  - Multi-path



## Existing QoS Routing Protocols

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- A quite ideal model
  - Ticket-based QoS routing protocol [JSAC99]
  - A Multi-Path QoS Routing Protocol in a Wireless Mobile Ad Hoc Network [ICN'01]
- CDMA-over-TDMA channel model
  - Lin approaches [JSAC99][INFOCOM01]
    - Uni-Path



## Motivation

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- Design a **new routing protocol** under CDMA-over-TDMA model with
  - High success rate of a QoS route
  - Mobility-tolerant capability



## Channel model Assumption

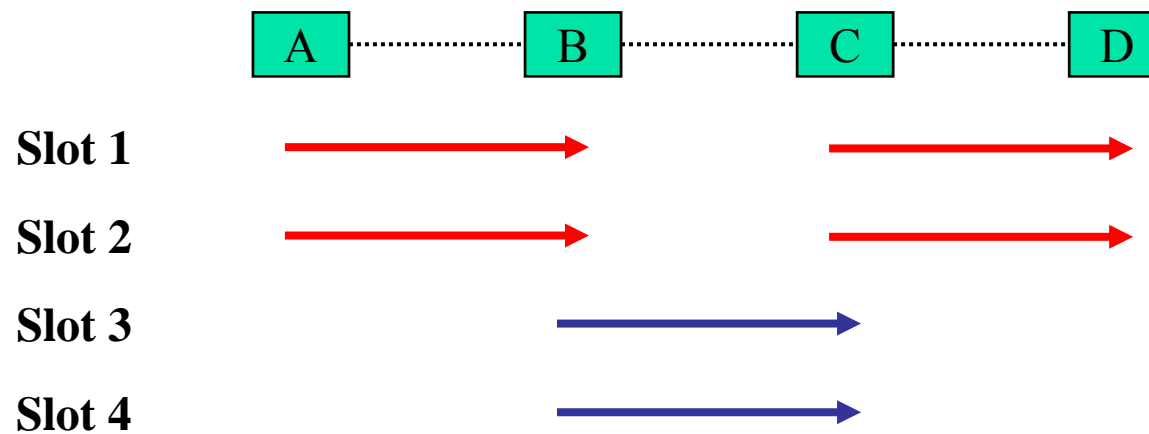
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- CDMA-over-TDMA channel model
  - Use an orthogonal code to overcome the hidden-terminal problem
  - The use of a time slot on a link is only dependent of it's one hop neighboring links



# CDMA-over-TDMA Model

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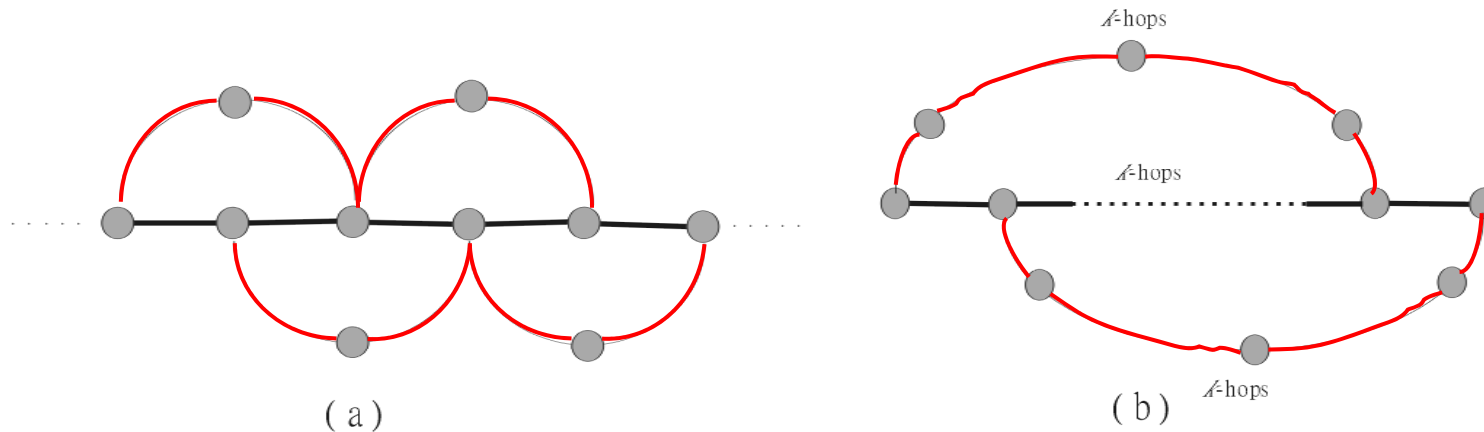


## Basic Idea

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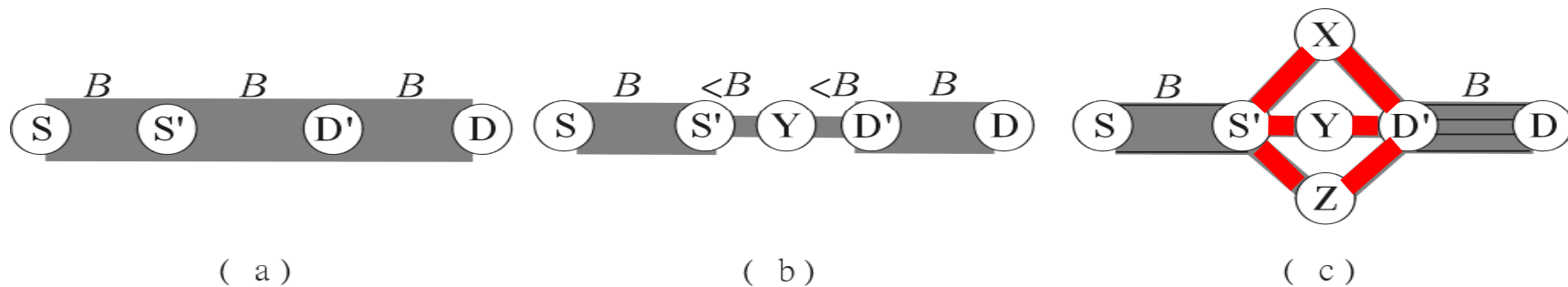
- Propose **spiral-multi-path routing** by combining
  - **spiral-path** routing
  - **multi-path** routing
- Aims
  - High **success rate** of a QoS route
  - With well **mobility-tolerant** capability

# MESH: Multi-Eye **Spiral-Hopping** Protocol in a Wireless Ad Hoc Network



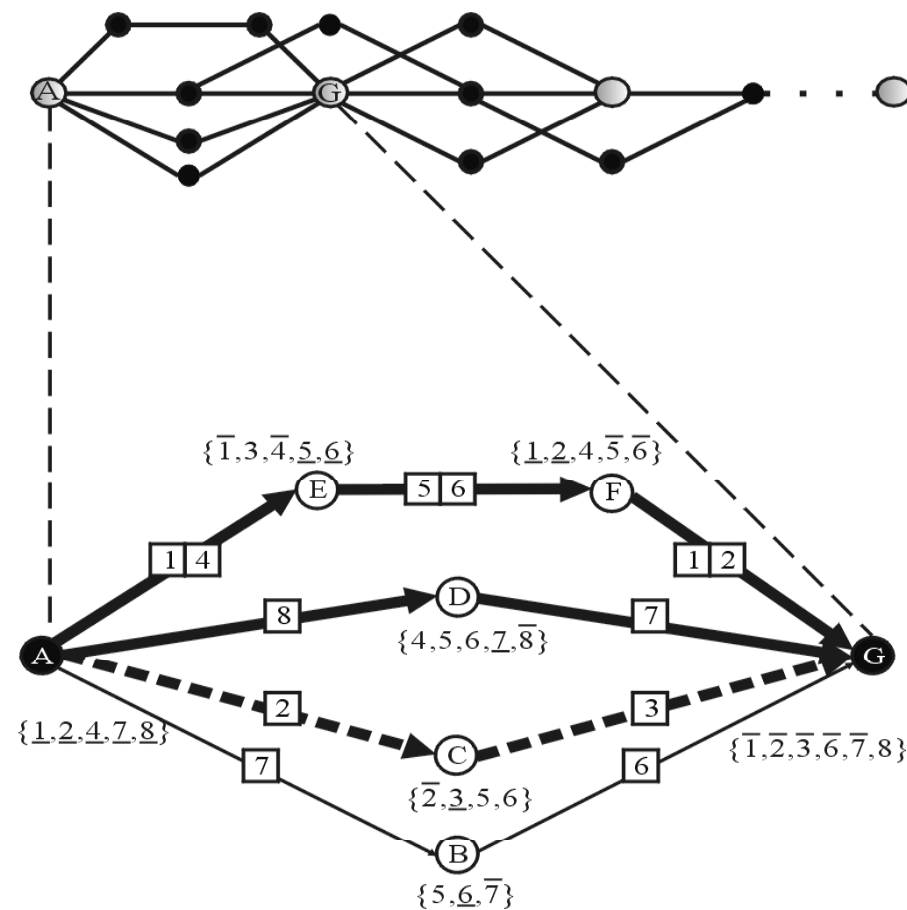
## A Spiral-path approach

# A Multi-Path QoS Routing Protocol in a Wireless Mobile Ad Hoc Network

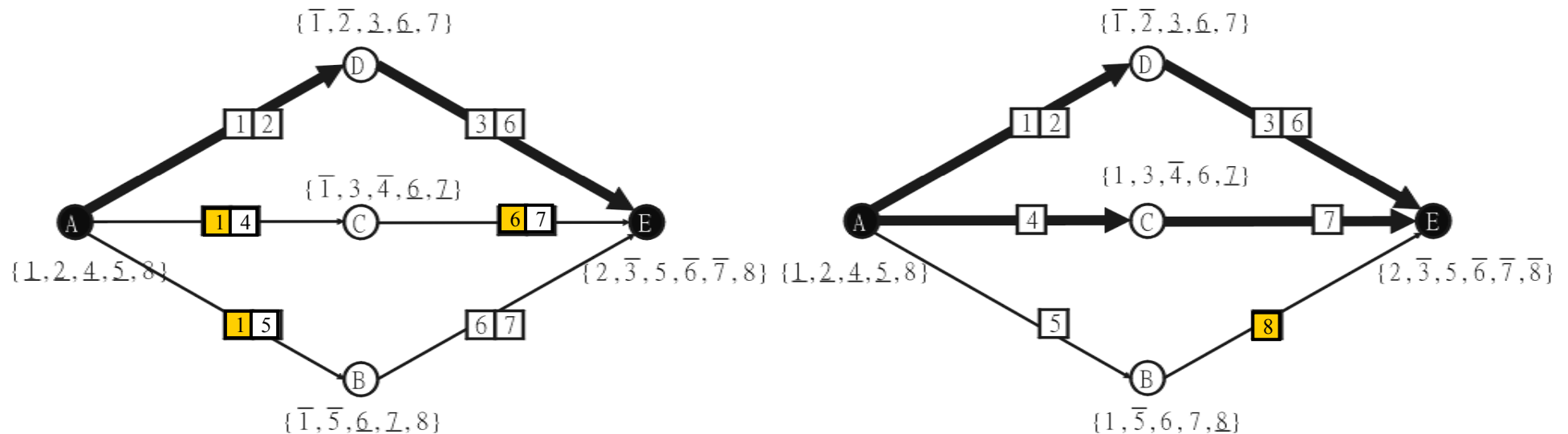


**A multi-path approach**

# Informal definition of *spiral-multi-path*



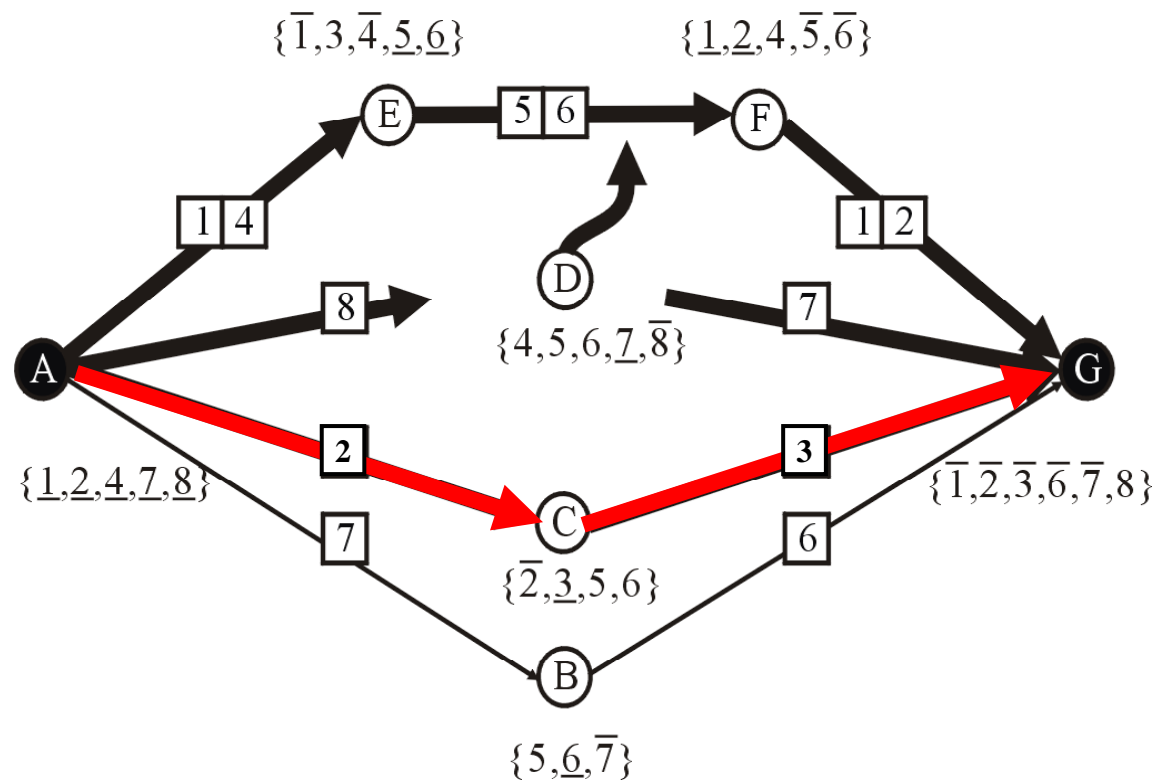
# Difference of slot reservation between uni-path and multi-path



(a) uni-path

(b) multi-path

# Mobility-tolerant capability





## Our SMPQ Protocol

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- Phase 1: **Keeping Link Bandwidth**
  - Keep the information of link bandwidth in MANET
- Phase 2: **QoS Route-Discovery Phase**
  - Find the *spiral-multi-path* and reserve time slots
- Phase 3: **QoS Route-Reply Phase**
  - Confirm a final *spiral-multi-path* and send a reply packet
- Phase 4: **QoS Route-Maintenance Phase**



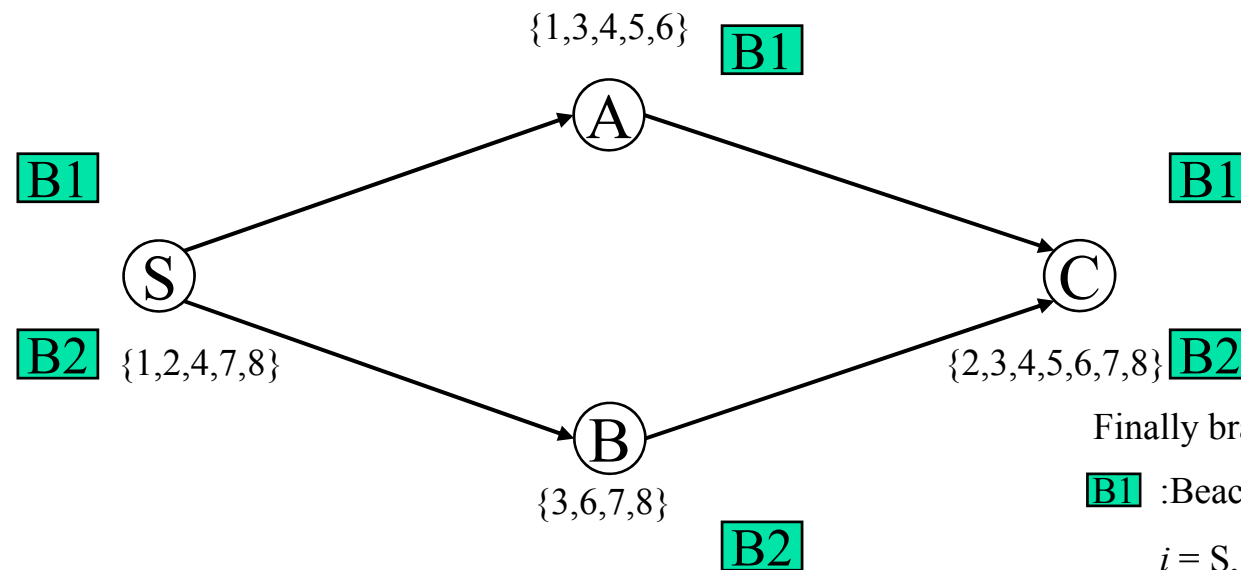
# Phase 1: Keeping Link Bandwidth

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- The purpose of this phase
  - Keep the information of link bandwidth in MANET
- **Branch Node**
  - If there exist at least two disjoint paths between two nodes, then these two nodes are said as branch nodes
- **Supernode**
  - The gateway nodes between a pair of branch nodes



## Identifying operation of *branch node*



Finally branch node C receive

**B1** :Beacon<sub>1</sub>(path\_record=[S,A,C], free\_slot<sub>*i*</sub>)

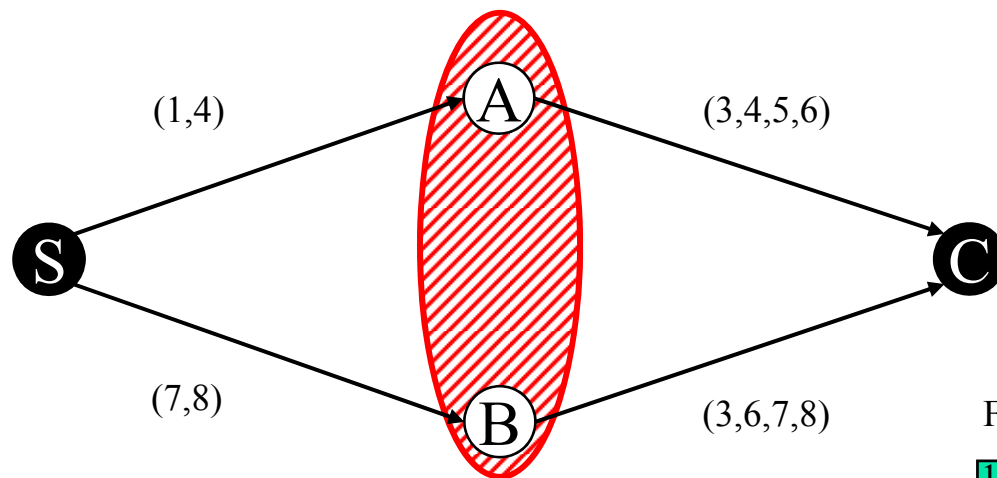
*i* = S, A, C

**B2** :Beacon<sub>2</sub>(path\_record=[S,B,C], free\_slot<sub>*i*</sub>)

*i* = S, B, C

**B1** :Beacon( hopnumber, path\_record, free\_slots<sub>path\_record[*i*]</sub> )

## Identifying operation of *Supernode*



Finally branch node C receive

**1.1** :Beacon<sub>1</sub>(path\_record=[S,A,C], free\_slot<sub>i</sub>)

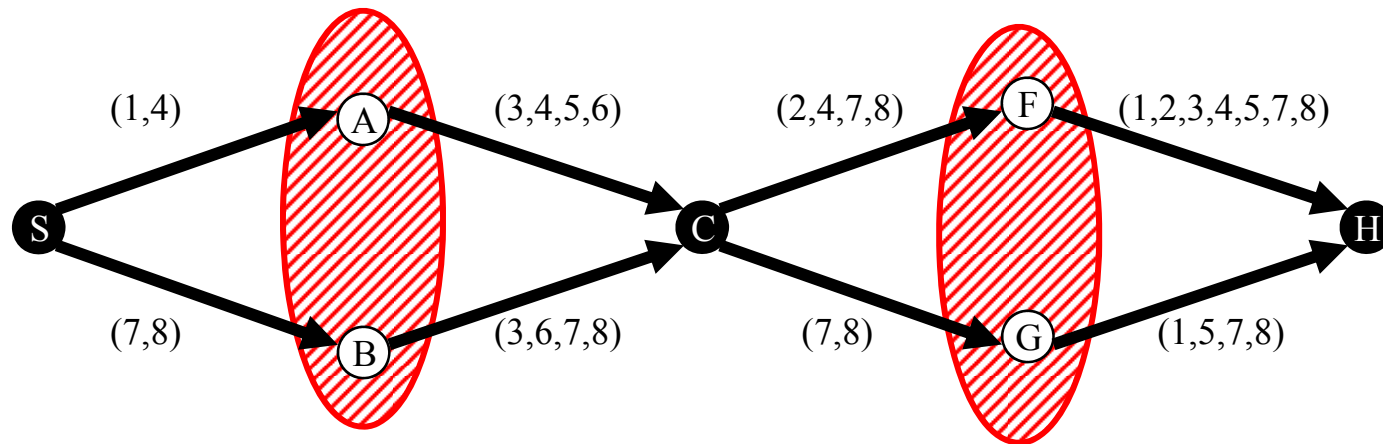
$i = S, A, C$

**1.2** :Beacon<sub>2</sub>(path\_record=[S,B,C], free\_slot<sub>i</sub>)

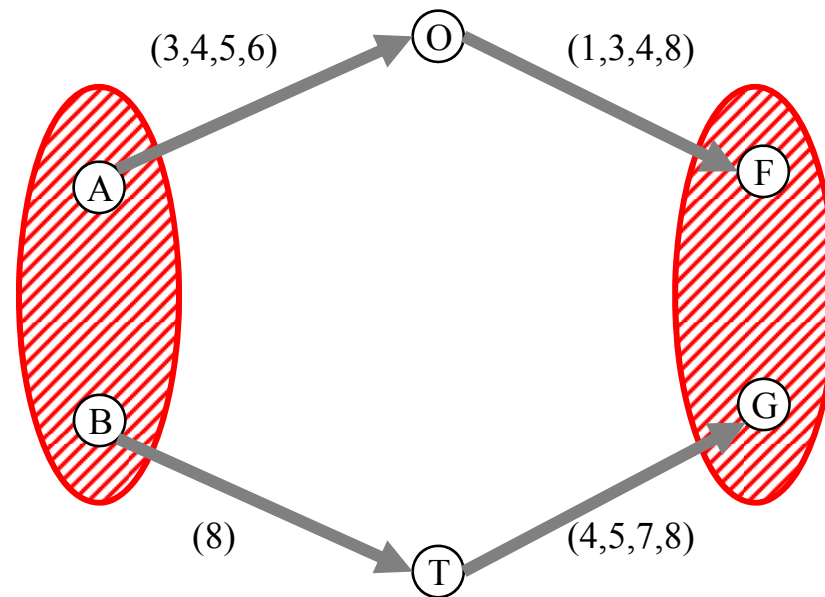
$i = S, B, C$

**1** :Beacon( hopnumber, path\_record, free\_slots<sub>path\_record[i]</sub> )

## Identifying operation of *branch supernode*



## Identifying operation of *branch supernode* and its link bandwidth





## Phase 2: QoS Route-Discovery Phase

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- The purpose of this phase
  - Reserve possible time slots during constructing the *spiral-multi-path*

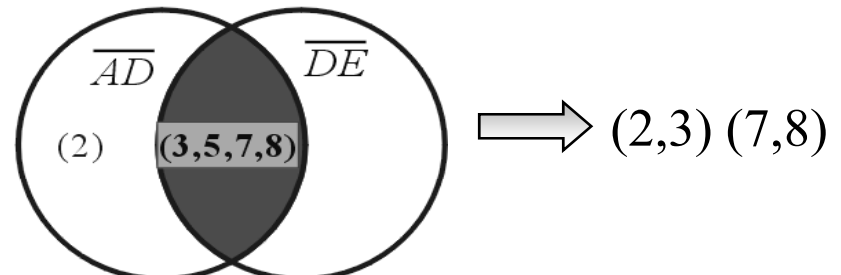
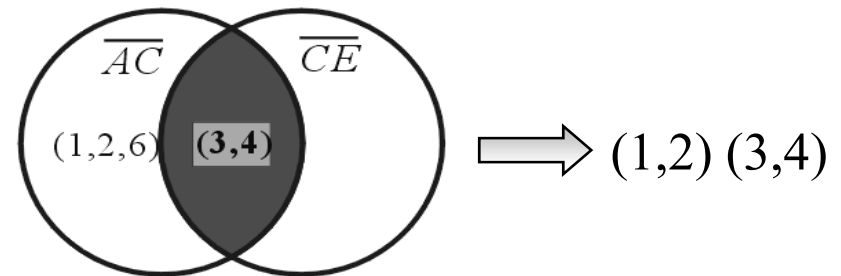
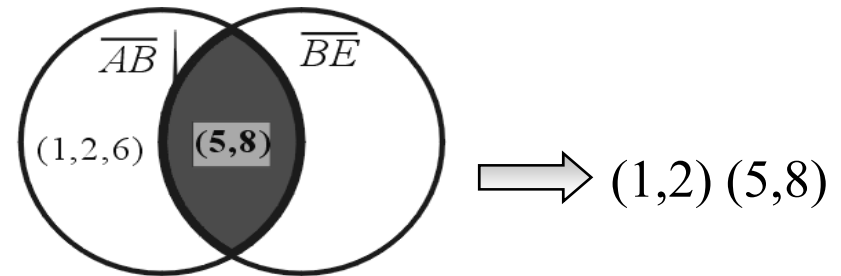
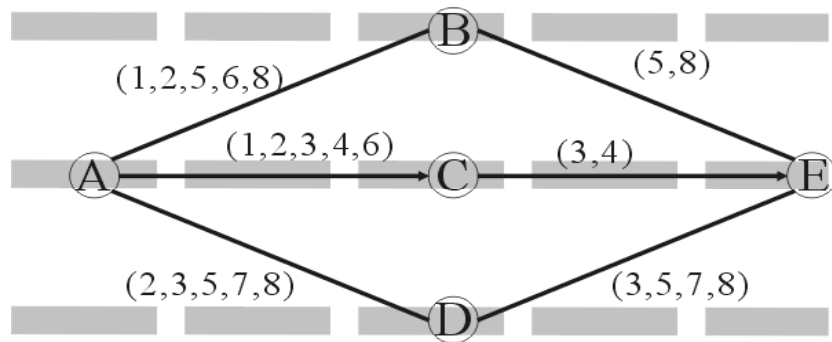


## A basic sub-path bandwidth reservation operation

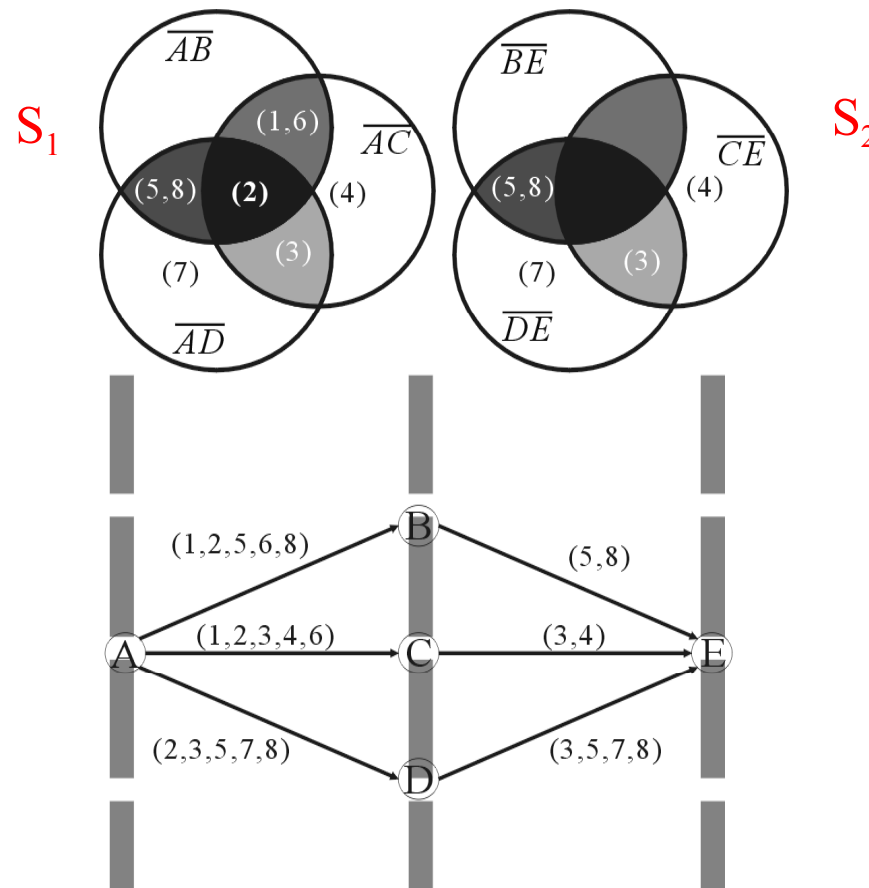
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- Check the available bandwidth of each uni-path
- A path with higher maximum sub-path bandwidth has priority
- The slots not in the intersection of links will be reserved first

# An example of bandwidth reservation operation

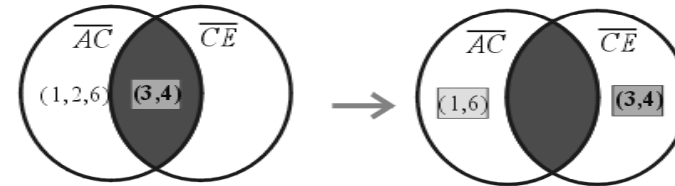
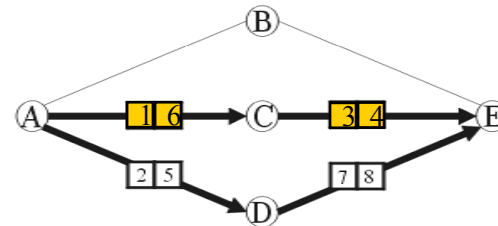
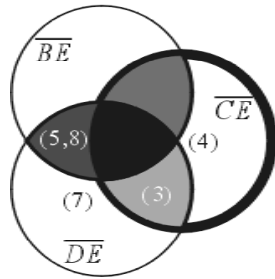
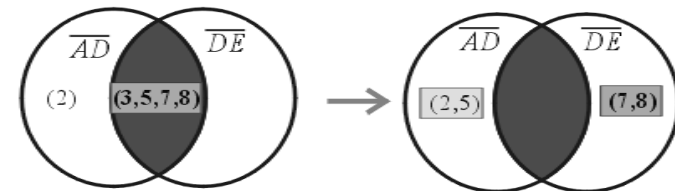
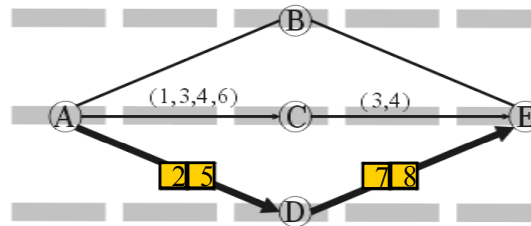
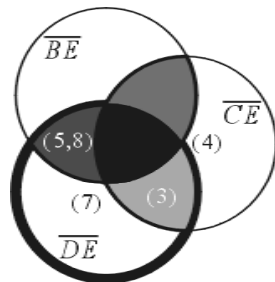


# An example of bandwidth reservation operation (con.)

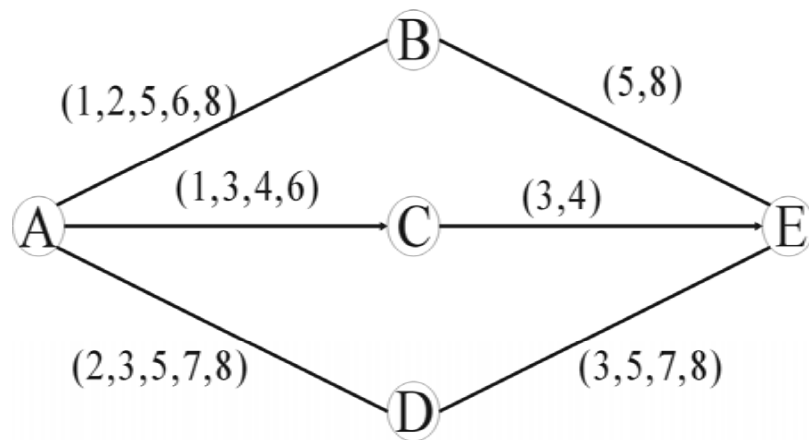




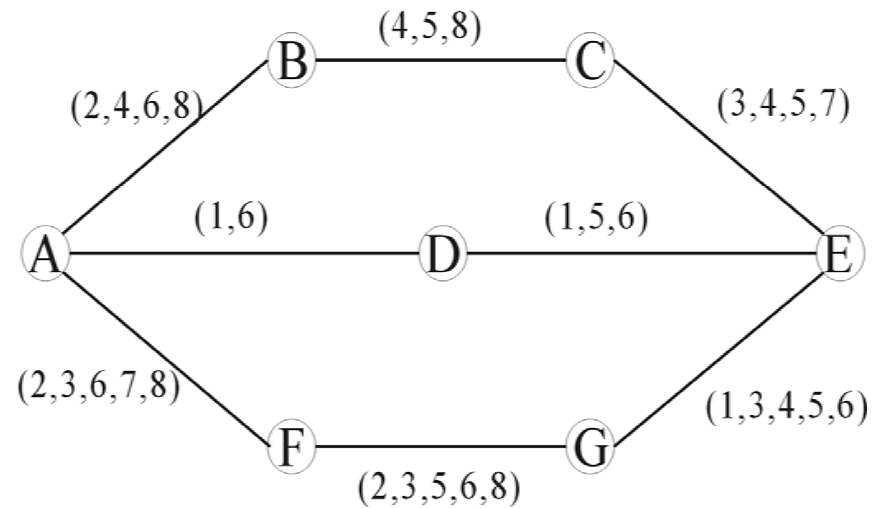
# A basic sub-path bandwidth reservation operation (con.)



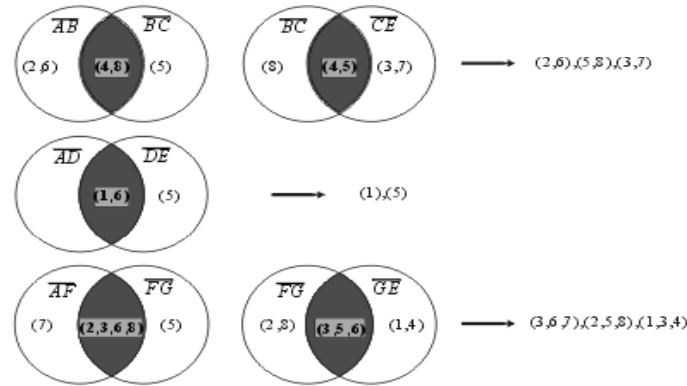
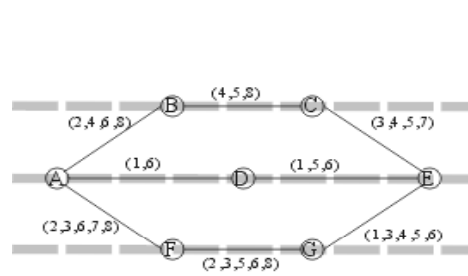
## Other example of basic sup-path bandwidth reservation operation



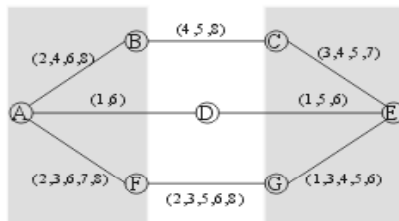
(a)



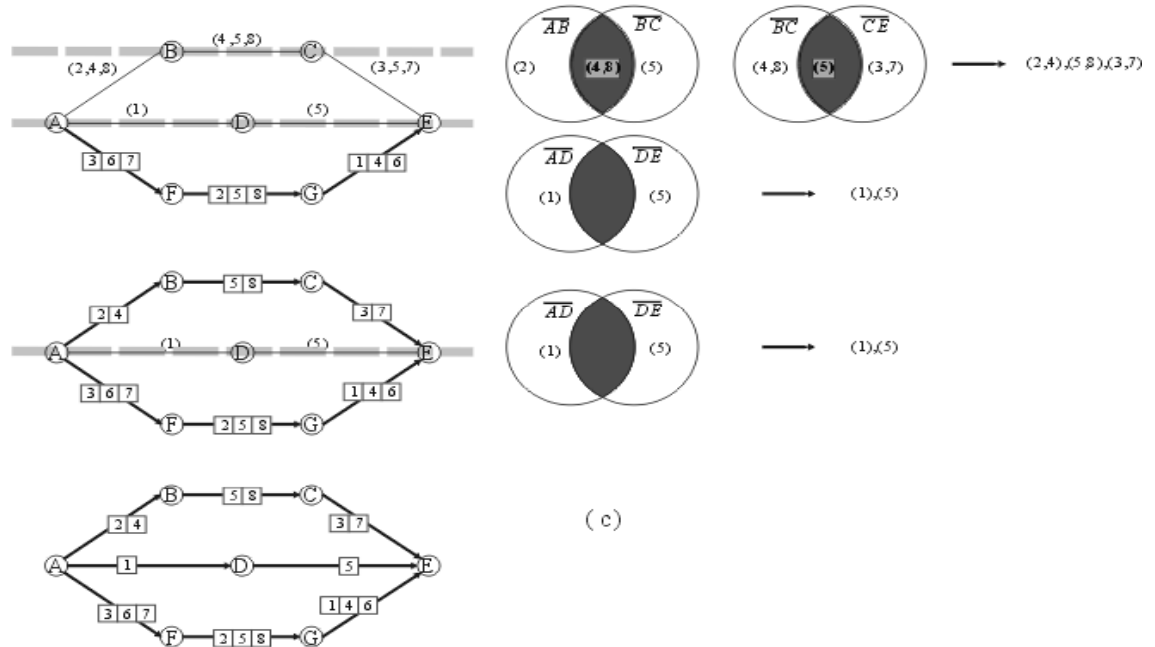
(b)



(a)



(b)



(c)

**Other example of basic sup-path bandwidth reservation operation**



## Path Bandwidth

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- Path bandwidth of a feasible path from S to D

$$\Rightarrow [\overline{\alpha}_1, \underline{\beta}_1, \overline{\alpha}_2, \underline{\beta}_2, \dots, \underline{\beta}_{k-1}, \overline{\alpha}_k], i \geq 1$$

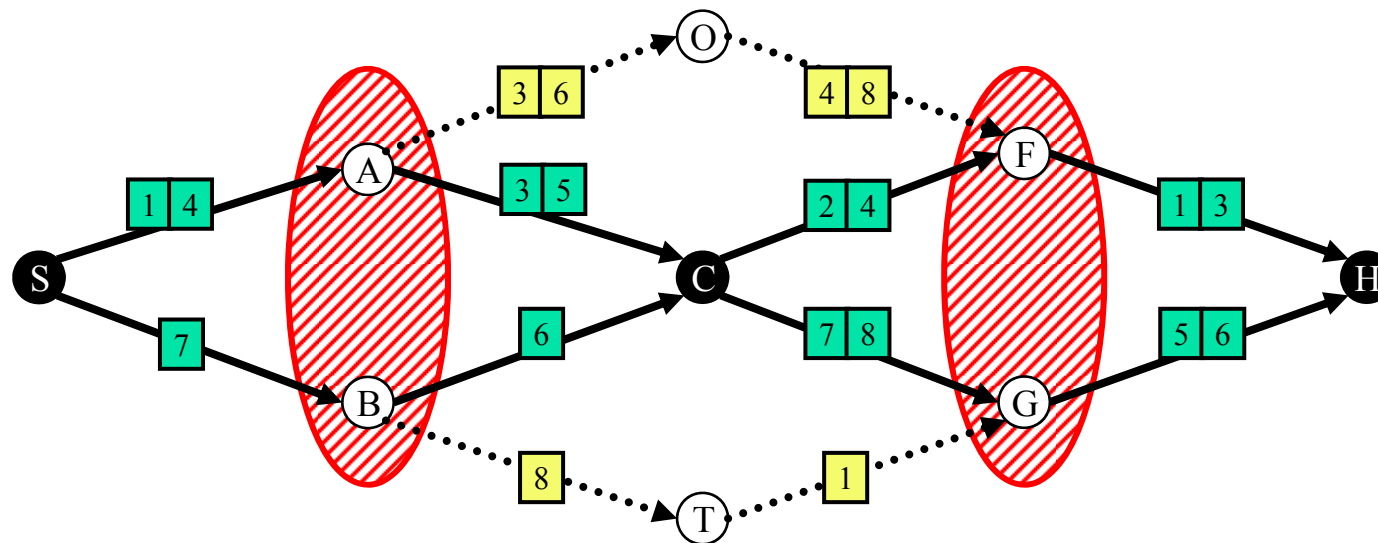
Where

$\overline{\alpha}_i$  is the reservable bandwidth between a pair of branch node

$\underline{\beta}_j$  is the reservable bandwidth between a pair of branch supernodes

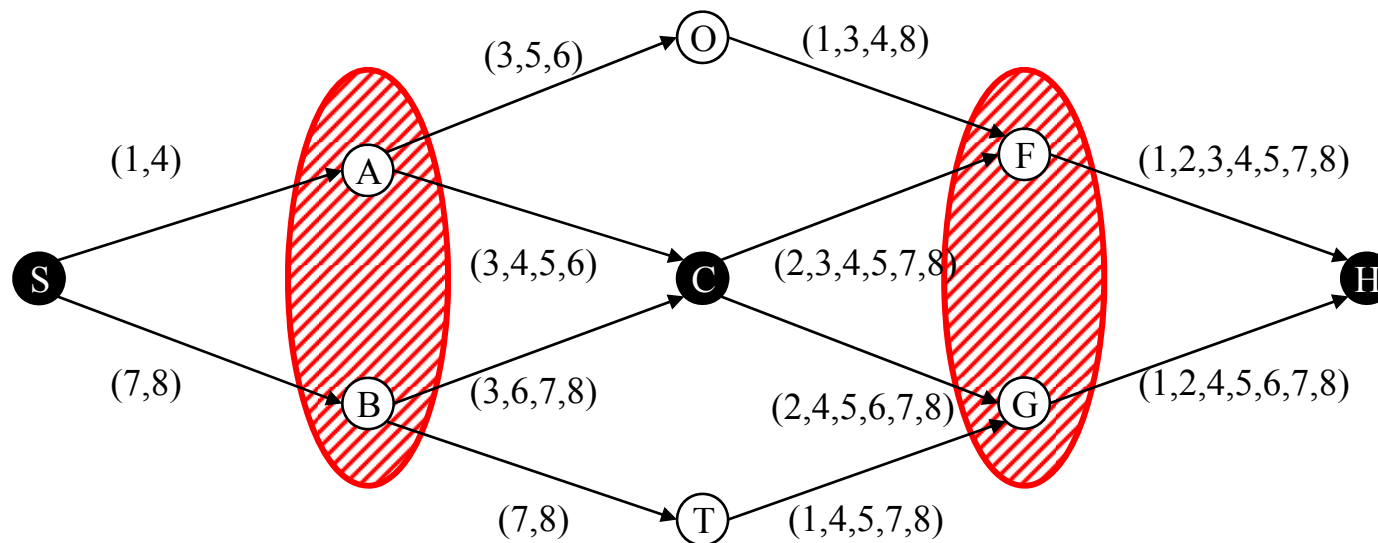
# Path Bandwidth—an example of $[\overline{3}, \underline{3}, \overline{4}]$

Requirement=2



# Route-discovery operation

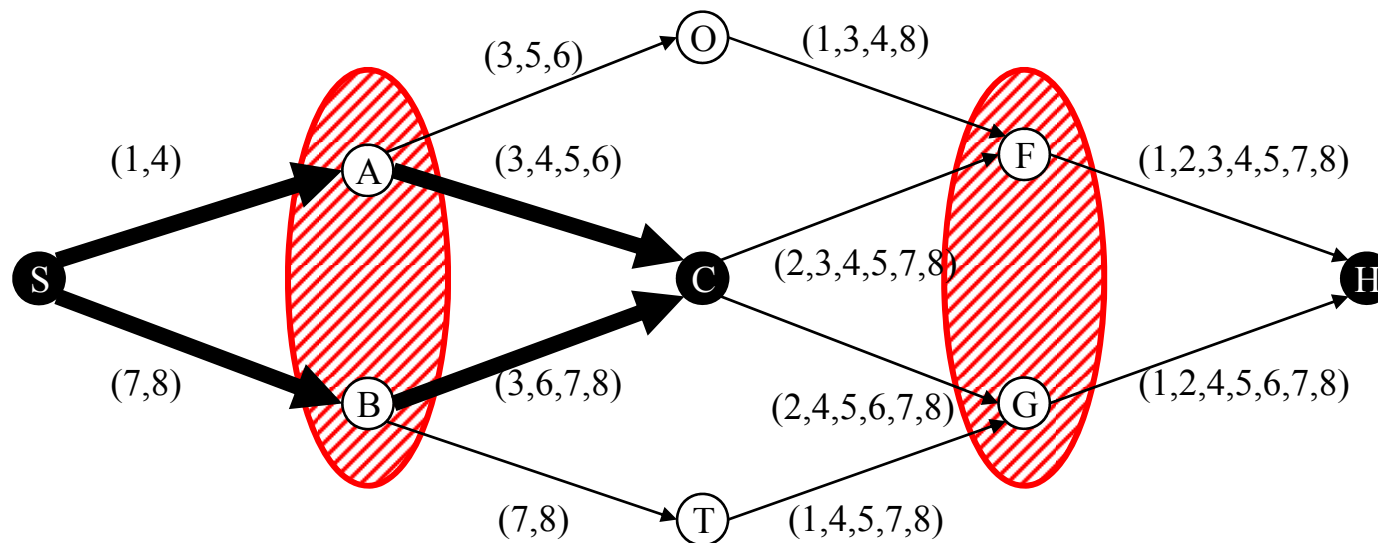
Requirement=2



An example of  $[\bar{3}, \underline{3}, \bar{4}]$

# Route-discovery operation

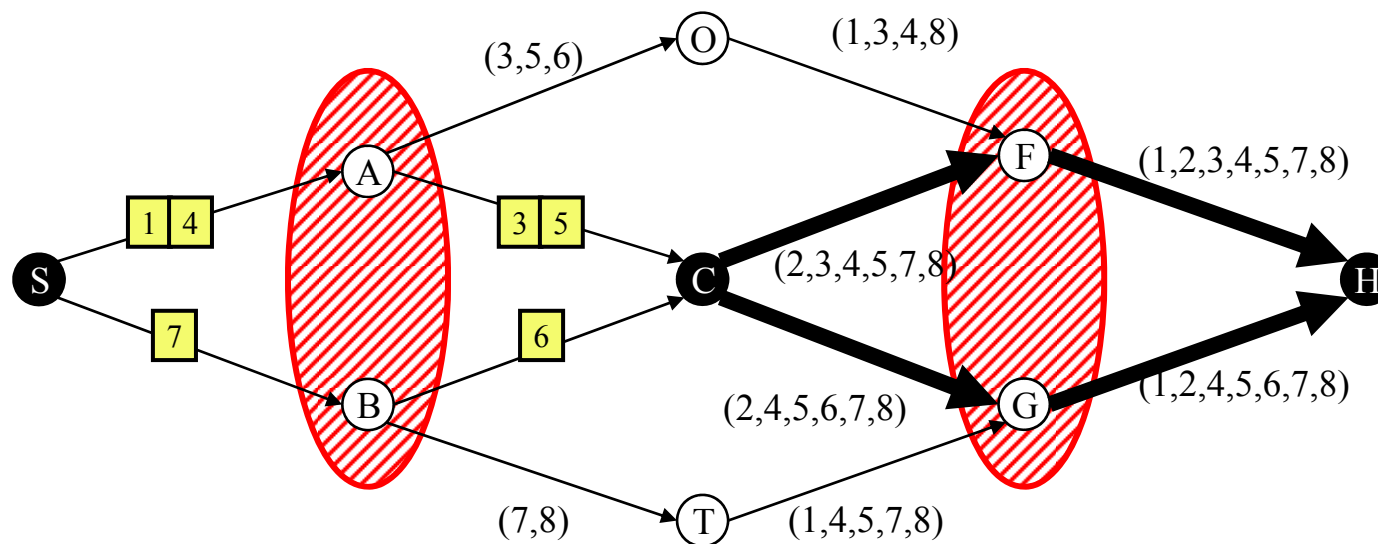
Requirement=2



An example of  $[\bar{3}, \underline{3}, \bar{4}]$

# Route-discovery operation

Requirement=2

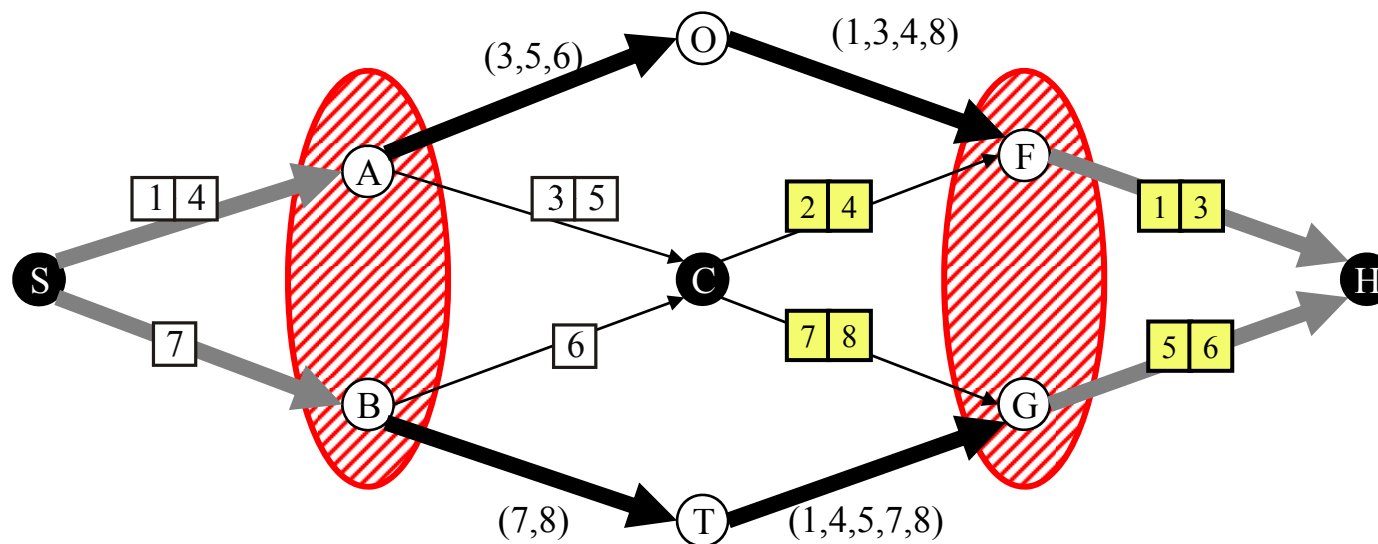


An example of  $[\bar{3}, \bar{3}, \bar{4}]$



# Route-discovery operation

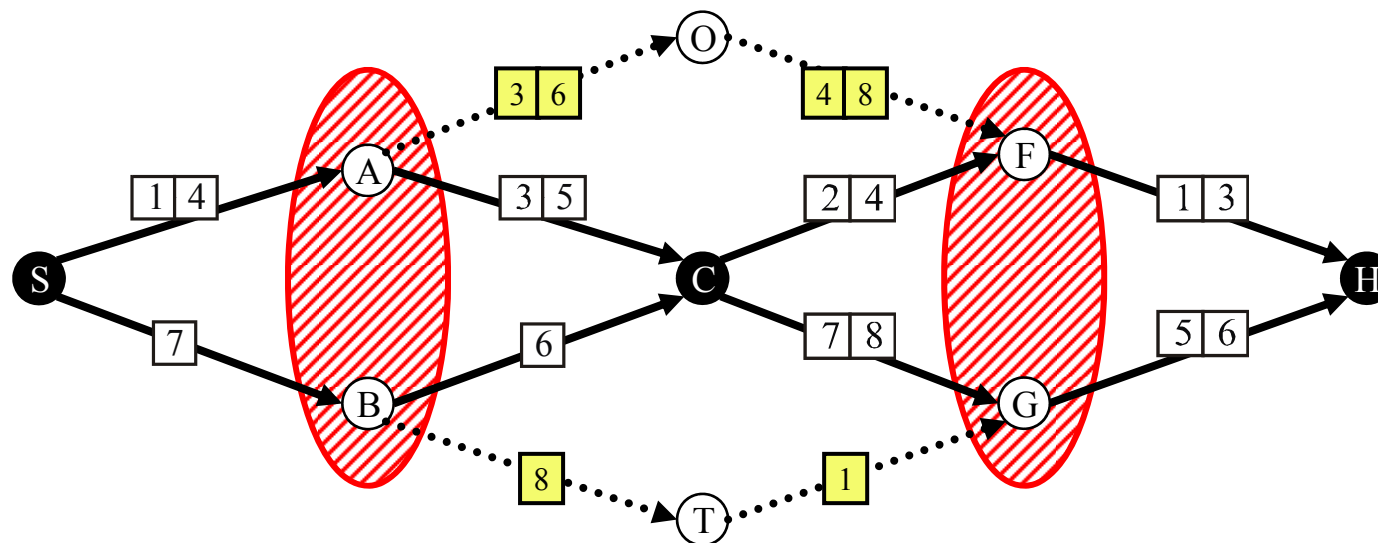
Requirement=2



An example of  $[\bar{3}, \bar{3}, \bar{4}]$

# Route-discovery operation

Requirement=2



An example of  $[\bar{3}, \underline{3}, \bar{4}]$

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## Phase 3: QoS Route-Reply Phase

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- The purpose of this phase
  - Confirm a final *spiral-multi-path* and send a reply packet



## Route-reply operation

■ **Path bandwidth**  $\Rightarrow [\overline{\alpha_1}, \underline{\beta_1}, \overline{\alpha_2}, \underline{\beta_2}, \dots, \underline{\beta_{k-1}}, \overline{\alpha_k}], i \geq 1$

■ **Mobility-tolerant capability for gateway node:**

A spiral-multi-path satisfies condition,

$$[\overline{\alpha_j}]_i \geq \gamma, \text{ for all } 1 \leq j \leq k$$

■ **Mobility-tolerant capability for branch node:**

A spiral-multi-path satisfies condition,

$$[\underline{\beta_j}]_i \geq \gamma, \text{ for all } 1 \leq j \leq k-1$$

where  $\gamma$  is the time slot requirement



## Route-reply operation (con.)

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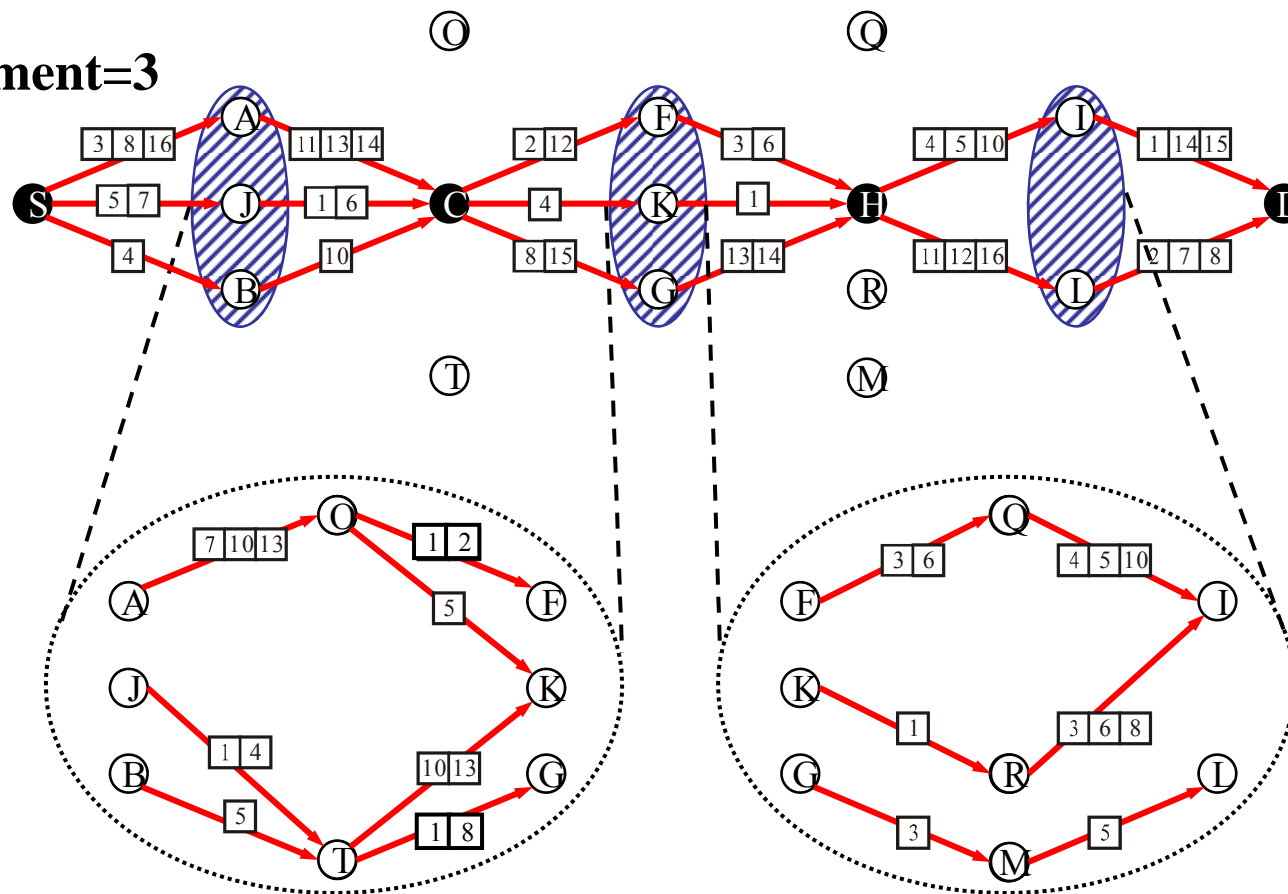
- **Selecting path decision:**

$$EB_{av} = \overline{EB}_{av} + \underline{EB}_{av} = \frac{\sum_{i=1}^k |\overline{\alpha}_i - \gamma|}{k} + \frac{\sum_{i=1}^{k-1} |\underline{\beta}_i - \gamma|}{k-1}$$

**The higher the value is, the high QoS route stability will be**

# An example of $[\bar{6}, \underline{6}, \bar{5}, \underline{5}, \bar{6}]$

Requirement=3





## An example of $[\bar{6}, \underline{6}, \bar{5}, \underline{5}, \bar{6}]$

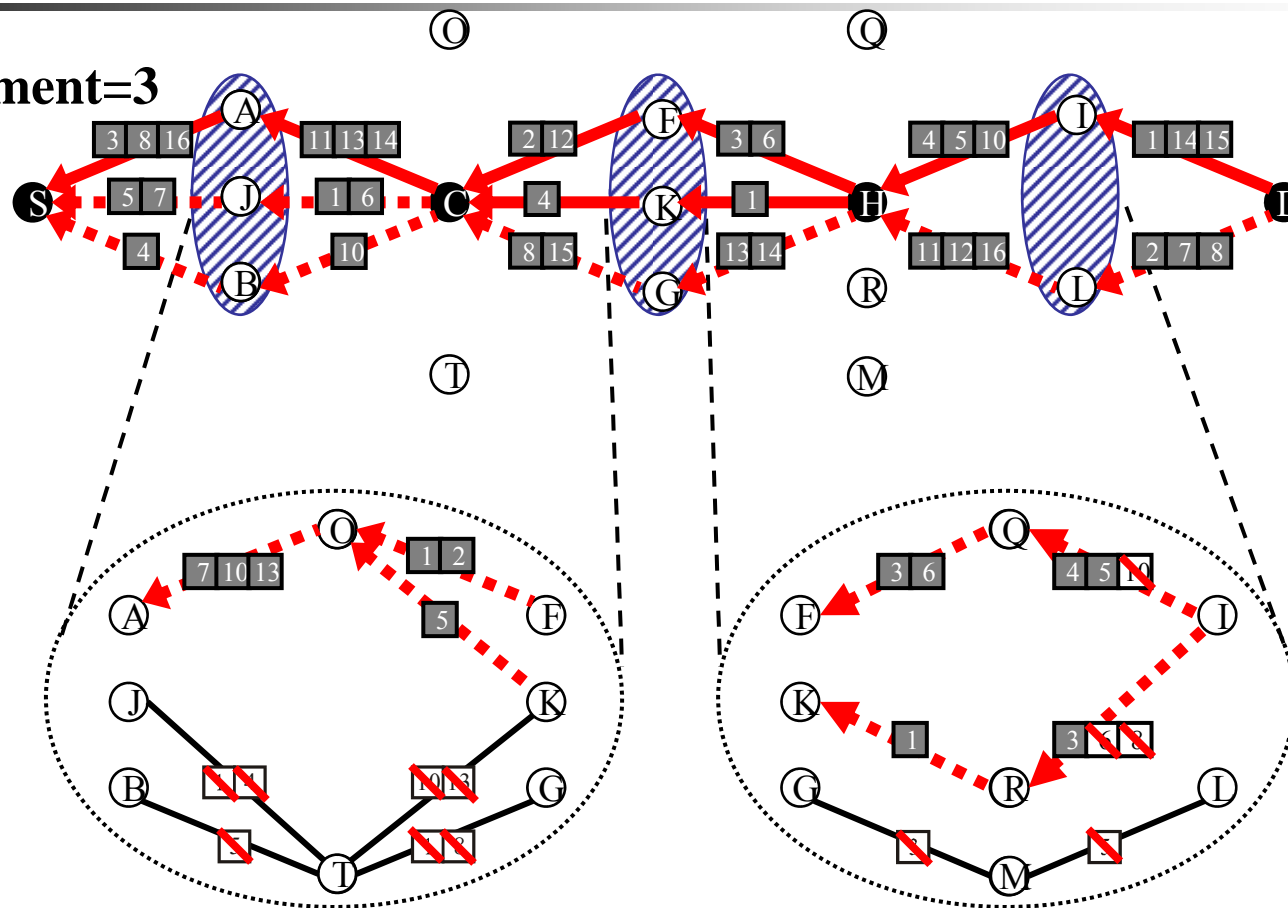
Ex.  $[\bar{6}, \underline{6}, \bar{5}, \underline{5}, \bar{6}]$

$$\Rightarrow EB_{av} = \overline{EB}_{av} + \underline{EB}_{av} = \frac{|6-3| + |5-3| + |6-3|}{3} + \frac{|6-3| + |5-3|}{2} = 5.16$$



# QoS Route-Reply operation

Requirement=3





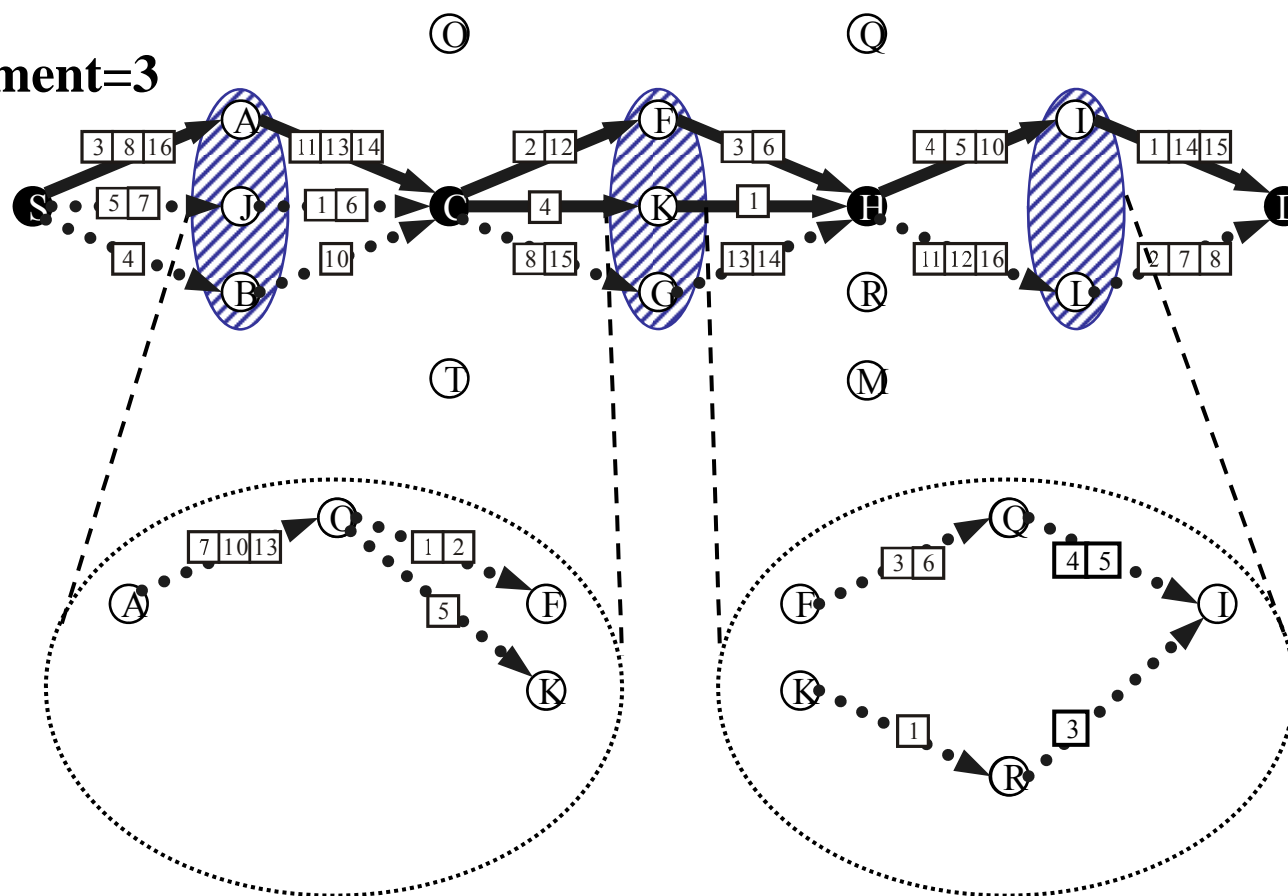
## Phase 4: QoS Maintenance Phase

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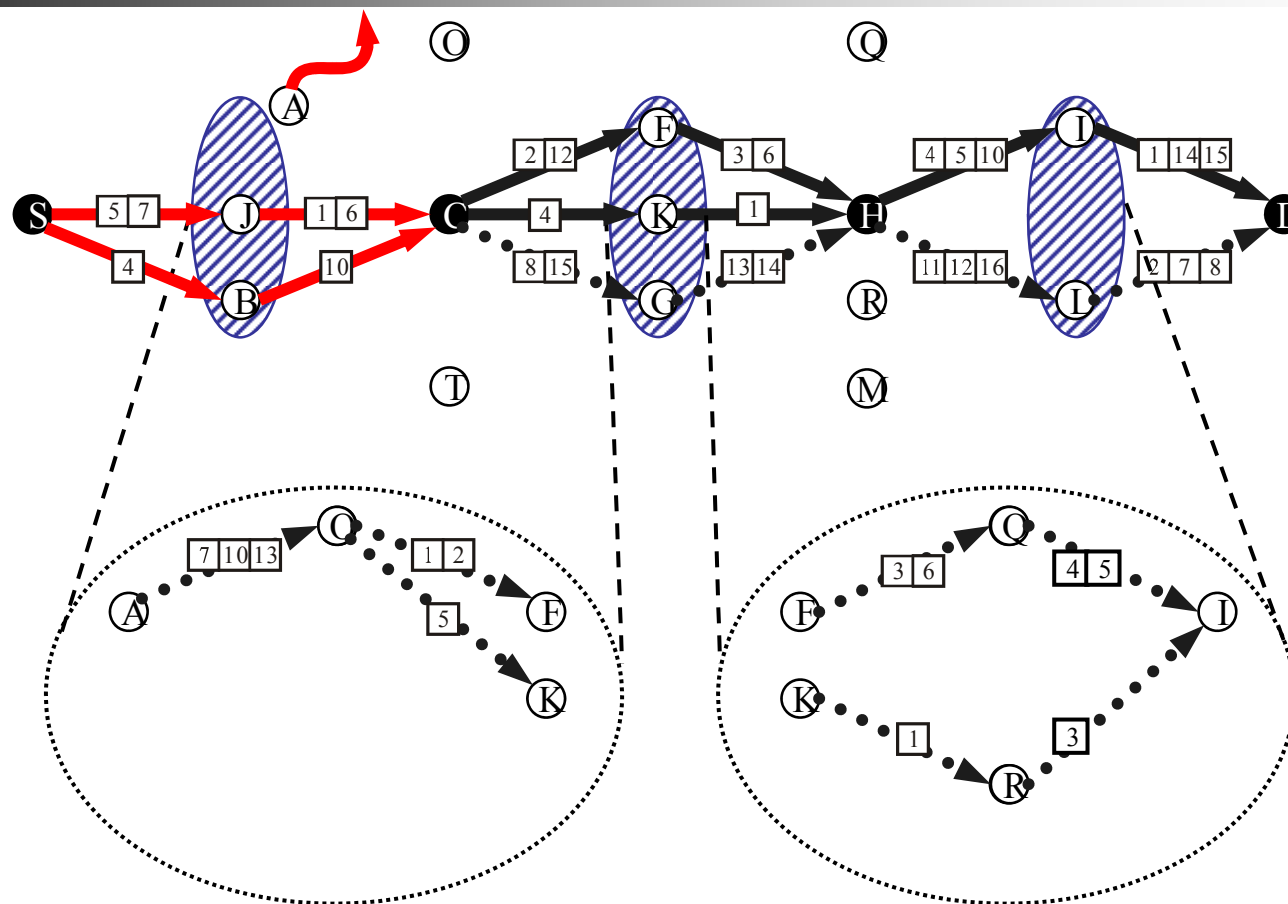
- Our main contribution is to provide the on-line route-recovery capability

# Tolerating failed gateway node

Requirement=3

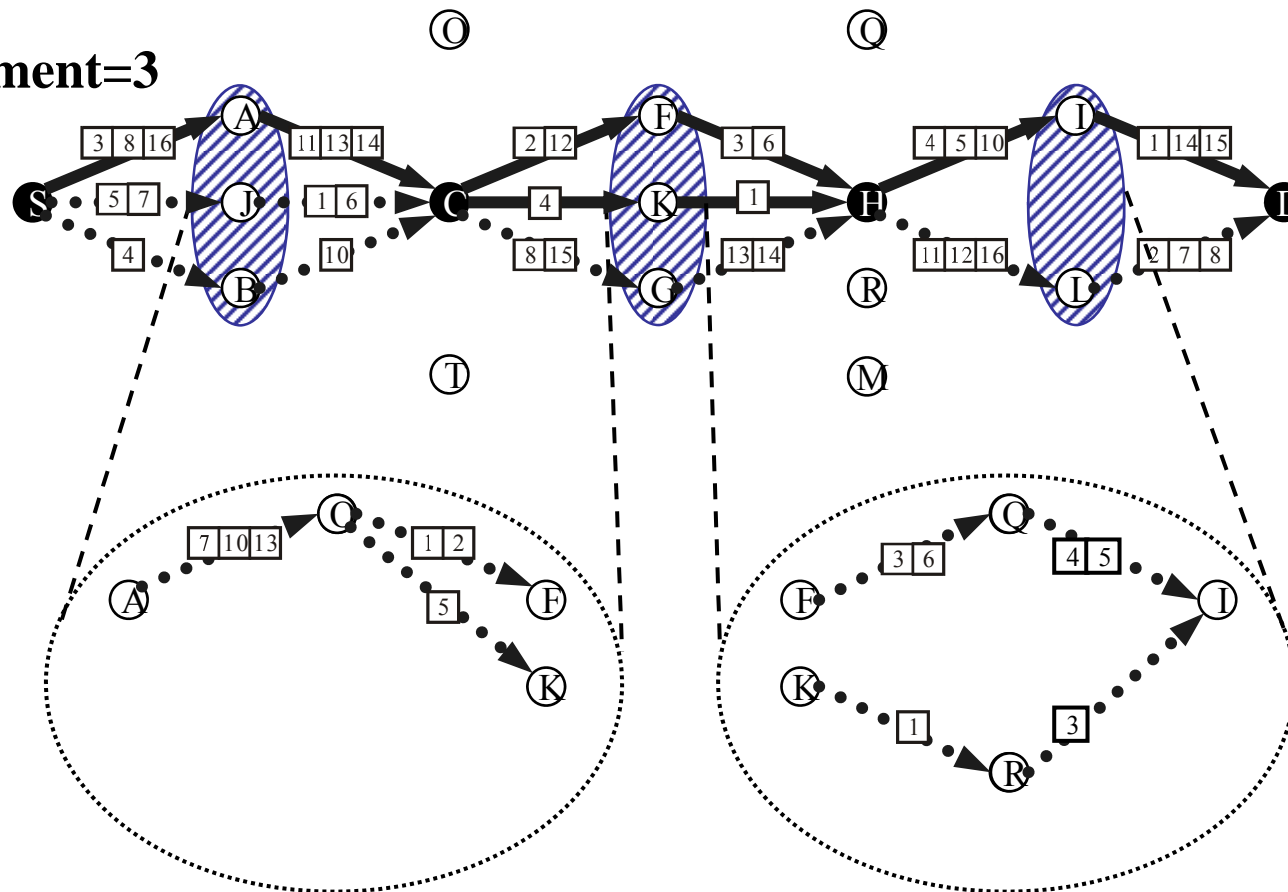


# Tolerating failed gateway node

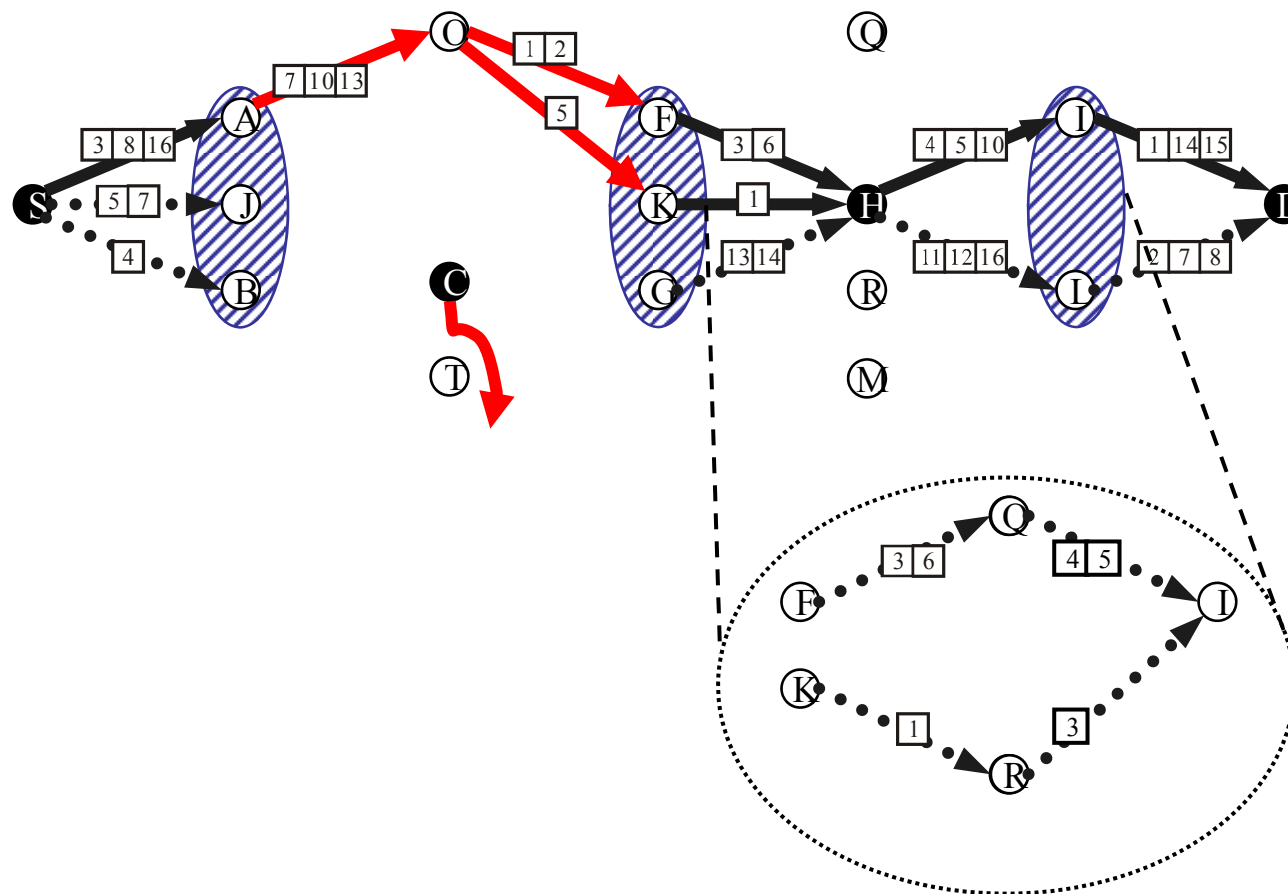


# Tolerating failed branch node

Requirement=3



# Tolerating failed branch node





## Experimental Results - The parameters of our simulation platform

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- The simulation platform is simulated in  $1000 \times 1000 \text{ m}^2$
- The number of mobile hosts is from 20 to 40
- The total number of time slots of each link is 8, 12 and 16
- Three different bandwidth requirement are 1, 2, 4, each called Lin-1, Lin-2, Lin-4 and SMPQ-1, SMPQ-2, and SMPQ-4
- The network bandwidth are low(25%), medium(50%), and high(75%)



## Experimental Results - Performance Metrics

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- **Success\_Rate(SR):**
  - the number of successful QoS route requests divided by the total number of QoS route requests from source to destination
- **Slot\_Utilization(SU):**
  - the average slot utilization of every link in all QoS routes
- **Overhead (OH):**
  - the number of packets used for constructing and maintaining the QoS route from source to destination



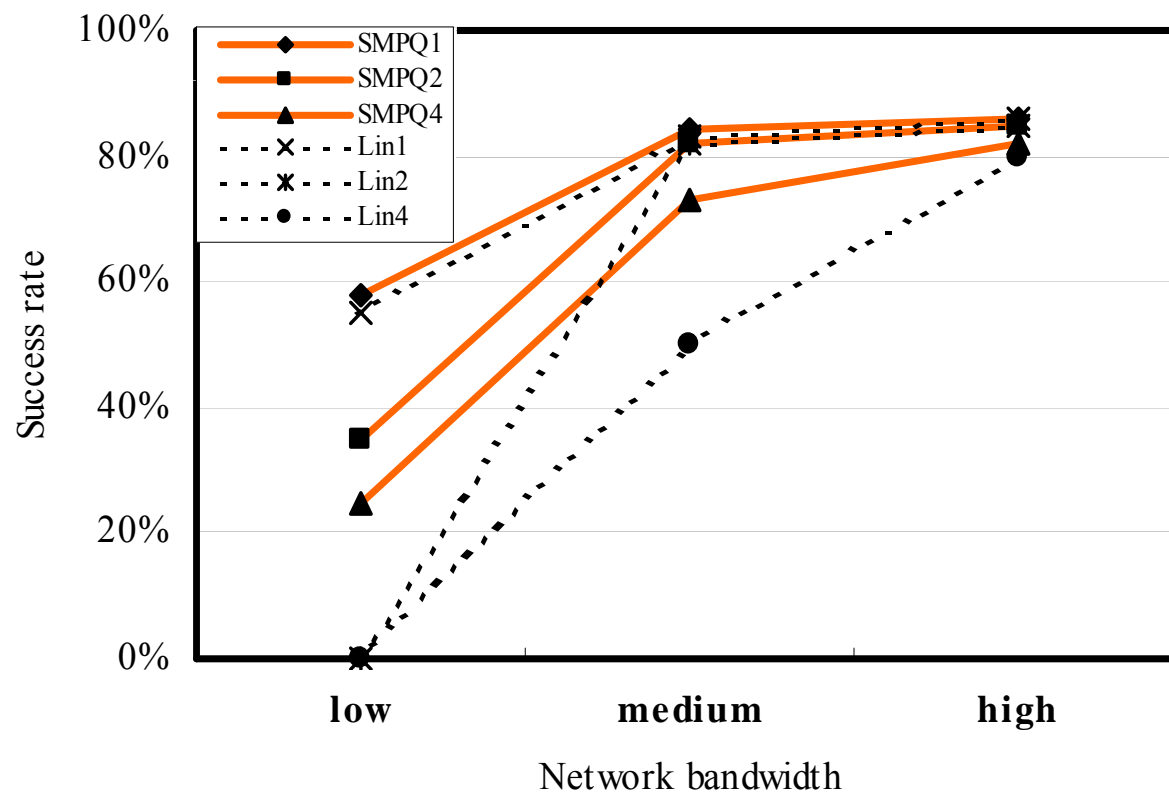


## Performance of *Success\_Rate(SR)*

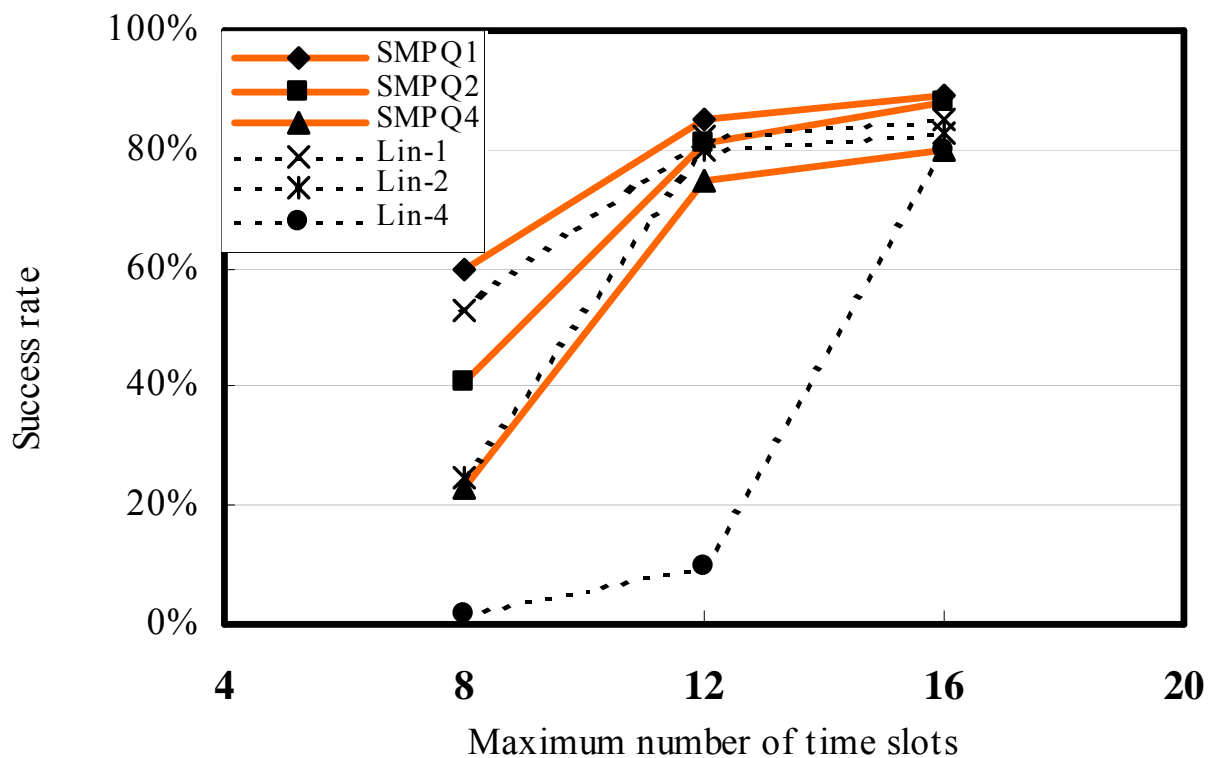
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- Effects of *Network Bandwidth*
  - Our scheme has higher SR than the Lin's when the network bandwidth is low
  - More Bandwidth, high SR
- Effects of *Maximum Number of Time Slots*
  - Our scheme has higher SR than the Lin 's
  - More Time Slots, high SR

## Effects of *Network Bandwidth*



## Effects of *Maximum Number of Time Slots*



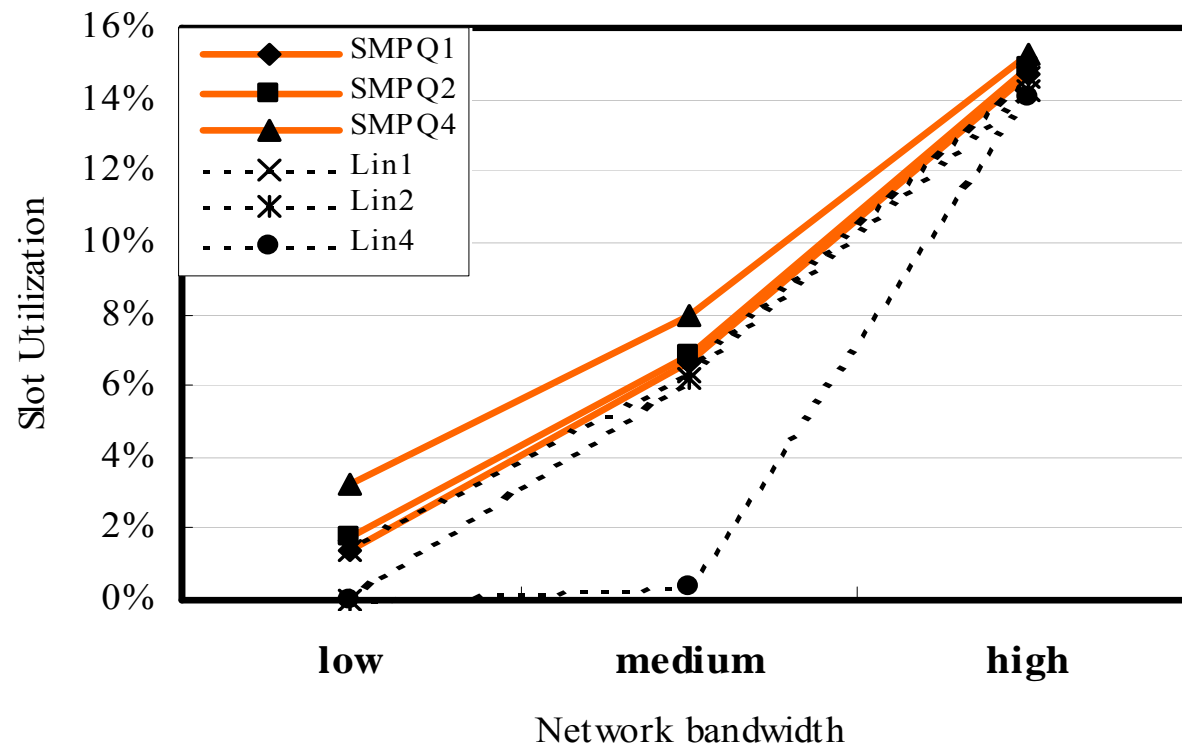


## Performance of *Slot\_Utilization(SU)*

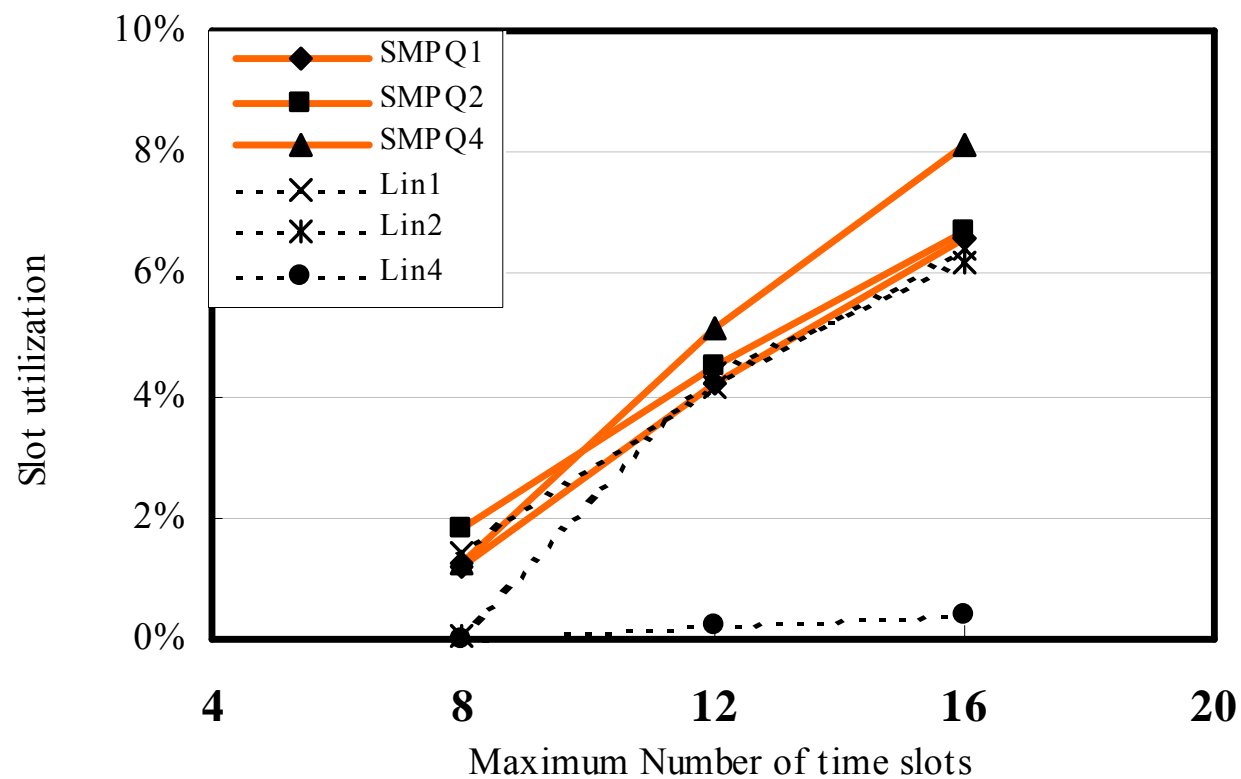
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- Effects of *Network Bandwidth*
  - Our scheme has higher SU than the Lin's
  - As  $\gamma$  increases, SU of
    - our scheme increases
    - Lin's scheme decreases
- Effects of *Maximum Number of Time slots*
  - Our scheme has higher SU than the Lin's

## Effects of *Network Bandwidth*



## Effects of *Maximum Number of Time Slots*



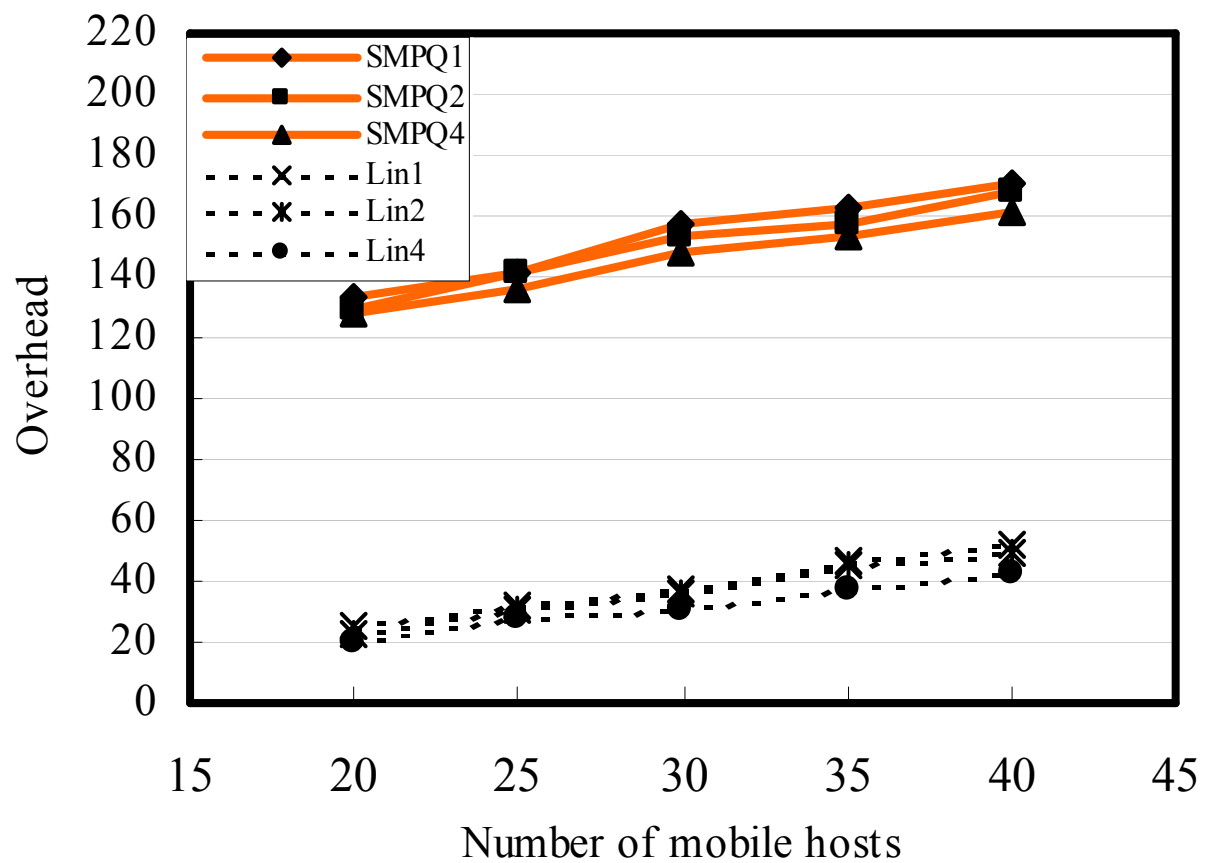


## Performance of *OverHead(OH)*

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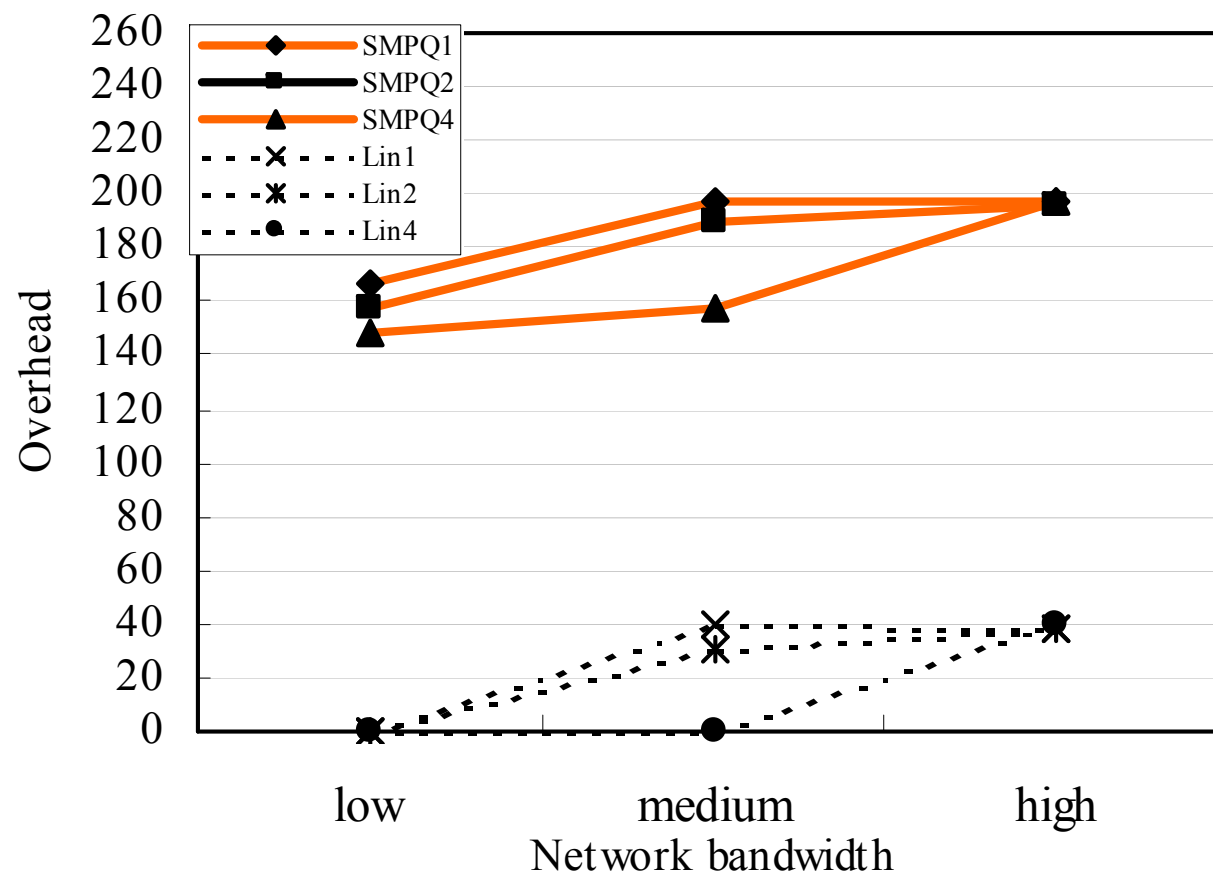
- Effects of *Number of Mobile Hosts*
  - Our scheme has higher OH than the Lin's
  - More Mobile Hosts, high OH
- Effects of *Network Bandwidth*
  - Our scheme has higher OH than the Lin's

## Effects of *Number of Mobile Hosts*





## Effects of *Network Bandwidth*





## Conclusions

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- An efficient On-demand QoS routing protocol is presented in a MANET by using
  - Spiral-path enhance the QoS route- robustness and route stability
  - Multi-path promotes the success rate of finding the QoS route