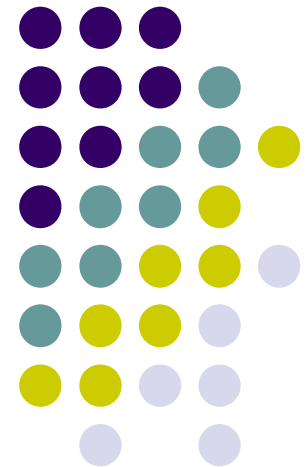


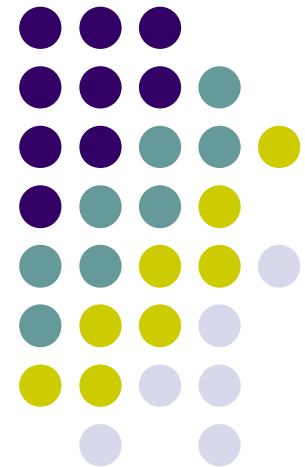
Chapter 7 : TDMA-Base MAC Protocol on VANETs



A Cross-Layer Multihop Data Delivery Protocol With Fairness Guarantees for Vehicular Networks

Gokhan Korkmaz, Eylem Ekici,
and Fusun Ozguner

*Department of Electrical and Computer Engineering,
Ohio State University*



*IEEE Transactions on Vehicular
Technology, TVT 2006*

Outline



- | Introduction
- | Controlled Vehicular Internet Access (CVIA)
- | Simulation
- | Conclusion



Introduction

- | Mobile wireless devices became the essential parts of our lives
- | “anytime, anywhere” connectivity gains a growing importance.
- | User spends hours in the traffic everyday
- | Internet access from vehicles is in great demand.

Introduction

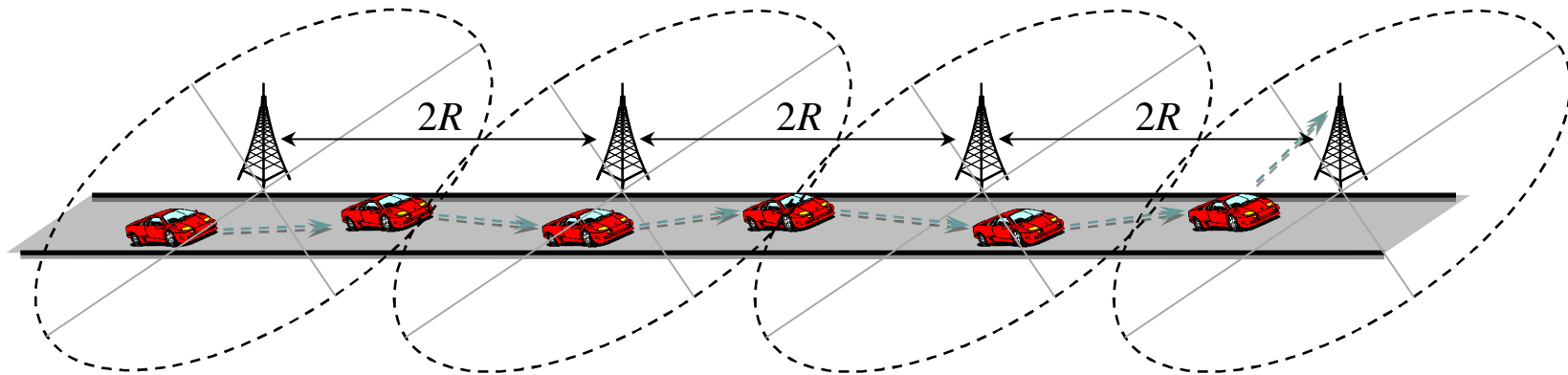


- | DSRC(Dedicated Short Range Communication) is one of the ITS(Intelligent Transport System) standards
- | DSRC systems use the IEEE 802.11 protocol as their MAC layer.



Introduction

- | A full coverage deployment of APs is costly
- | Multihopping with the IEEE 802.11 protocol suffers from several problems
 - | low throughput
 - | starvation of packets originating from vehicles far away from gateways.





Goal

- | To increase the end-to-end throughput.
- | Achieving fairness in bandwidth usage between vehicles.
- | Mitigating the hidden node problem .
- | Avoiding contention.



Network Assumption

- | Vehicular network that accesses the Internet through fixed IGWs (Internet gateways) along the road.
- | Gateways send periodic service announcements to indicate the availability of the service in their service area.
- | The uplink and the downlink packets are transmitted over two frequency-separated channels.



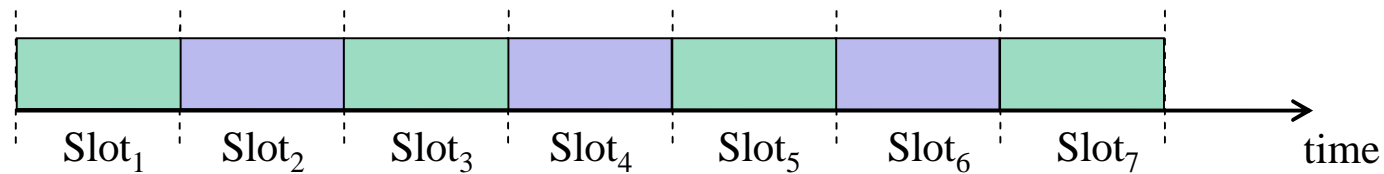
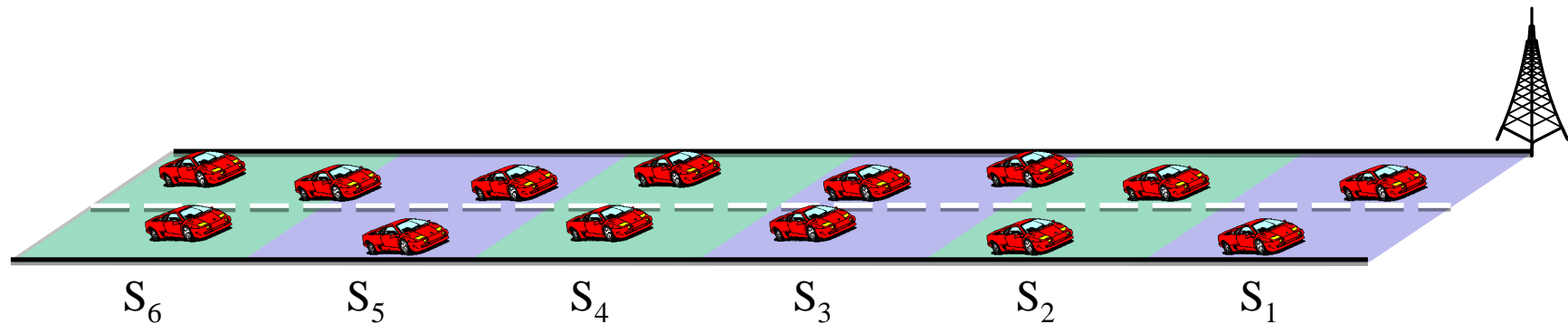
Network Assumption

- | Vehicles are equipped with GPS devices used for time synchronization and obtaining vehicle positions.
- | Vehicle positions obtained via GPS are exchanged among one-hop neighbors.



Basic Idea

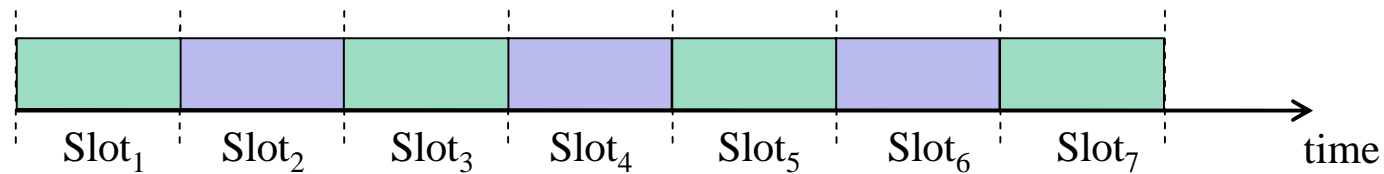
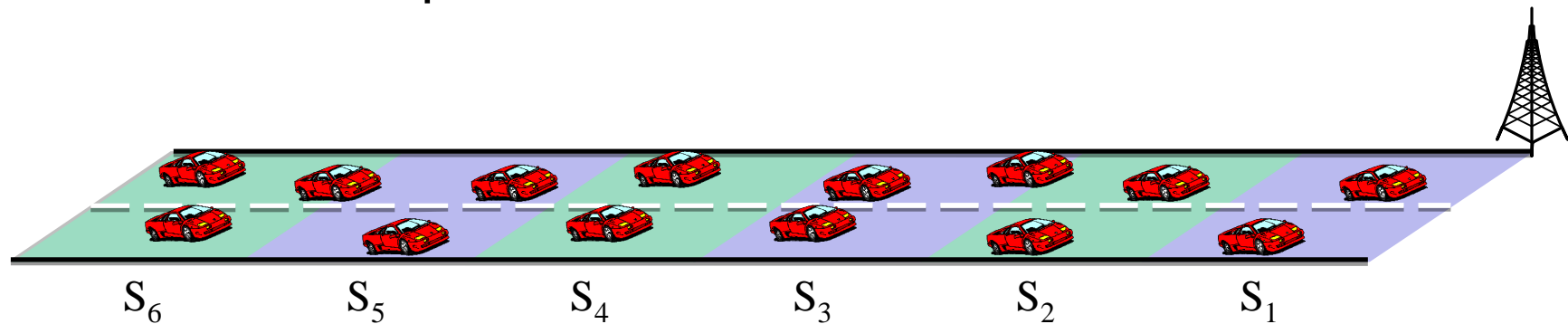
- I Divides the time into slots and the service area of the gateway into segments.



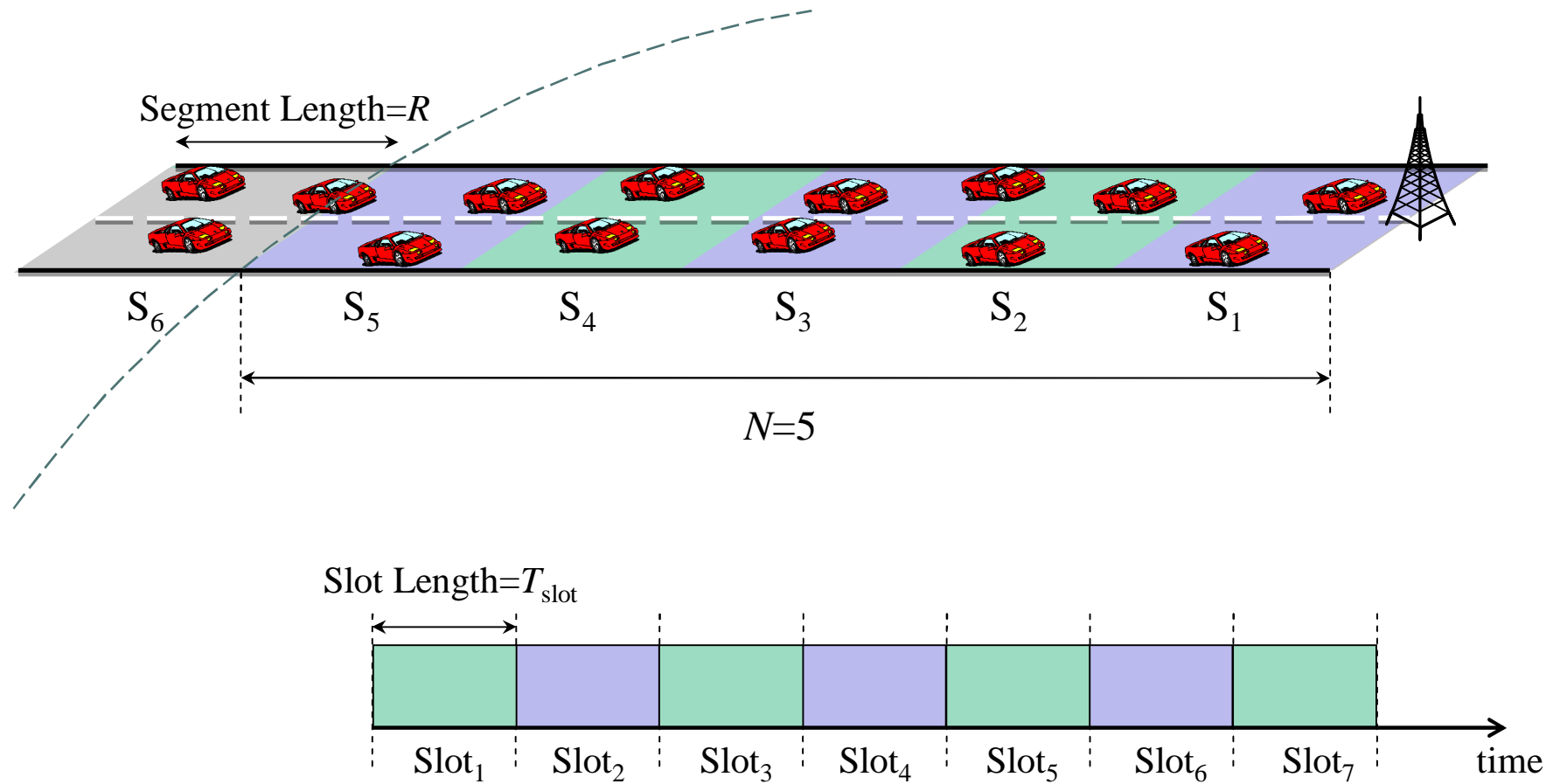


Basic Idea

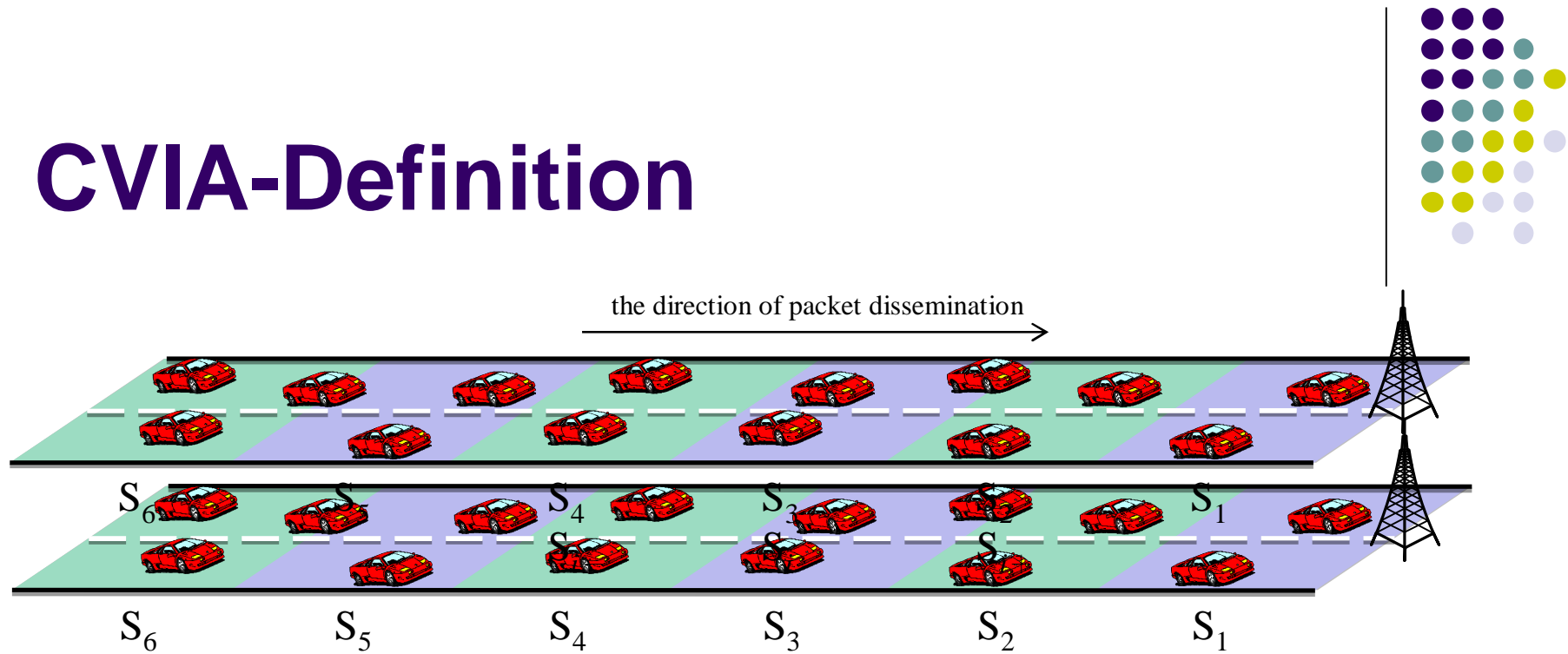
- I Controls time slots the vehicles are allowed to transmit in, how the vehicles access the channel, and to which vehicles the packets are sent.



CVIA-Definition



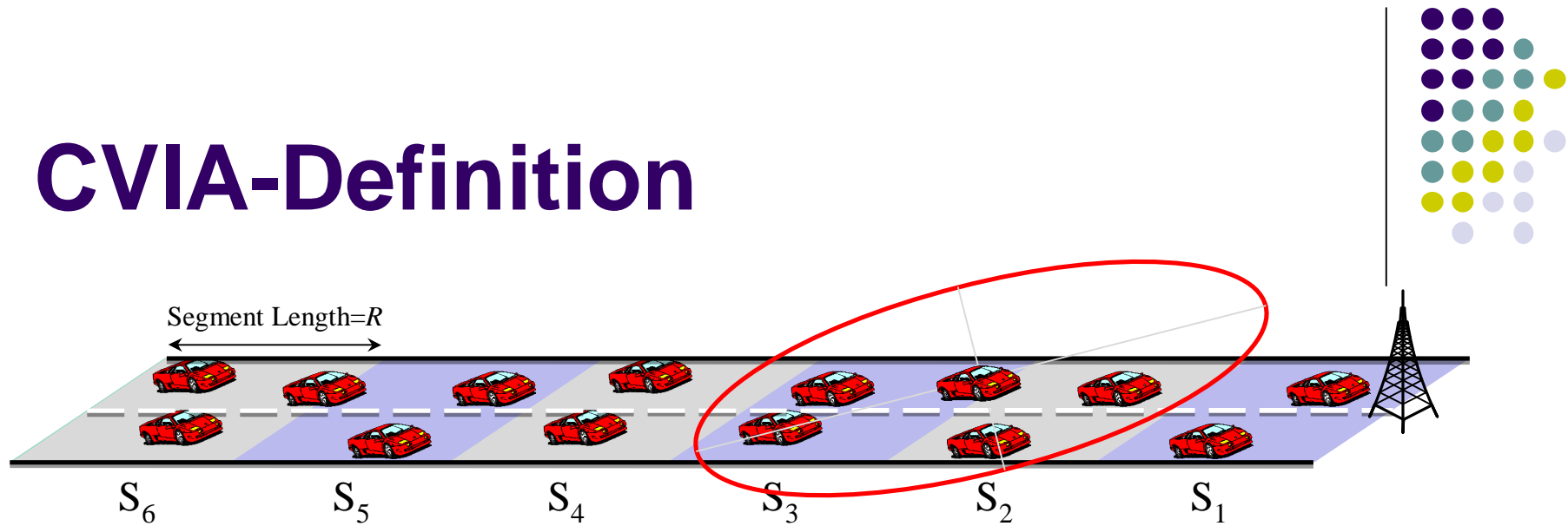
CVIA-Definition



S_{i+} : The neighboring segment in the same direction of the packet dissemination.

S_{i-} : The neighboring segment in the opposite direction of the packet dissemination.

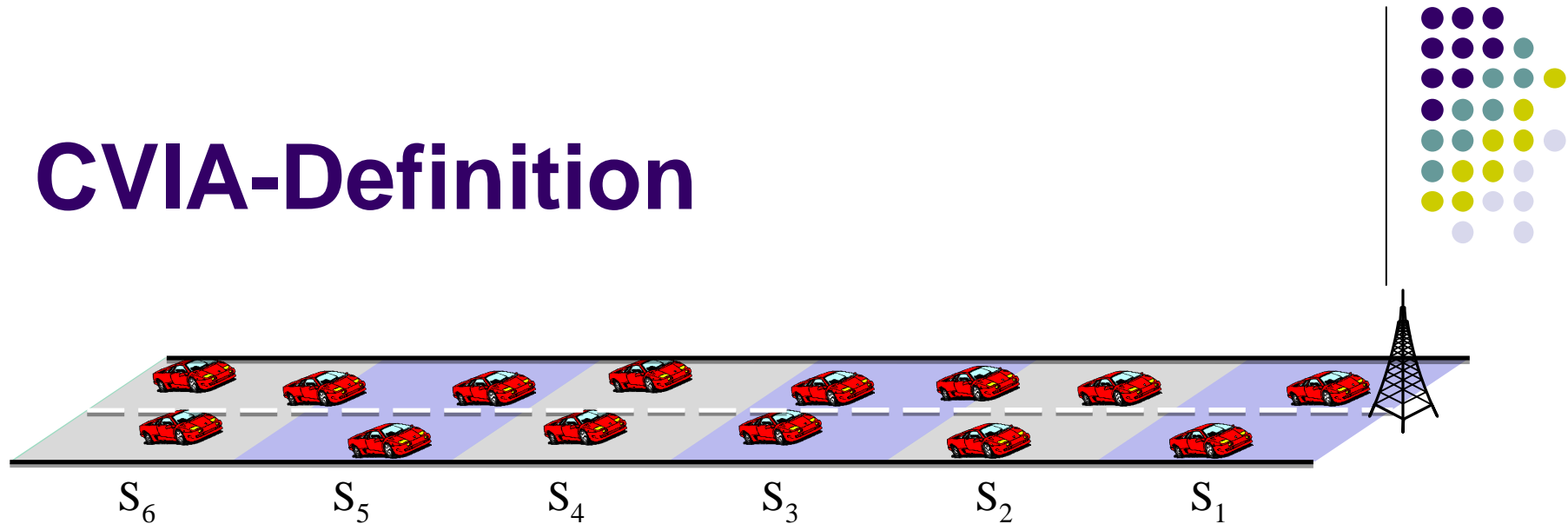
CVIA-Definition



$$\text{Interference parameter } (r) : r = \left\lceil \frac{\text{Interference Range}}{R} \right\rceil + 1$$

$$\text{Interference Range} = R \Rightarrow \left\lceil \frac{\text{Interference Range}}{R} \right\rceil = 1 \Rightarrow r = 2$$

CVIA-Definition



Interference parameter (r) : $r = \left\lceil \frac{\text{Interference Range}}{R} \right\rceil + 1$

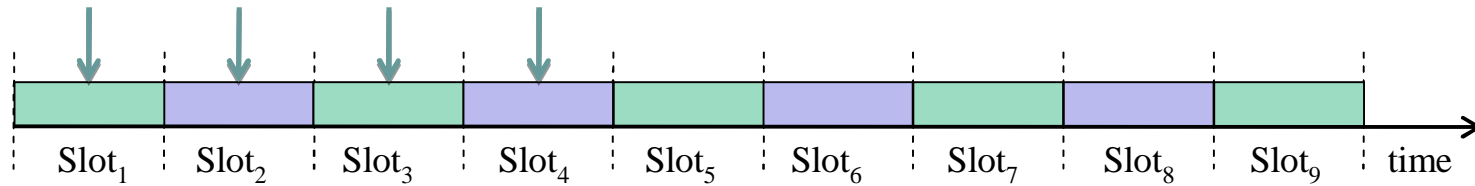
Active segment : S_i is active in TS_j if $(i \bmod r) = (j \bmod r)$

For example, $r=2$

when the current time slot is TS_5 ,

the segments S_1 , S_3 , and S_5 become active.

CVIA-Definition



$r=2$



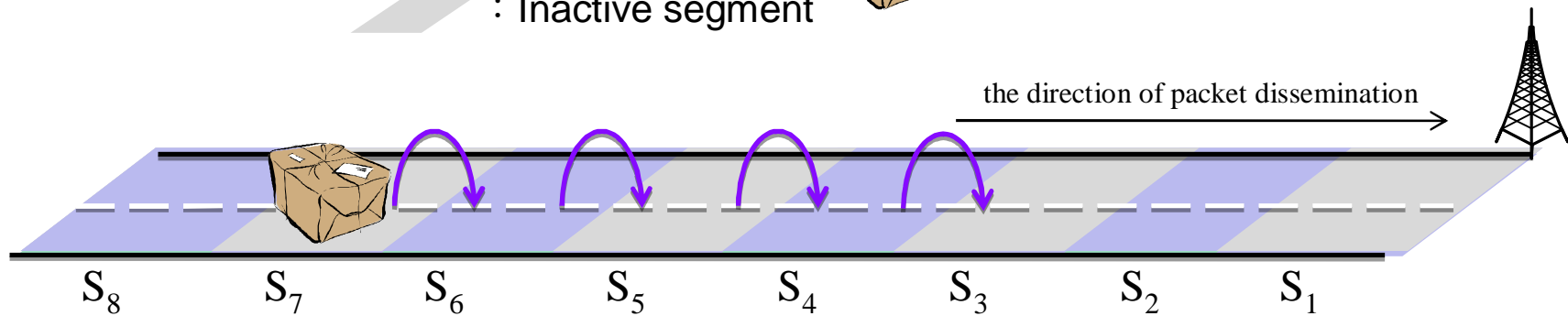
: Active segment



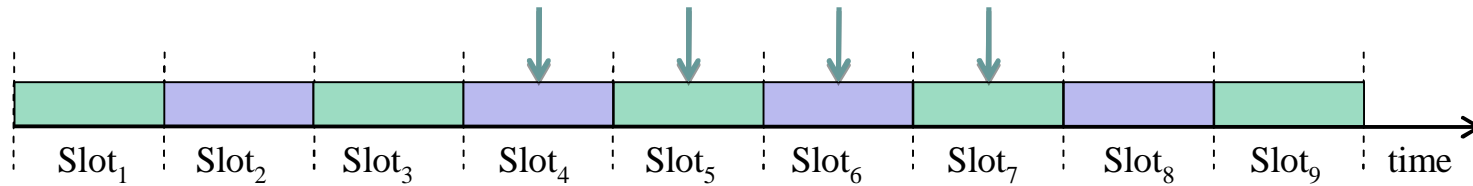
: Inactive segment



: Packets



CVIA-Definition



$r=2$



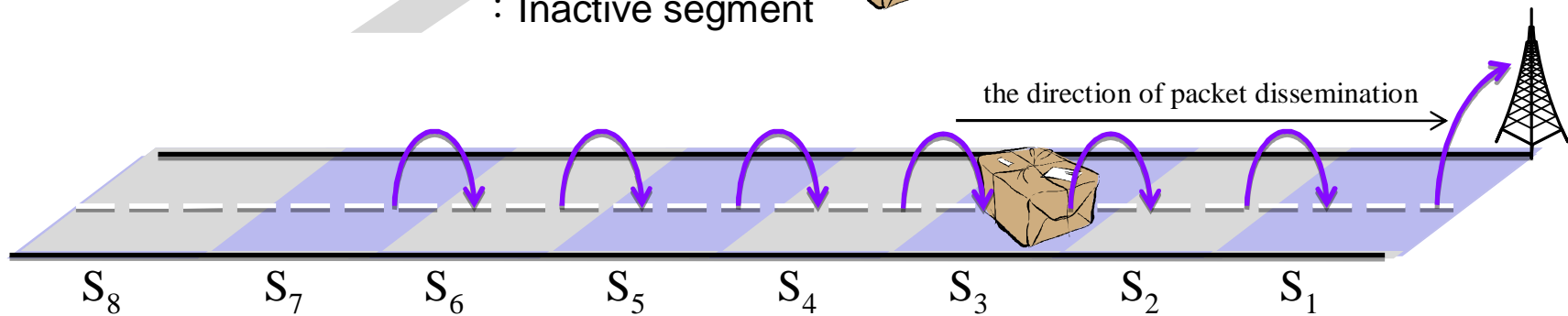
: Active segment



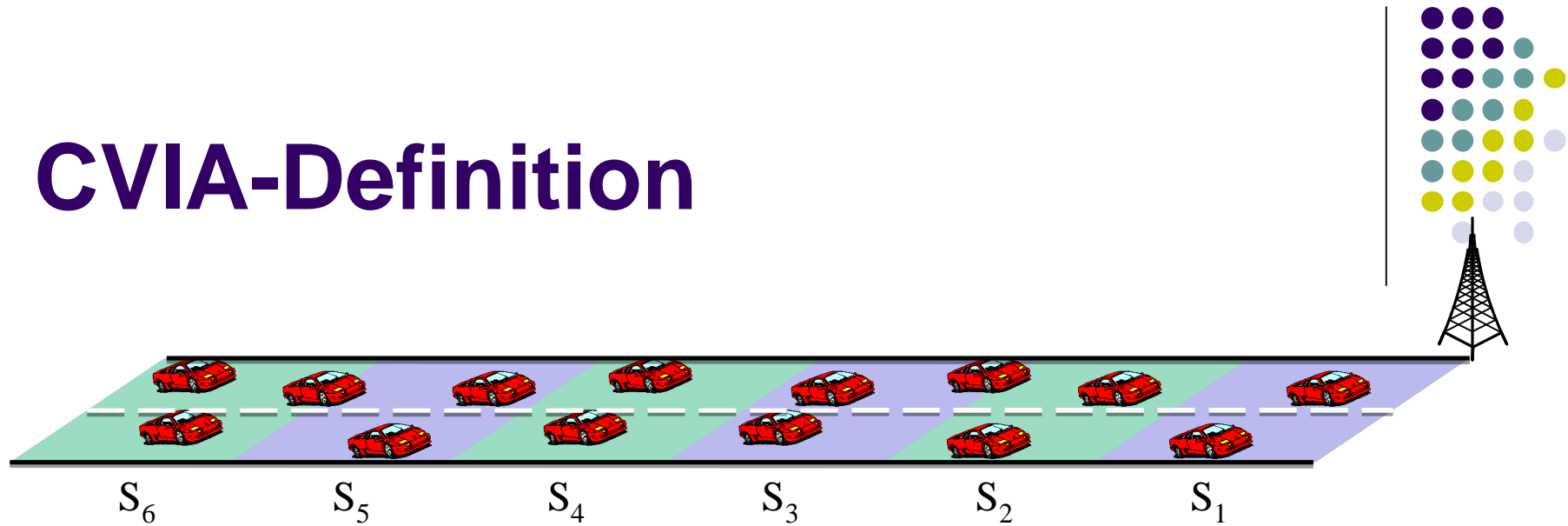
: Inactive segment



: Packets



CVIA-Definition



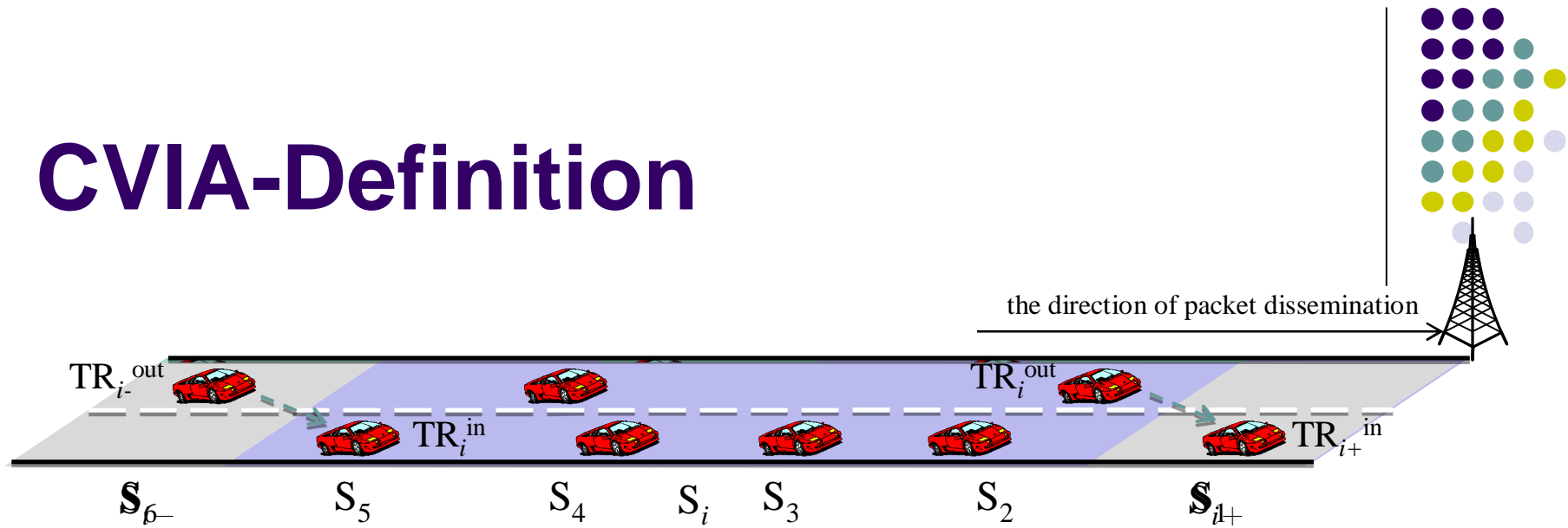
Inbound temporary router (TR_i^{in}) : The vehicle closet to S_{i-} .

All packets entering a segment go through TR_i^{in} .

Outbound temporary router (TR_i^{out}) : The vehicle closet to S_{i+} .

All packets leaving the segment go through TR_i^{out} .

CVIA-Definition

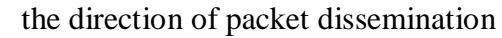


Inbound temporary router (TR_i^{in}) : The vehicle closet to S_{i-} .

All packets entering a segment go through TR_i^{in} .

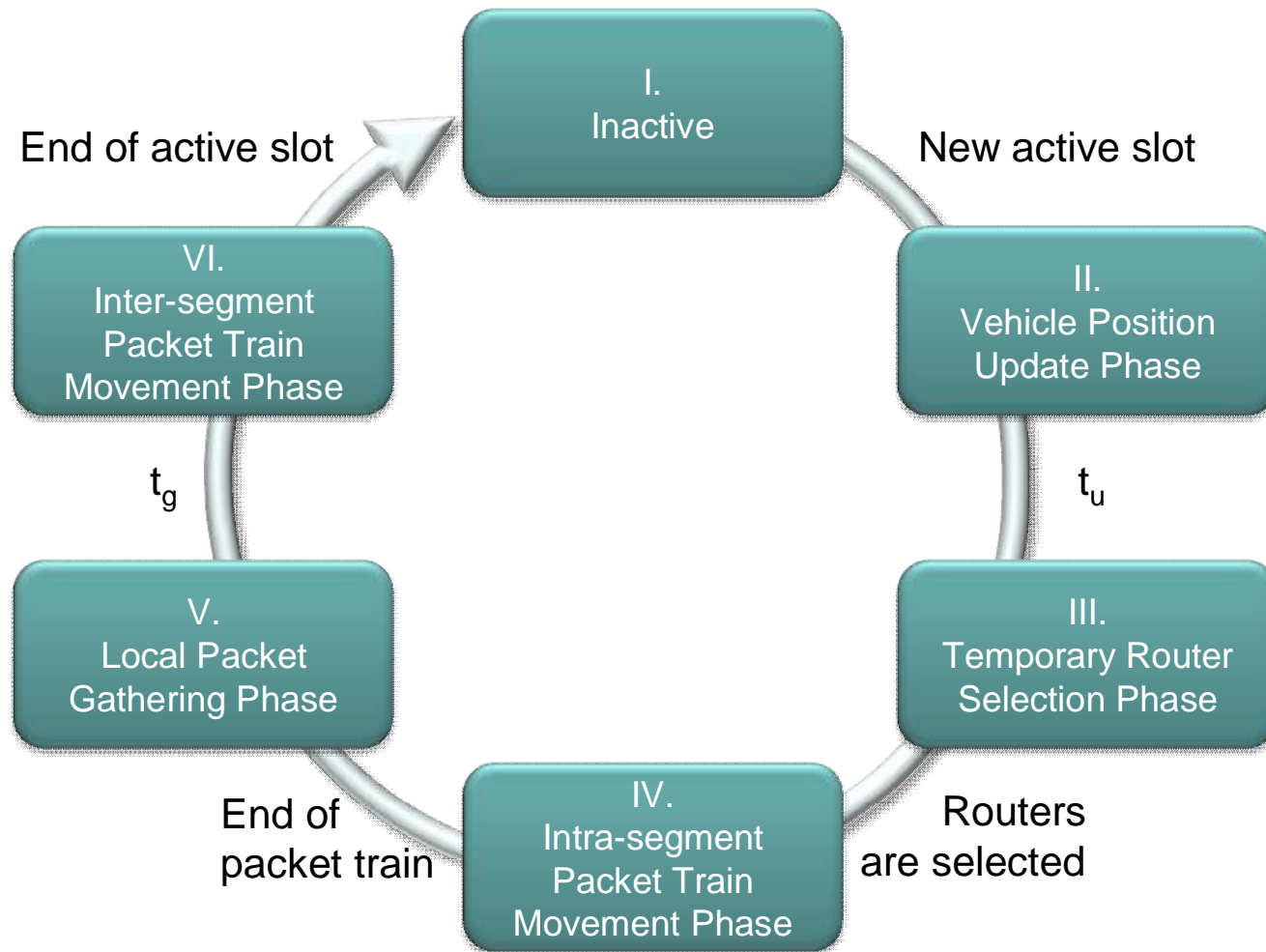
Outbound temporary router (TR_i^{out}) : The vehicle closet to S_{i+} .

All packets leaving the segment go through TR_i^{out} .

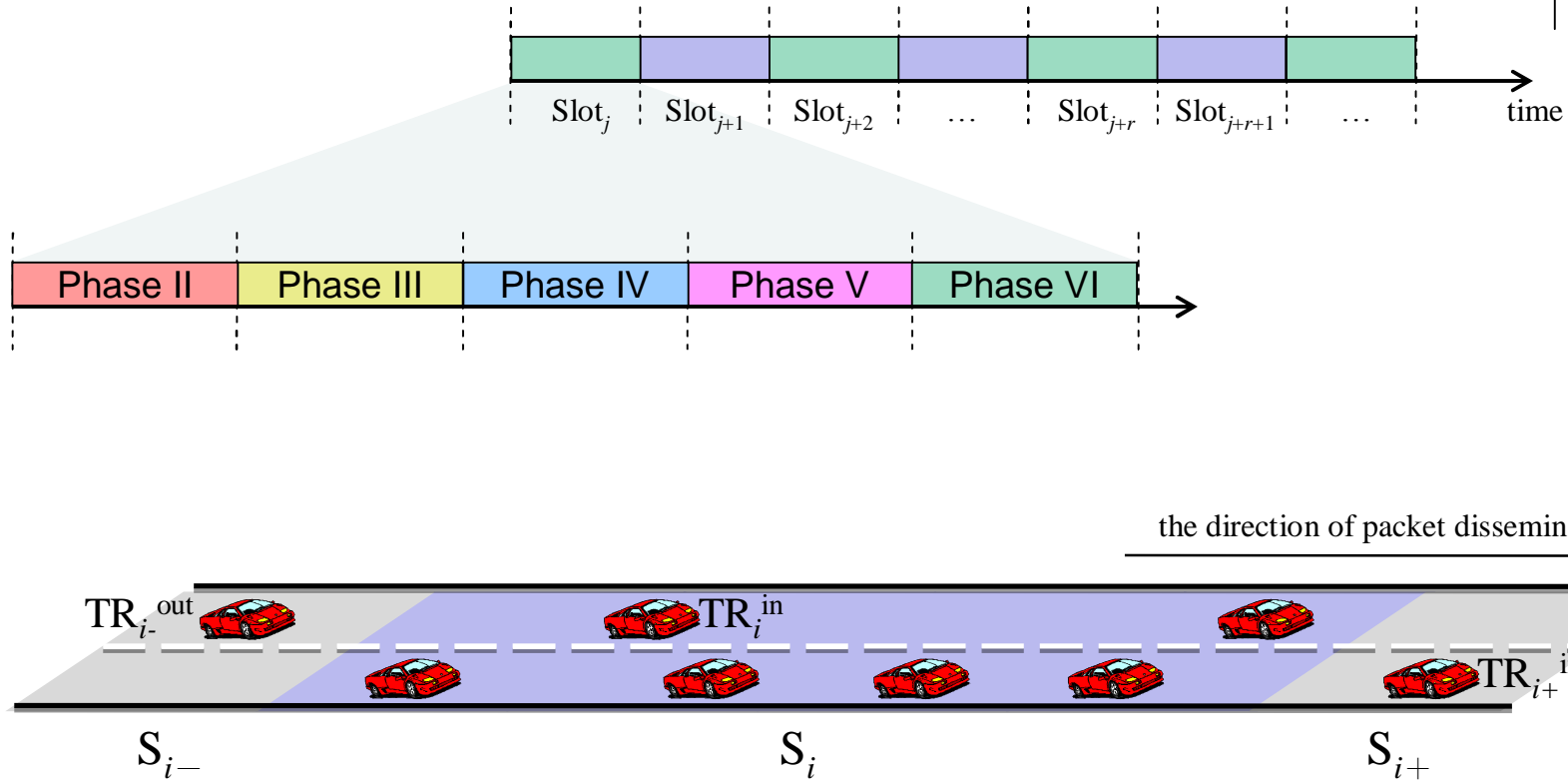


1. TR_i^{in} deliver packet train to TR_i^{out} .
2. TR_i^{out} gather local packets.
3. TR_i^{out} move out packet train to TR_{i+}^{out} .

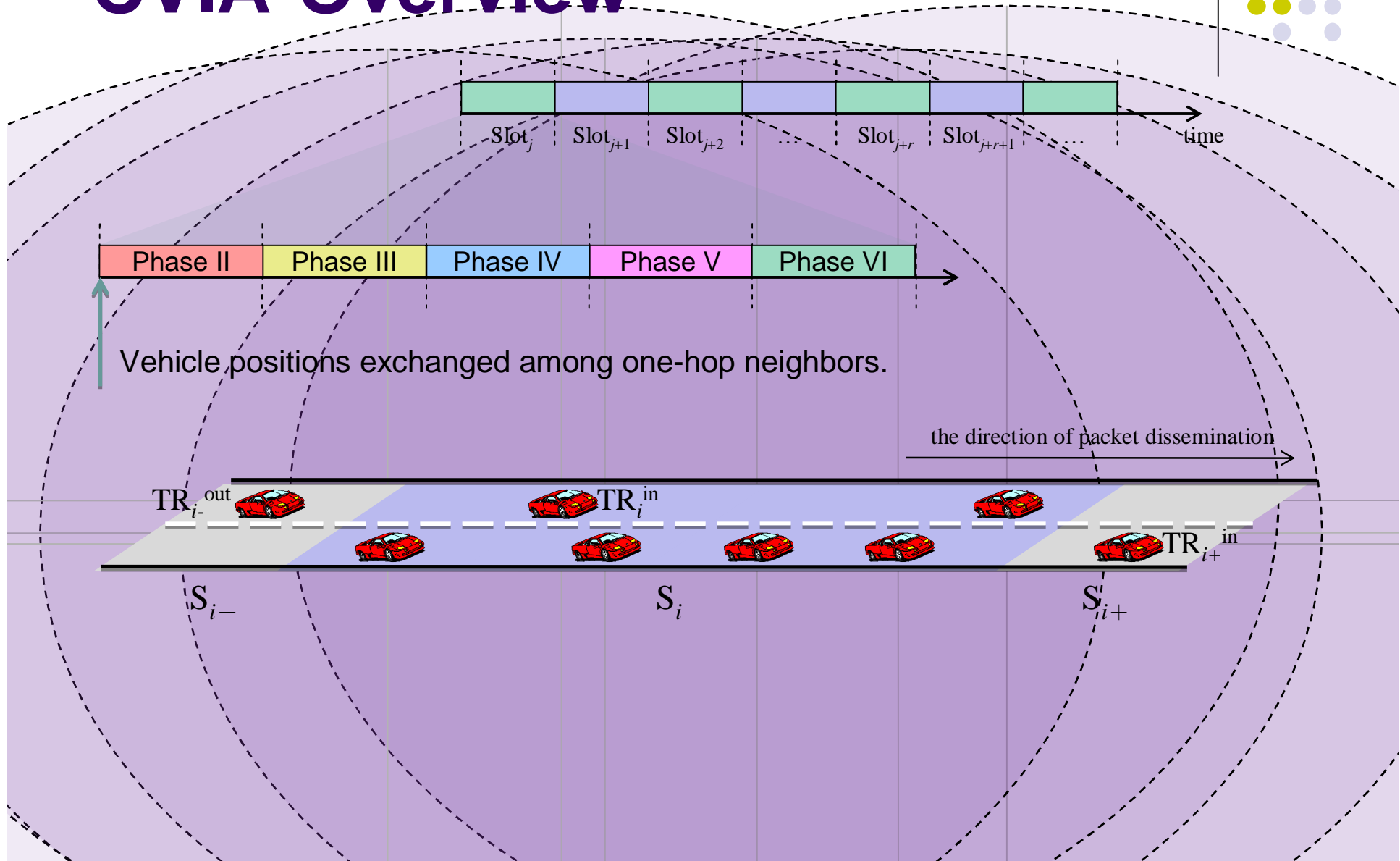
CVIA-Overview



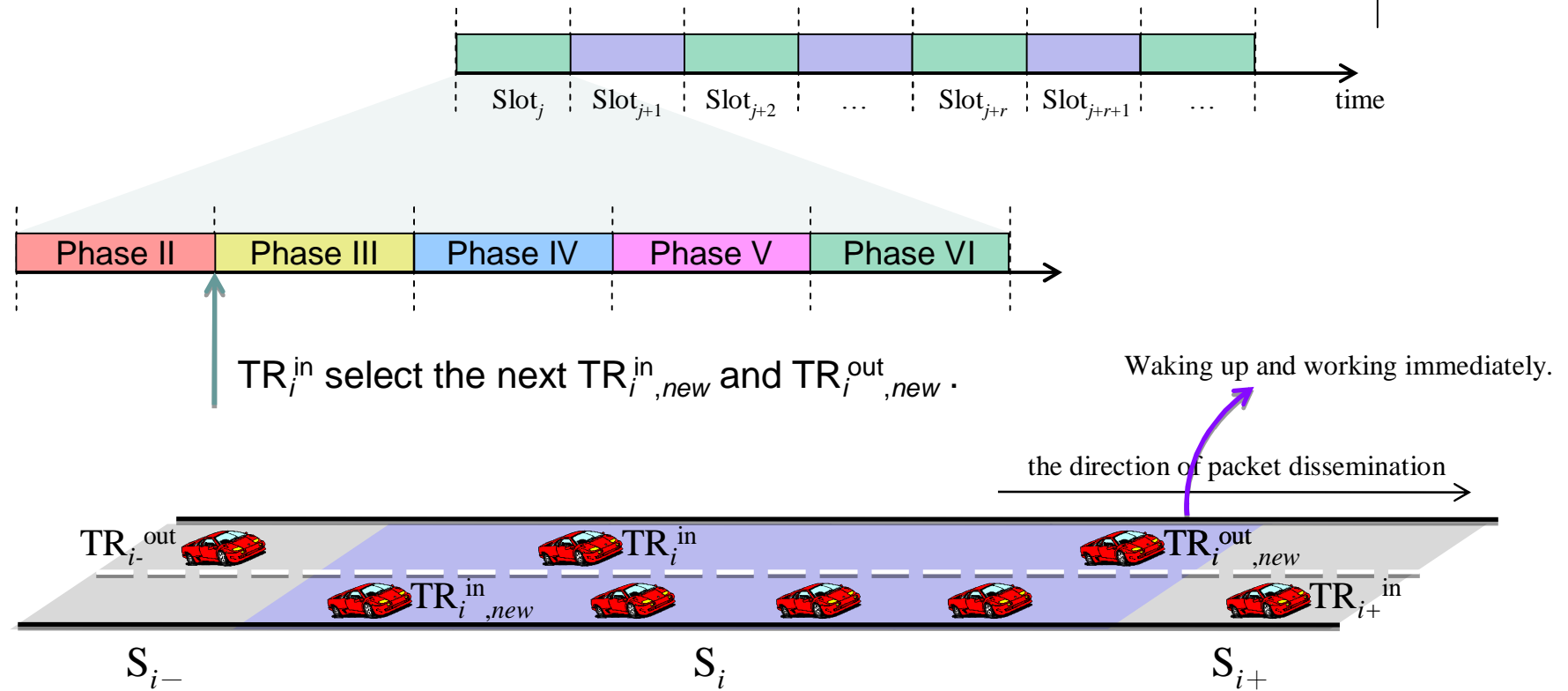
CVIA-Overview



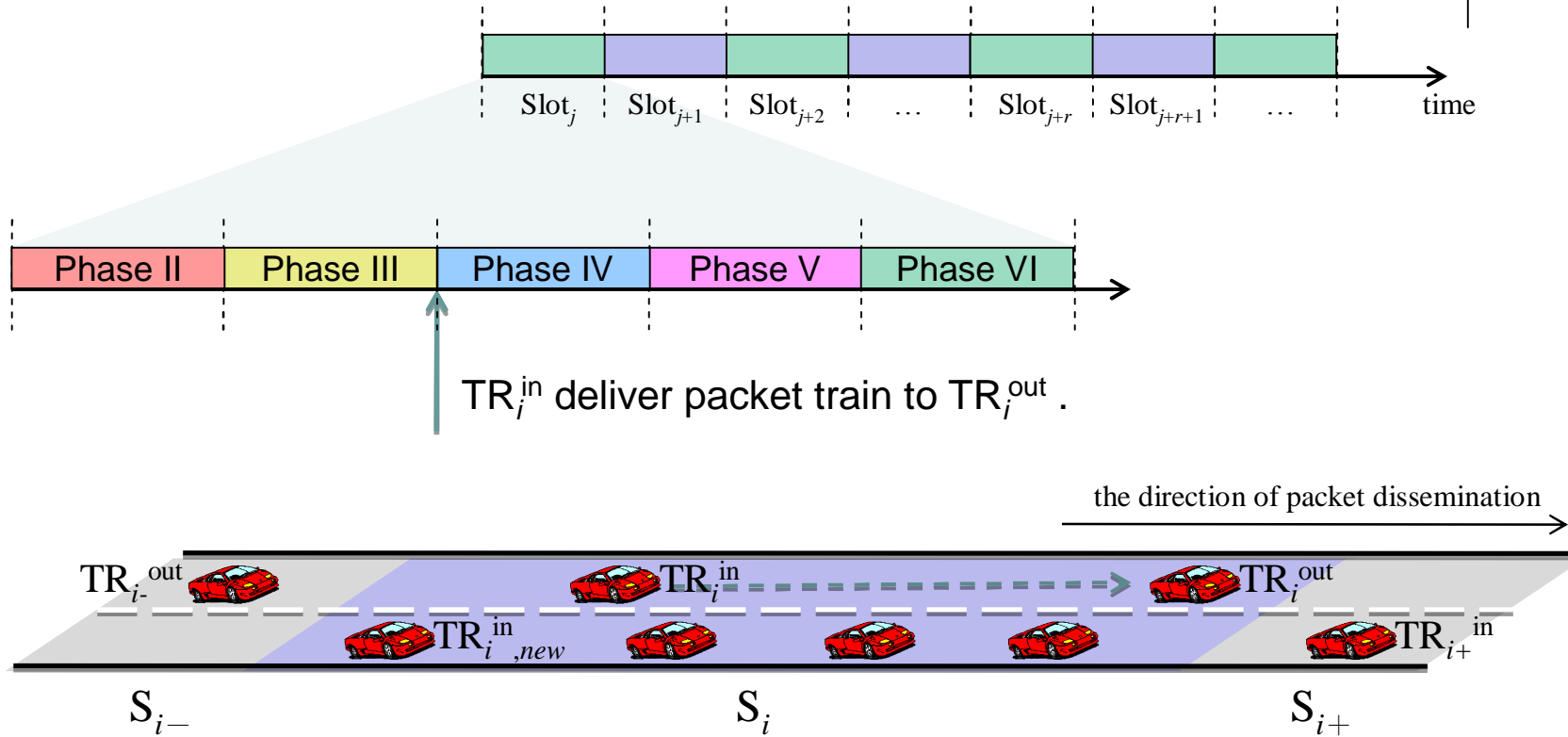
CVIA-Overview



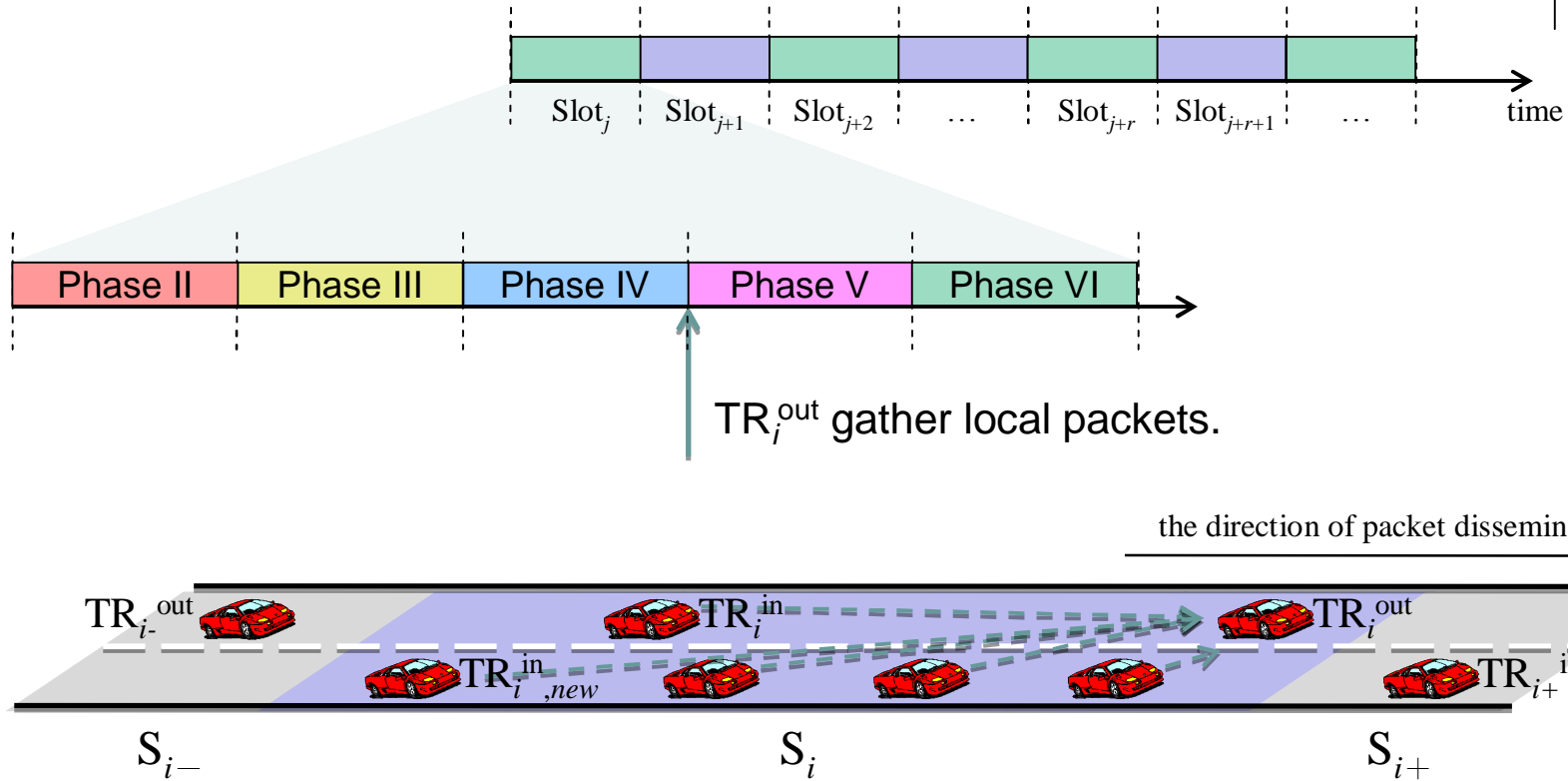
CVIA-Overview



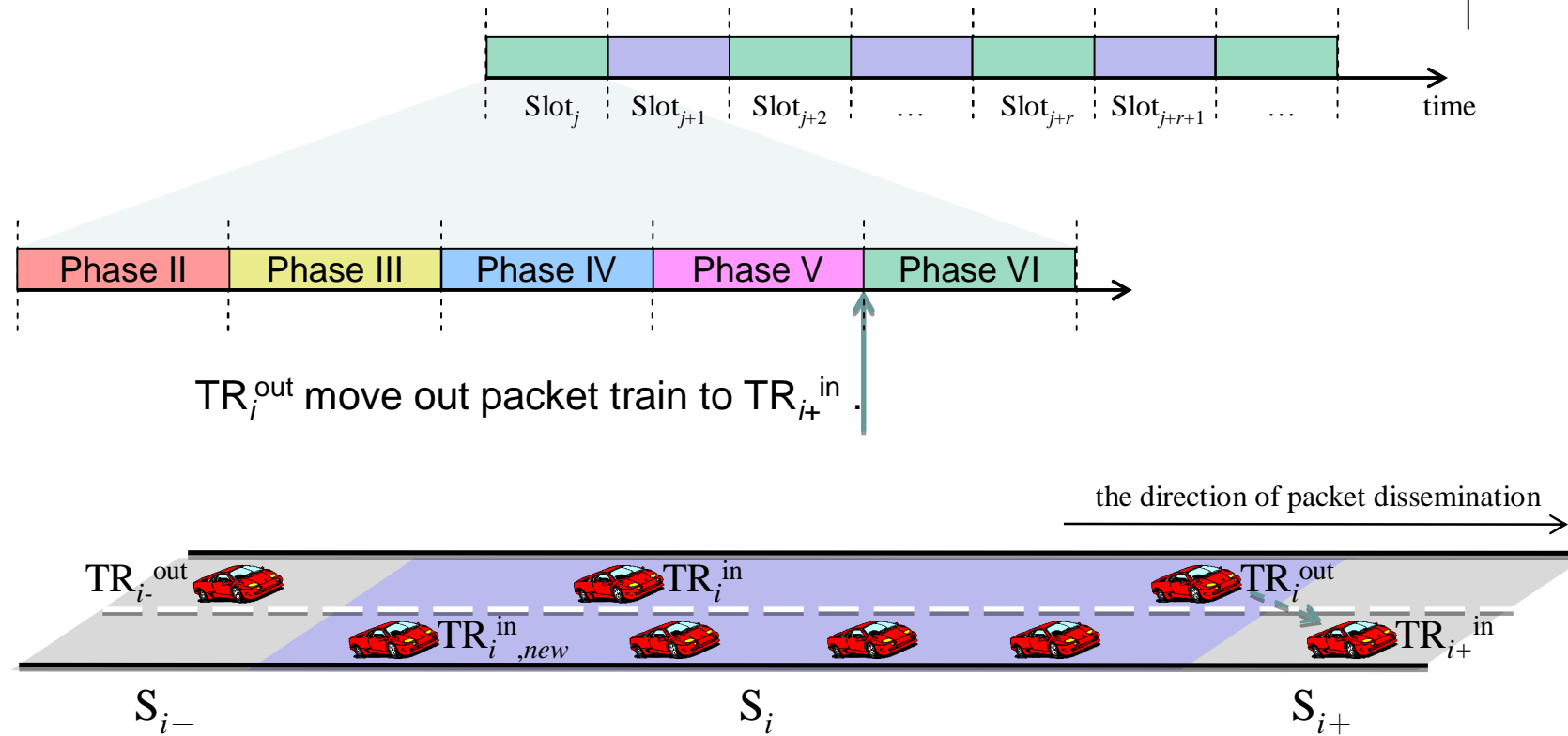
CVIA-Overview



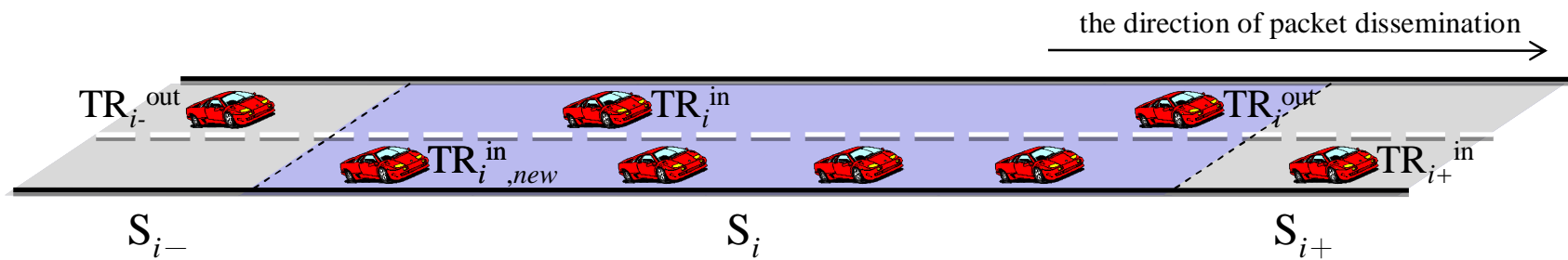
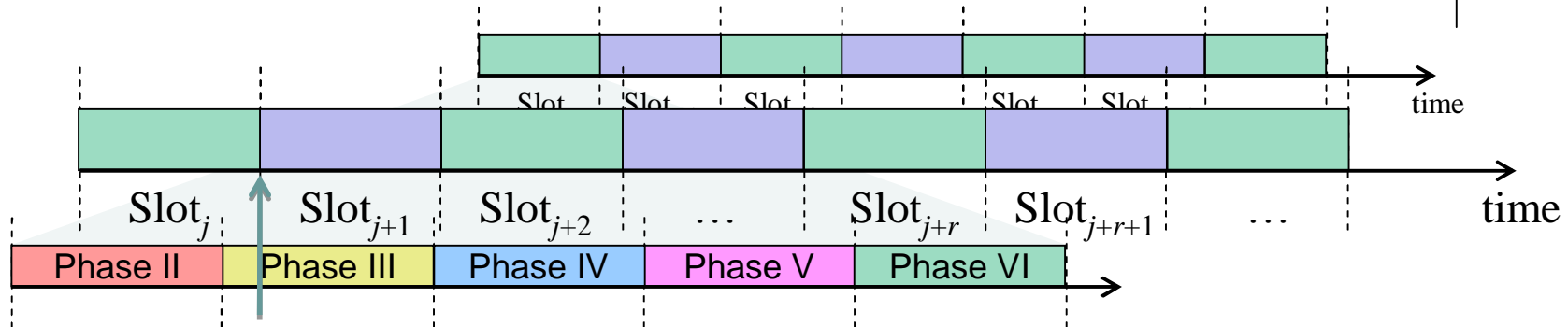
CVIA-Overview



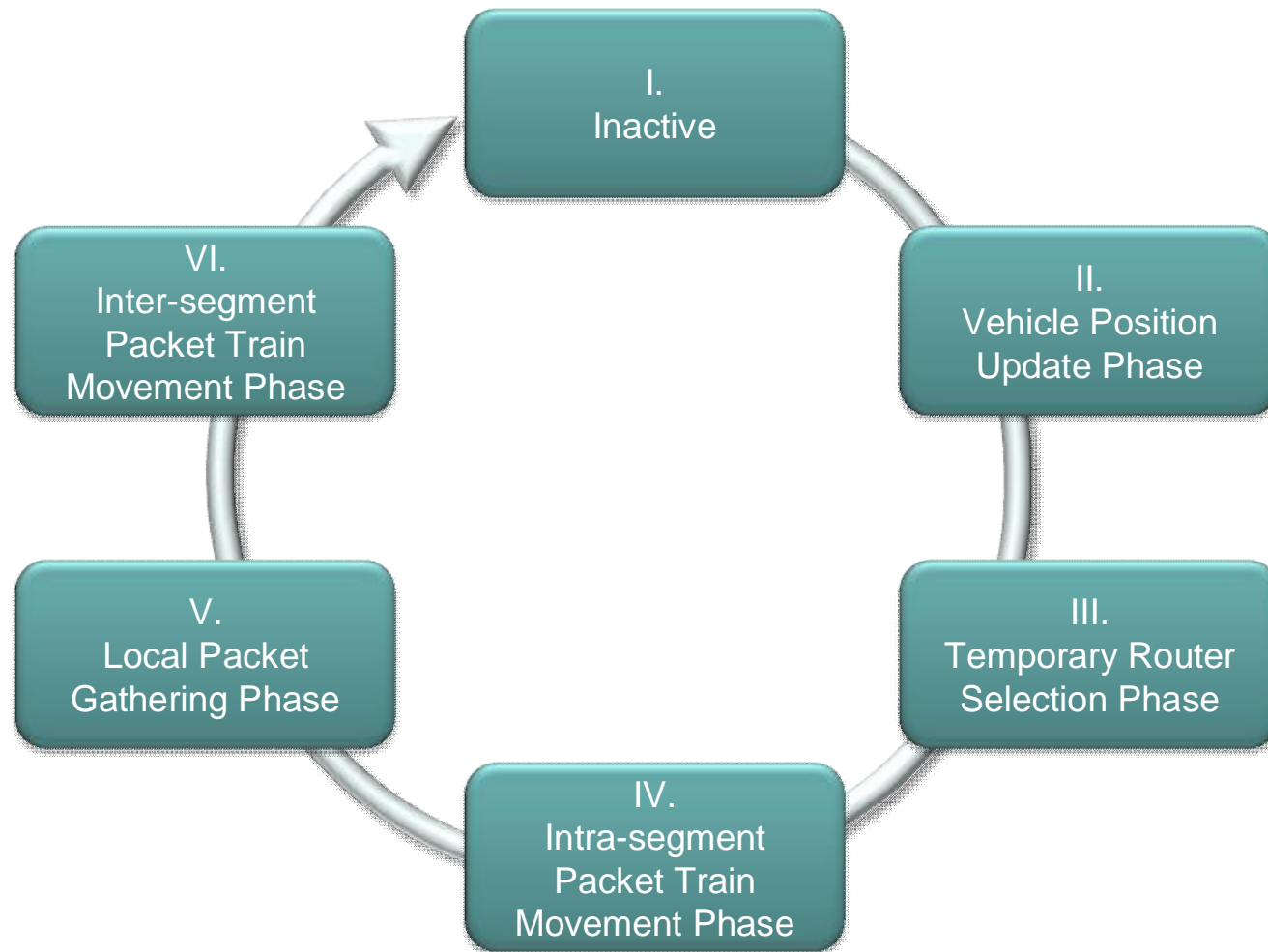
CVIA-Overview



CVIA-Overview



CVIA



CVIA - Inactive



Vehicles in inactive segments do not access the channel.

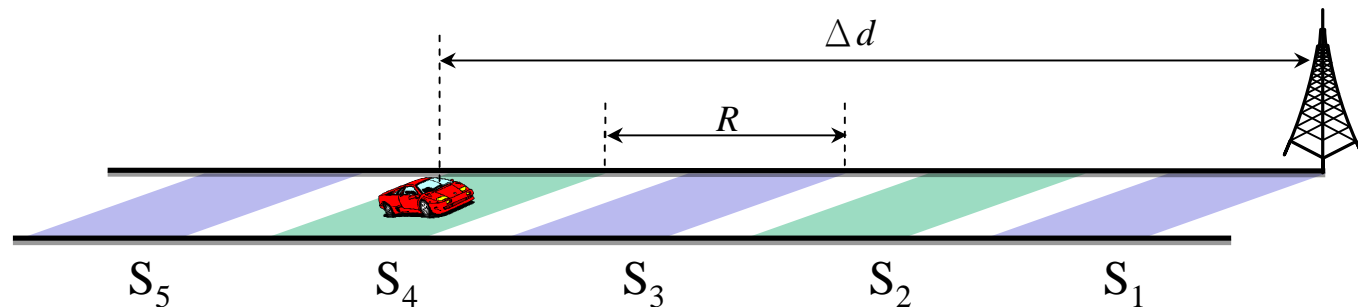
S_i is active in TS_j if $(i \bmod r) = (j \bmod r)$

$i = \left\lceil \frac{\Delta d}{R} \right\rceil$, Δd : The distance of the vehicle to the gateway.

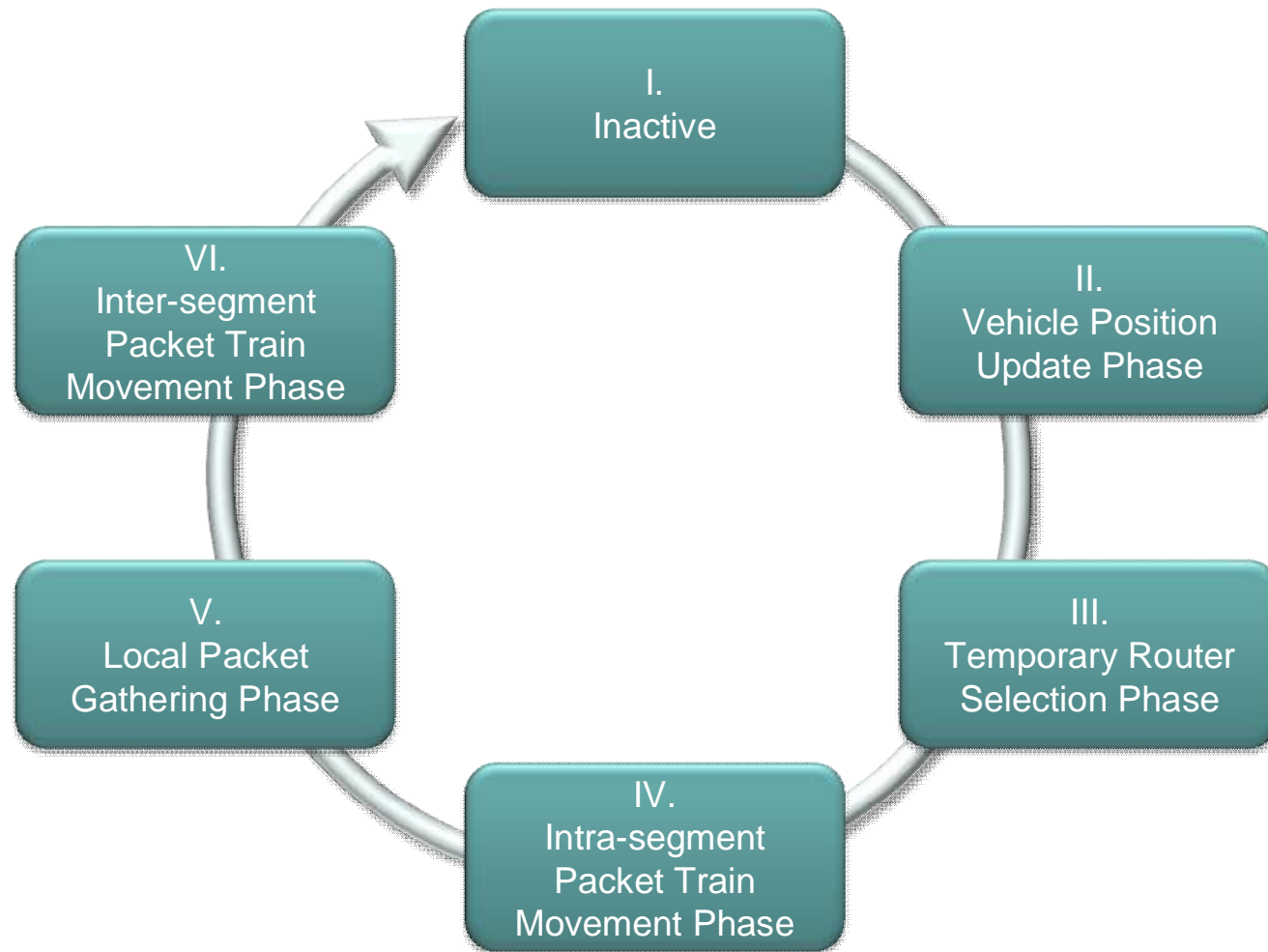
$j = \left\lceil \frac{\Delta t}{T_{slot}} \right\rceil$, Δt : The time passed since an absolute reference point.

Vehicles obtain their positions and synchronize their clocks using GPS.

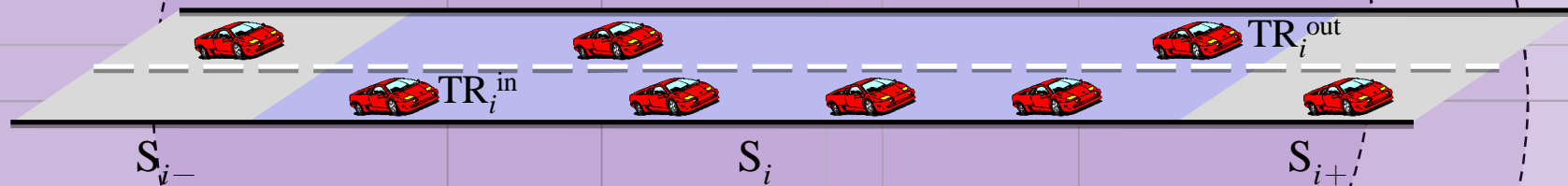
Vehicles learn the positions of gateways from a digital road map database or the service announcement packets broadcast periodically by gateways.



CVIA



CVIA - Vehicle Position Update Phase



t_u : A time interval reserved for position update packets (PUPs).

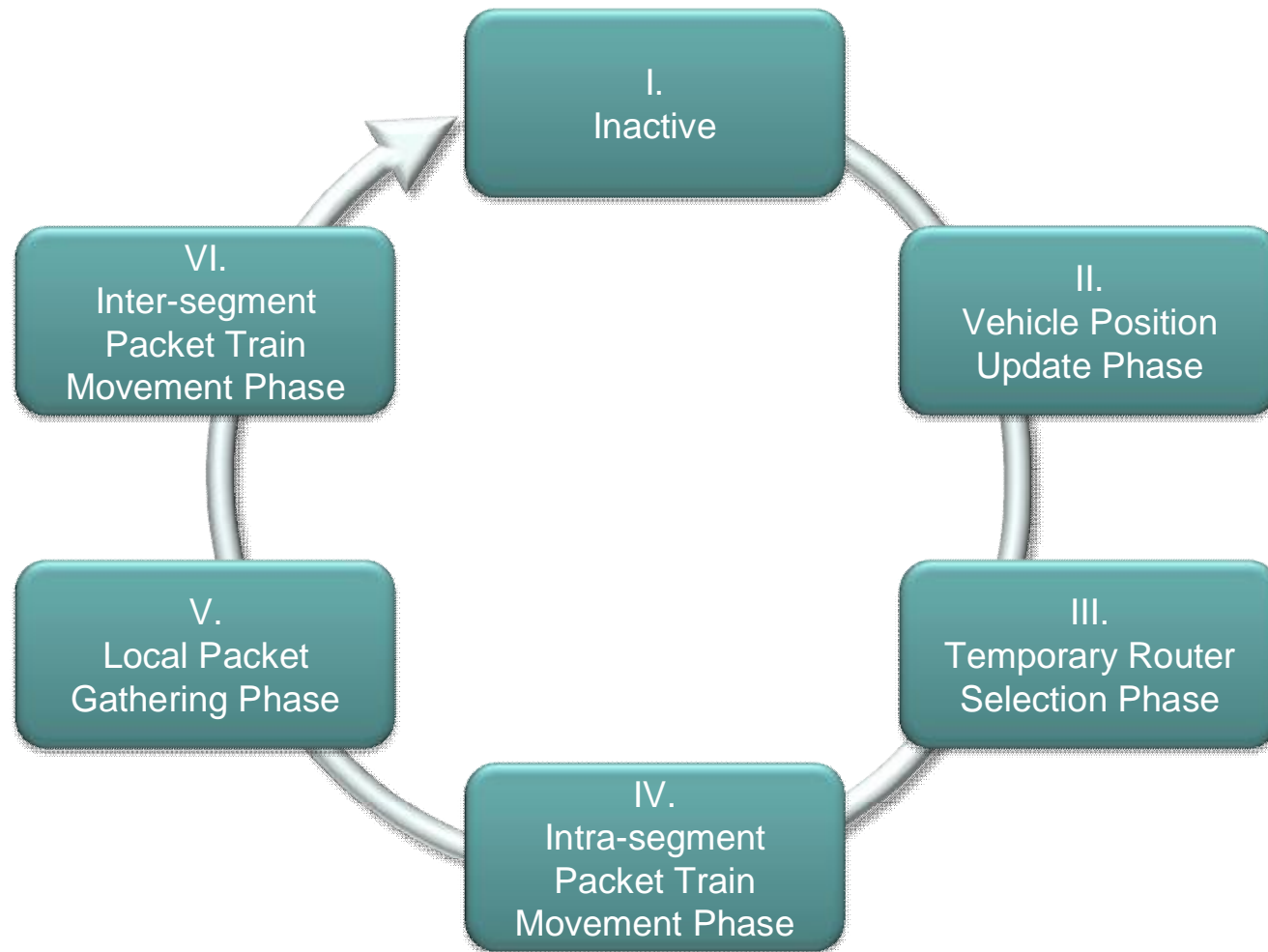
t_{PUP} : The duration of a PUP.

Vehicles pick a random waiting time (RWT) from $[0, t_u - t_{PUP})$.

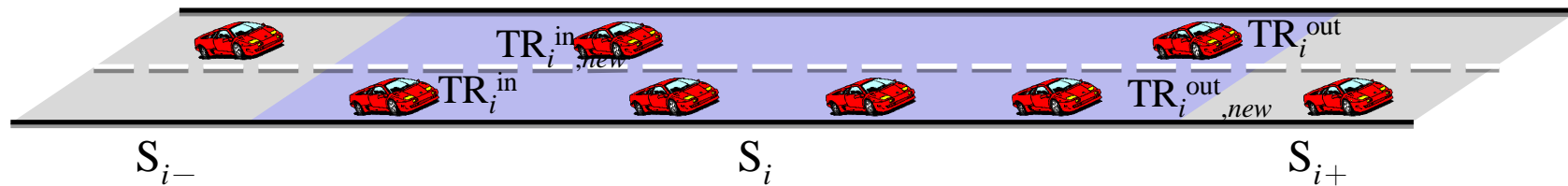
After an RWT, vehicles access the channel using DCF method of the IEEE 802.11 protocol.

The length of a PUP is short, RTS/CTS handshake is not used before sending PUP.

CVIA



CVIA - Temporary Router Selection Phase



New TR_i^{in} and TR_i^{out} are selected by TR_i^{in} .

The selected routers are called $TR_{i,new}^{in}$ and $TR_{i,new}^{out}$ until they become active.

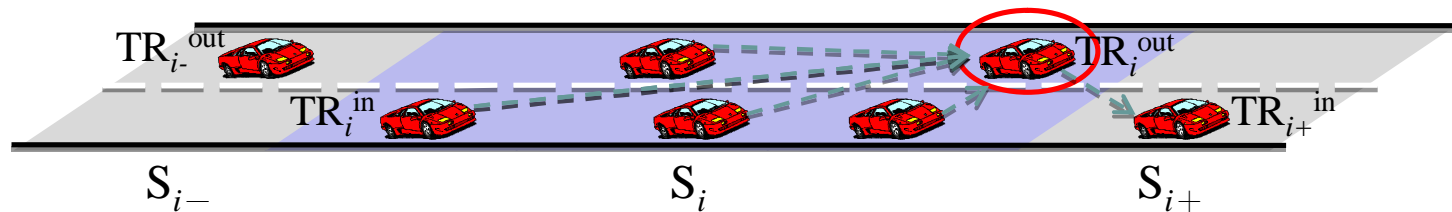
Use “router lifetime” and “safe area” concept.



CVIA - Temporary Router Selection Phase

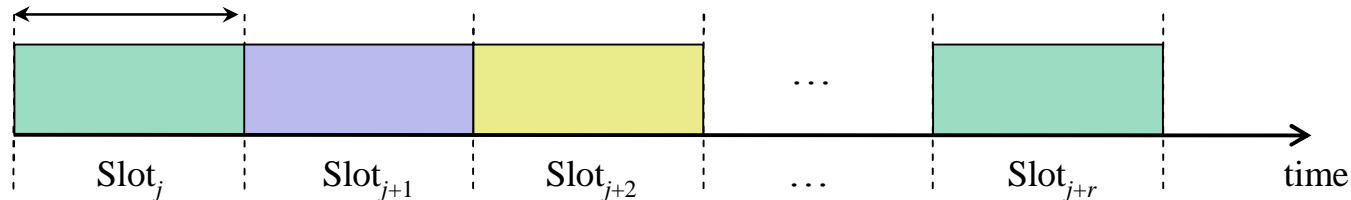
TR_i^{out} is responsible for

1. Receiving packet train relayed by TR_i^{in} .
2. Gathering local packets.
3. Creating a new packet train and sending this train out of S_i to TR_{i+}^{in} .



TR_i^{out} becomes active immediately and stays active until the end of TS_j .

lifetime of TR_i^{out} is T_{slot}



$$x_{margin+,v} = (\Delta t_{pu,v} + 1 \cdot T_{slot}) \cdot V_{max}$$

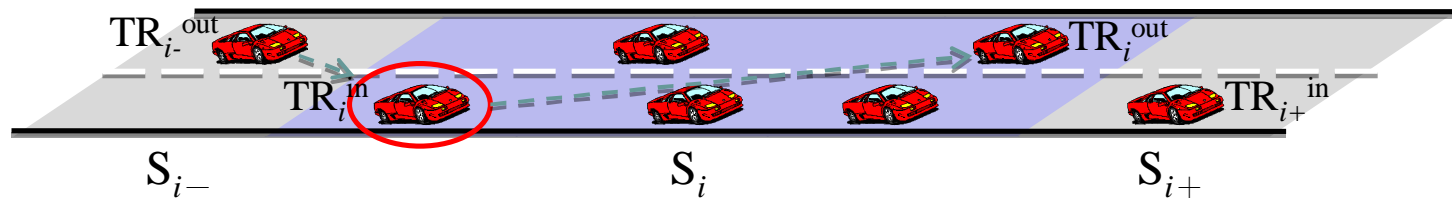
$\Delta t_{pu,v}$: The elapsed time since the last position update from vehicle v .

CVIA - Temporary Router Selection Phase

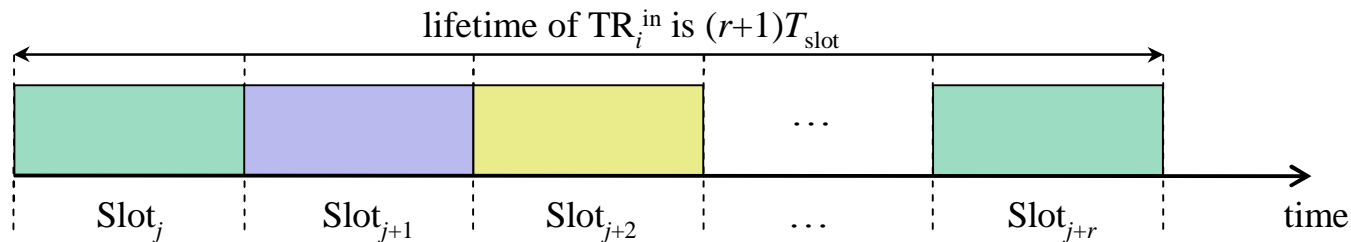


TR_i^{in} is responsible for

1. Receiving packet train coming from TR_{i-}^{out} in TS_{j+1} .
2. Selecting and announcing $TR_{i,new}^{in}$ and $TR_{i,new}^{out}$ at the beginning of TS_{j+r} .
3. Relaying the packet train to TR_i^{out} in TS_{j+r} (next active slot).



$TR_{i,new}^{in}$ becomes active throughout TS_{j+1} and TS_{j+r} .



$$x_{\text{margin-},v} = (\Delta t_{pu,v} + (r+1) \cdot T_{\text{slot}}) \cdot V_{\text{max}}$$

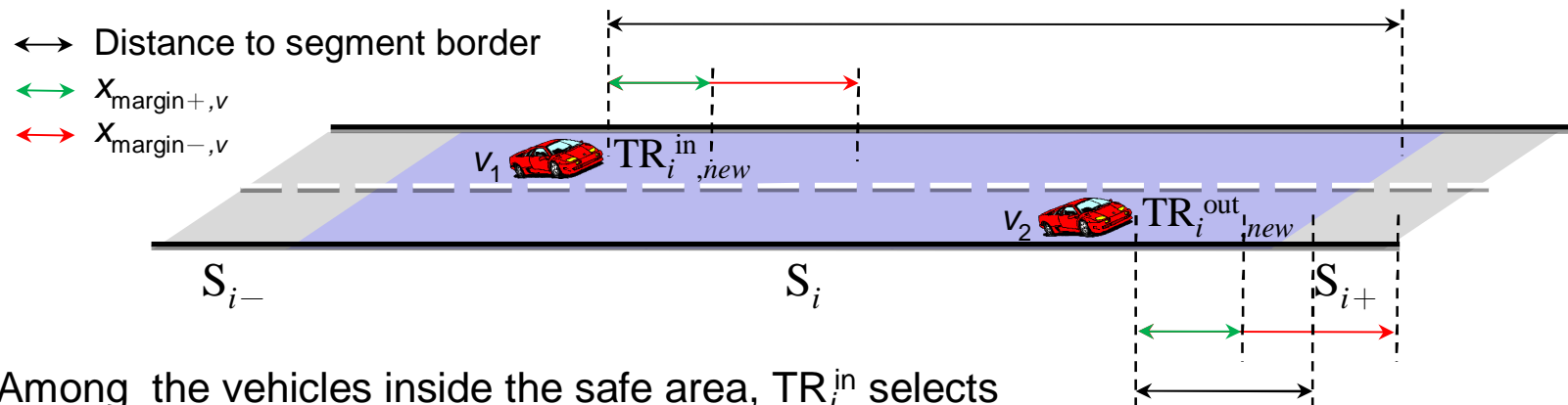
$\Delta t_{pu,v}$: The elapsed time since the last position update from vehicle v .



CVIA - Temporary Router Selection Phase

If vehicle v 's distance to segment borders is more than $x_{\text{margin}+,v}$,
the vehicle can safely be selected as $\text{TR}_{i,\text{new}}^{\text{out}}$.

If vehicle v 's distance to segment borders is more than $x_{\text{margin}-,v}$,
the vehicle can safely be selected as $\text{TR}_{i,\text{new}}^{\text{in}}$.



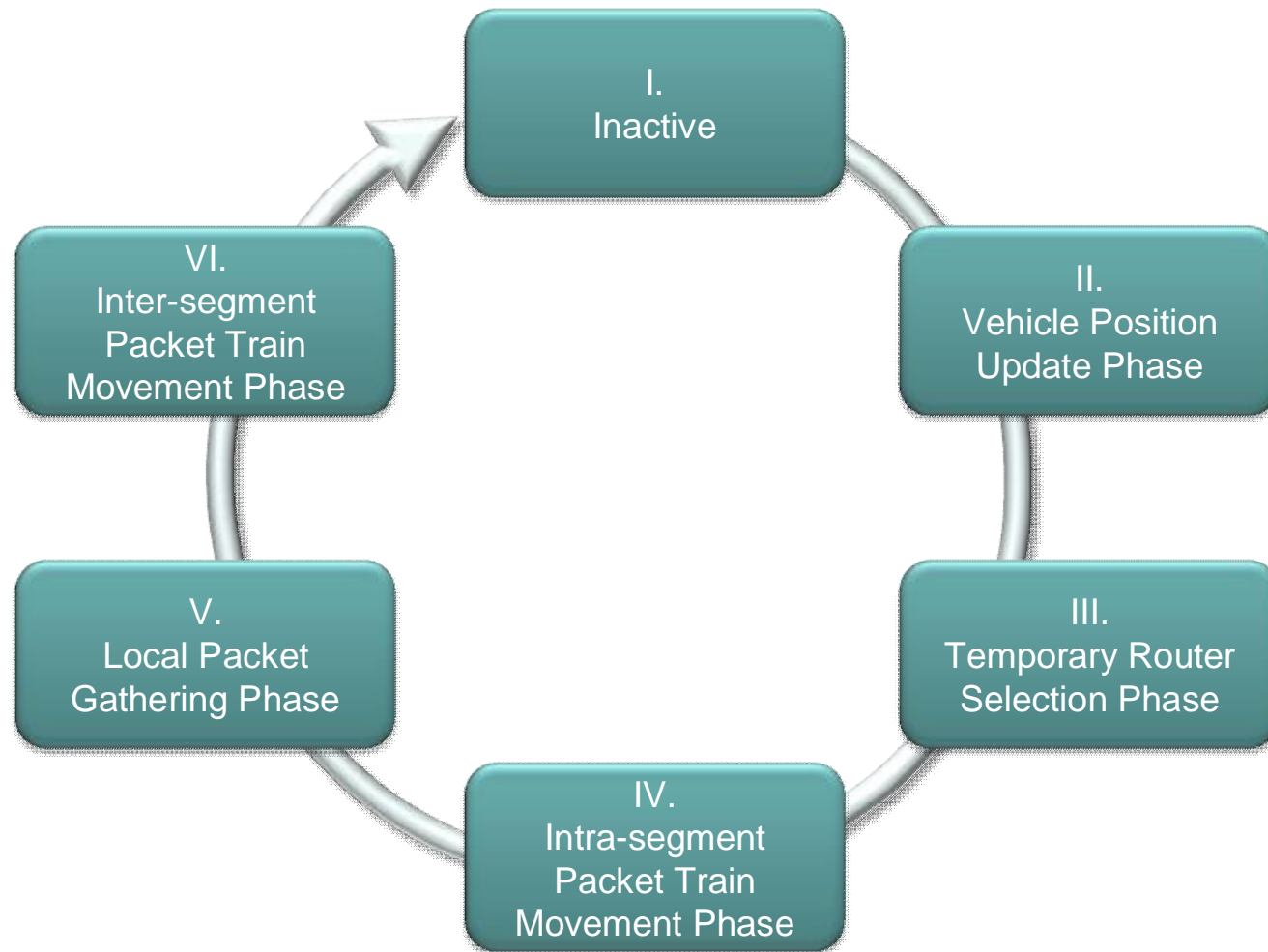
Among the vehicles inside the safe area, $\text{TR}_{i,\text{new}}^{\text{in}}$ selects

the vehicle closest to S_{i+} as $\text{TR}_{i,\text{new}}^{\text{out}}$ and

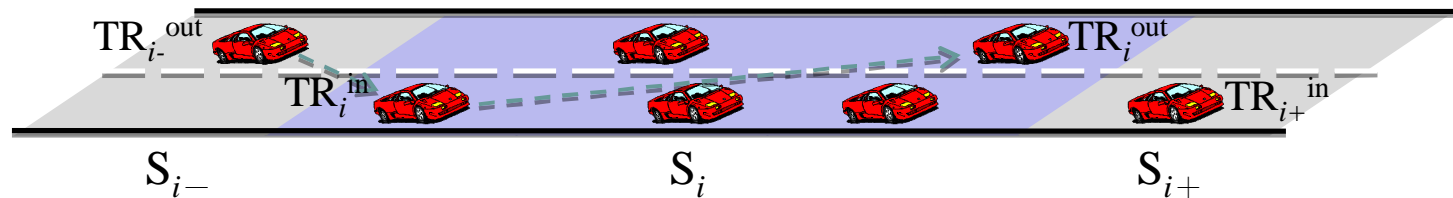
the vehicle closest to S_{i-} as $\text{TR}_{i,\text{new}}^{\text{in}}$

$\text{TR}_{i,\text{new}}^{\text{in}}$	v_1
$\text{TR}_{i,\text{new}}^{\text{out}}$	v_1, v_2

CVIA



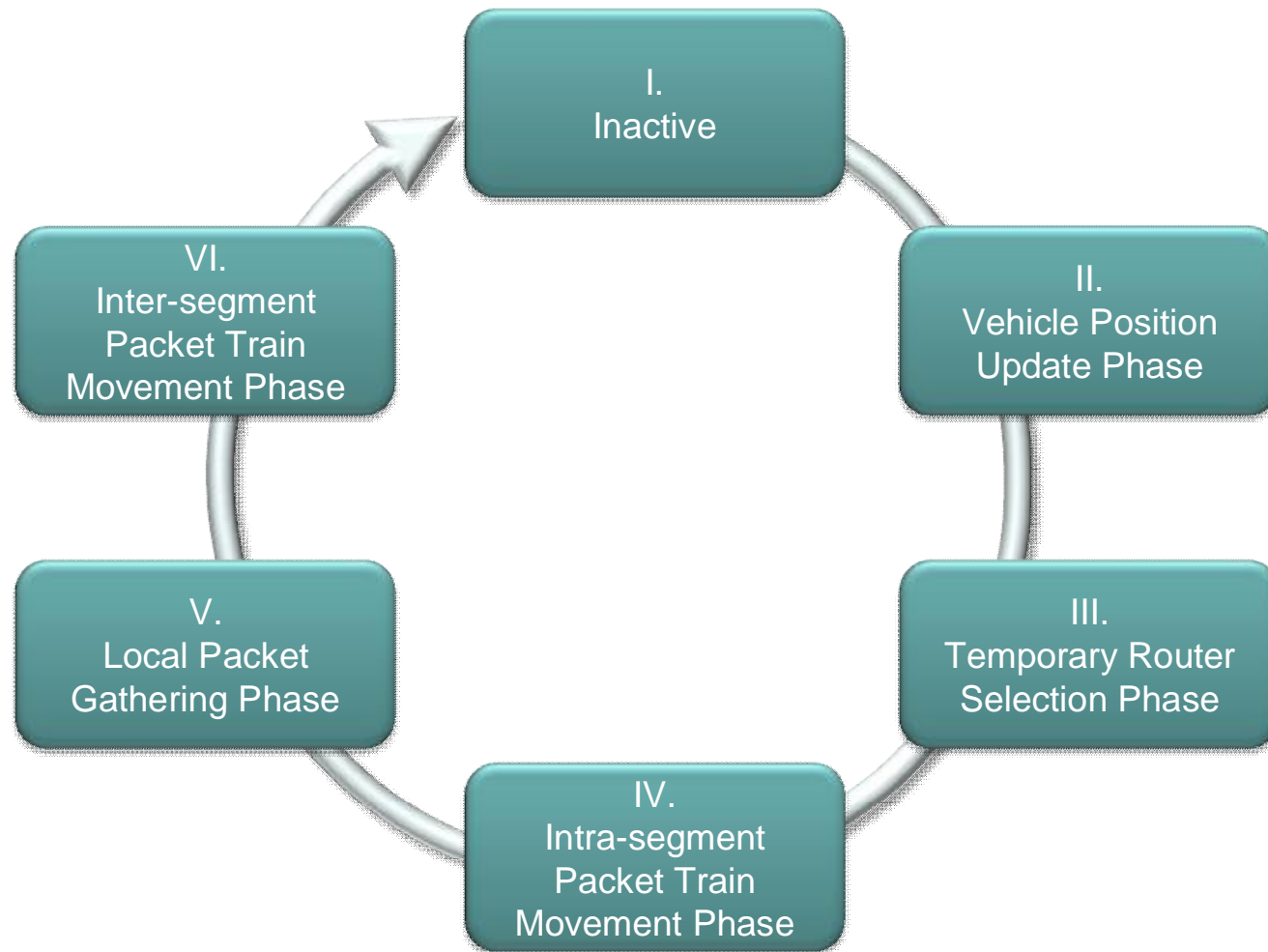
CVIA - Intra-segment Packet Train Movement Phase



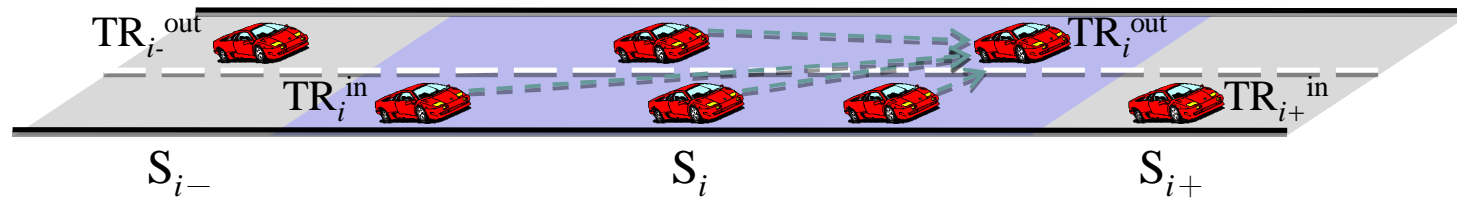
TR_i^{in} delivering the packet train coming from S_{i-} to TR_i^{out} .

To avoid contention, TR_i^{in} has the highest access priority and waits only SIFS duration before accessing the channel.

CVIA



CVIA - Local Packet Gathering Phase

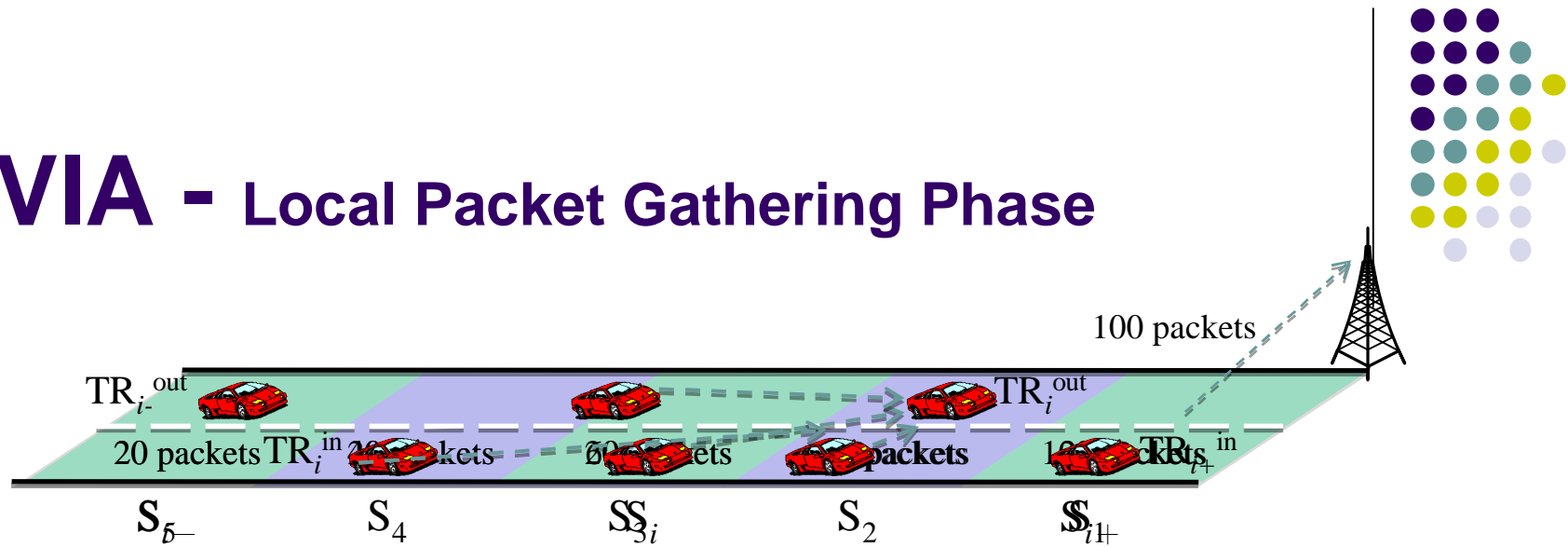


Vehicles directly send their packets to TR_i^{out} .

Each vehicle starts this phase with a random backoff counter.

TR_i^{out} has the highest channel access priority in this phase.

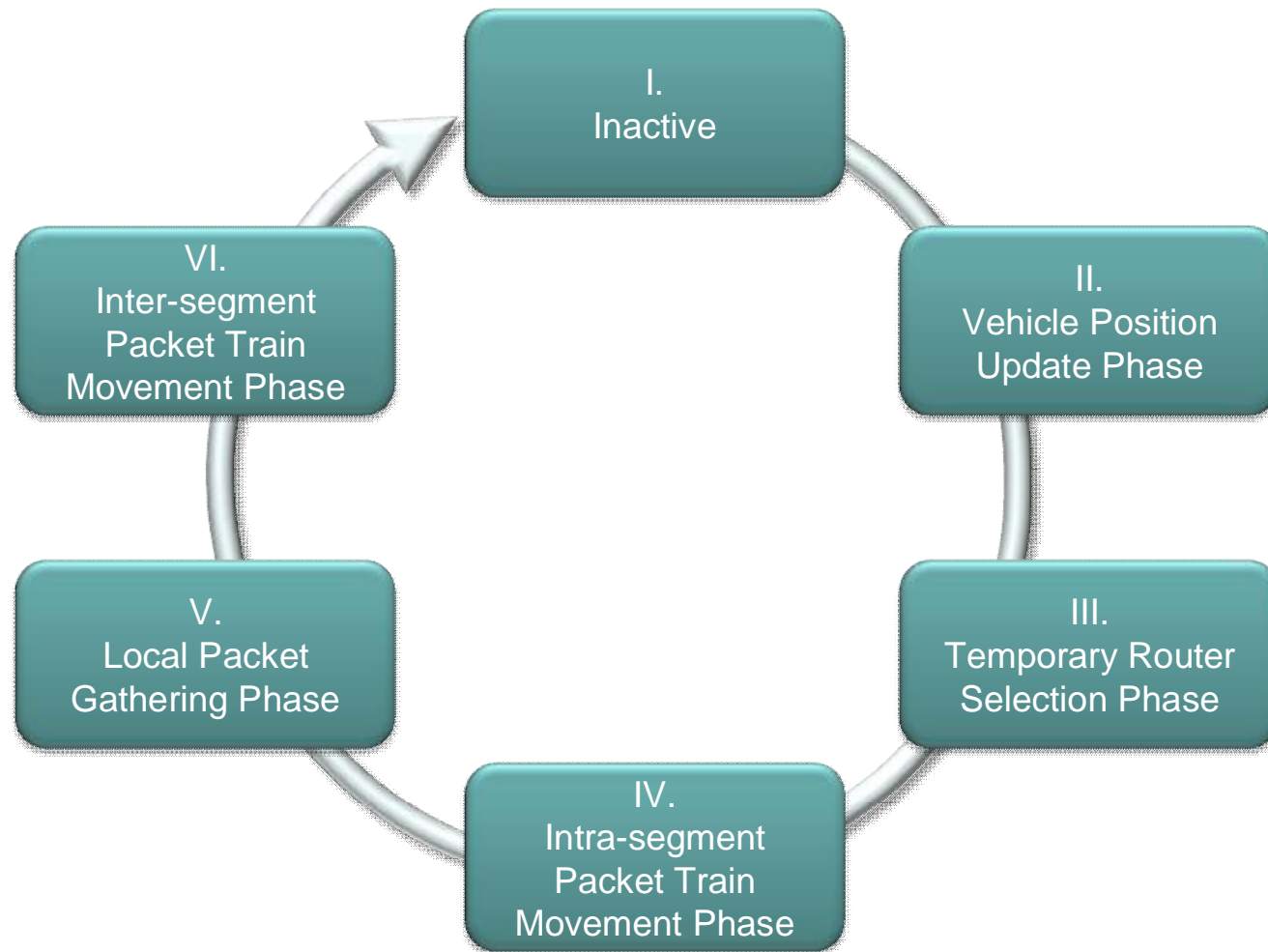
CVIA - Local Packet Gathering Phase



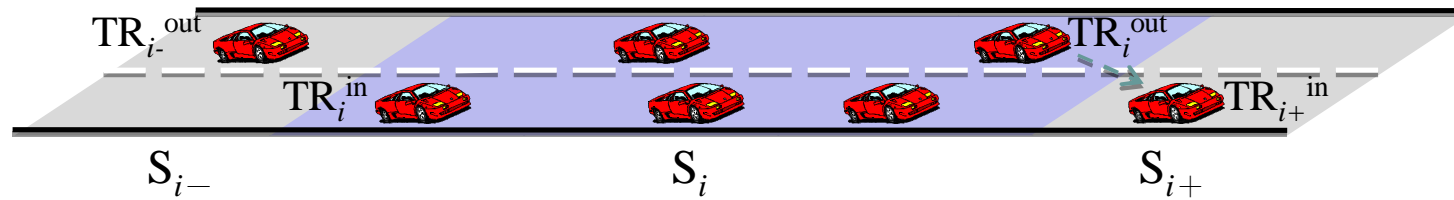
TR_i^{out} accesses the channel and ends the LPG phase in 2 cases :

1. When the total number of packets queued up in TR_i^{out} reaches $\left(\frac{C}{N}\right)(N-i+1)$
2. When the time left in the active slot is just enough to move out all packets in the queues.

CVIA



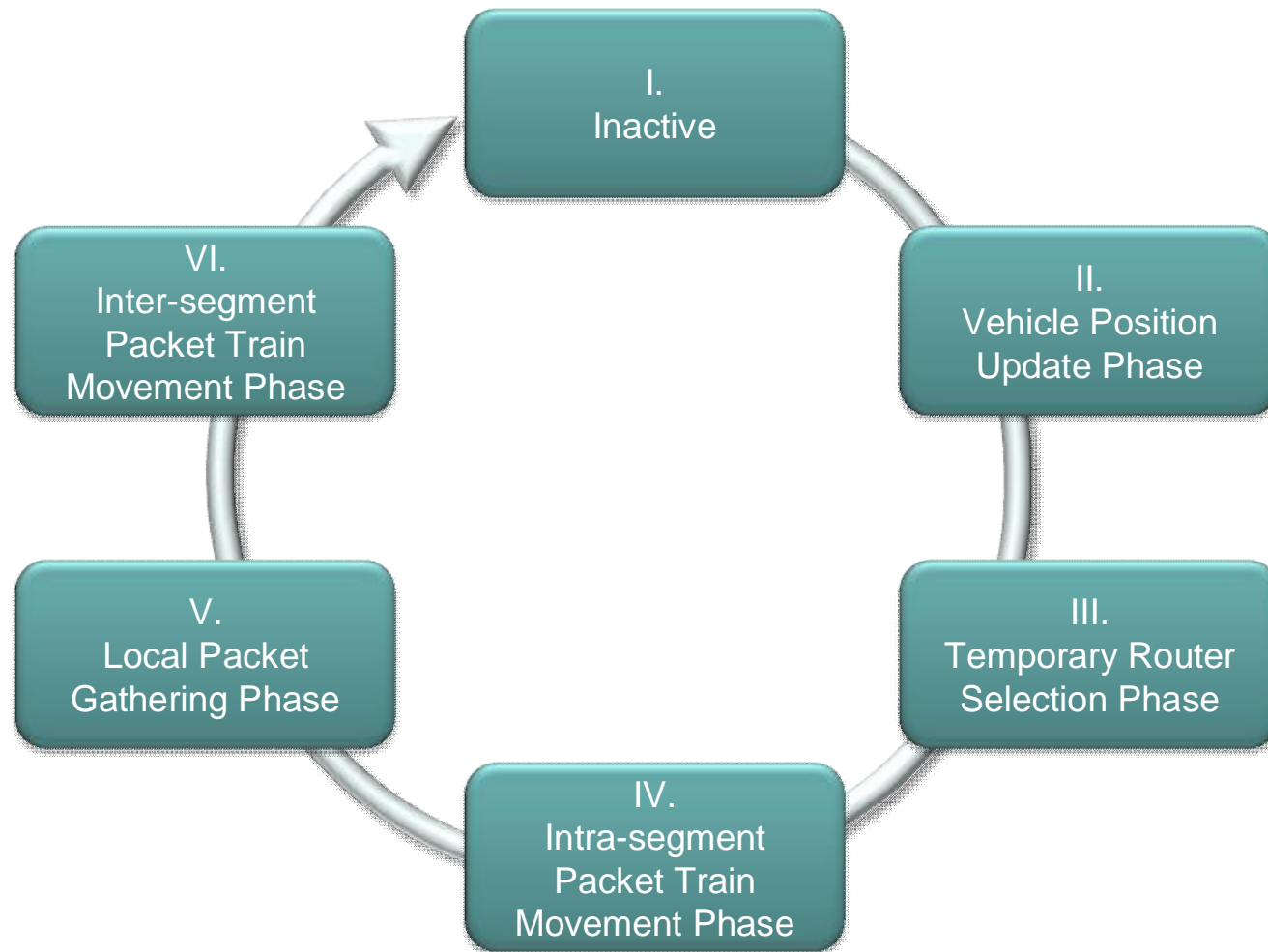
CVIA - Inter-segment Packet Train Movement Phase



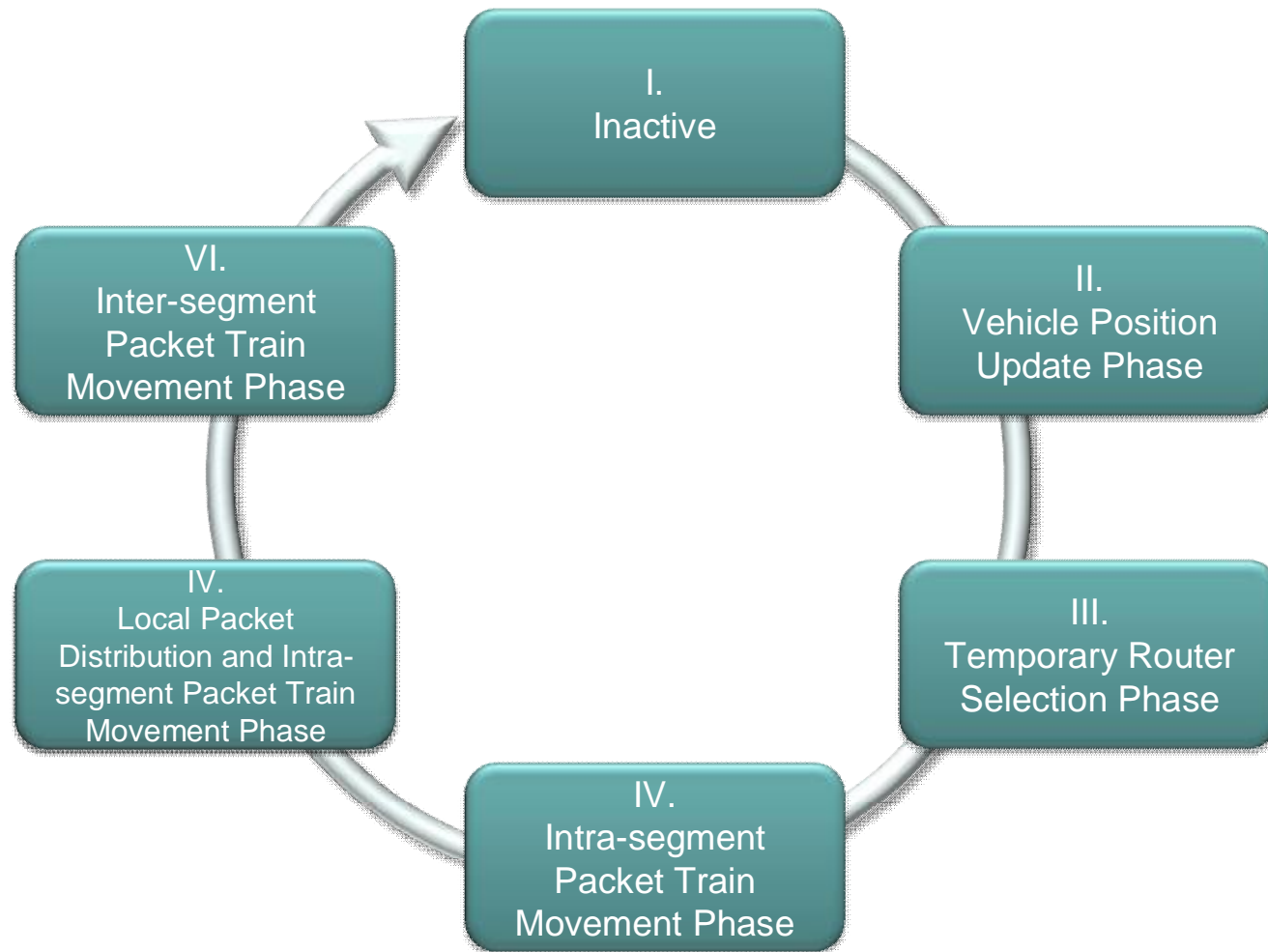
TR_i^{out} forms a packet train and moved to TR_{i+}^{in} before the end of TS_j .

There is only one TR_{i+}^{in} in S_{i+} , TR_i^{out} does not need to specify the vehicle ID of the destination in the data train.

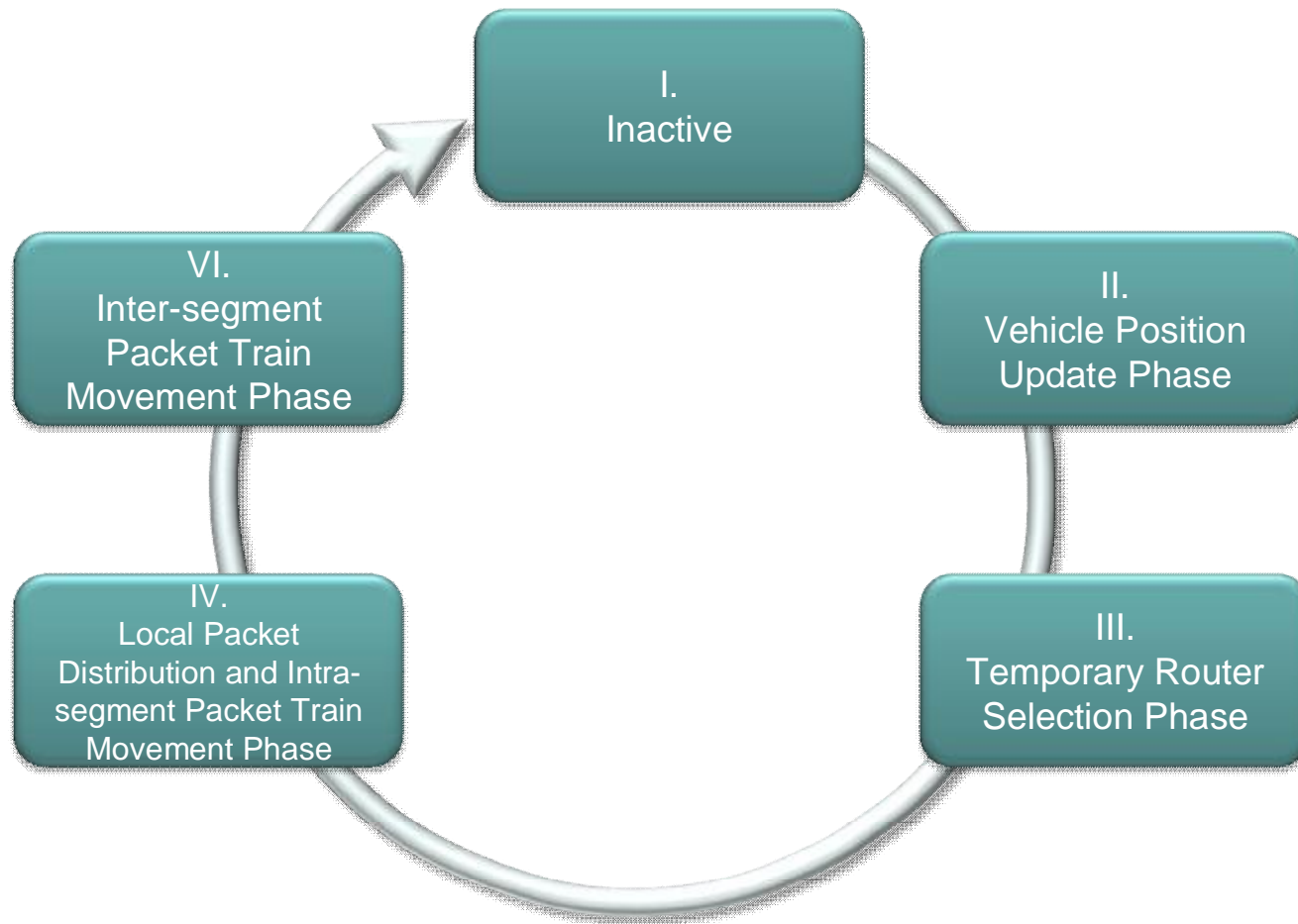
CVIA



CVIA-State diagram-Download

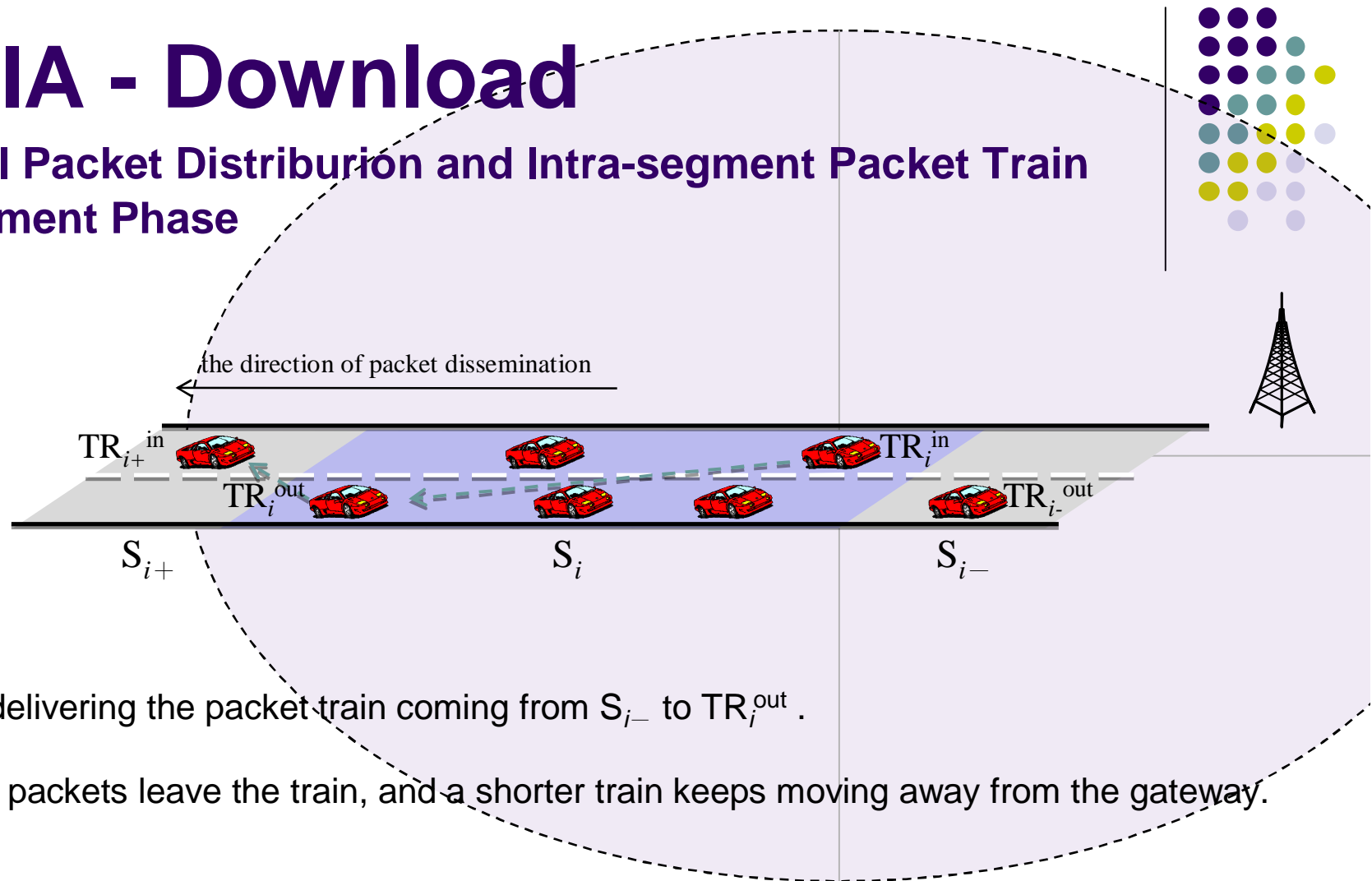


CVIA-State diagram-Download



CVIA - Download

Local Packet Distribution and Intra-segment Packet Train Movement Phase



TR_i^{in} delivering the packet train coming from S_{i-} to TR_i^{out} .

Some packets leave the train, and a shorter train keeps moving away from the gateway.



Performance Evaluation

Simulation scenarios	
Density	34 vehicles/km
Speed	90 \pm 5 km/h (do not change during simulation)
Transmission range	350 m
Data rate	27 Mbps
Payload	2312 or 500 bytes
Base protocol	802.11a
Interference range to transmission range ratio	1
Maximum number of packet retrials	10
Simulation time	10 s
Simulation repetitions	20

Performance Evaluation



CVIA parameters	
T_{slot}	100 ms
N	4 or 8
r	2

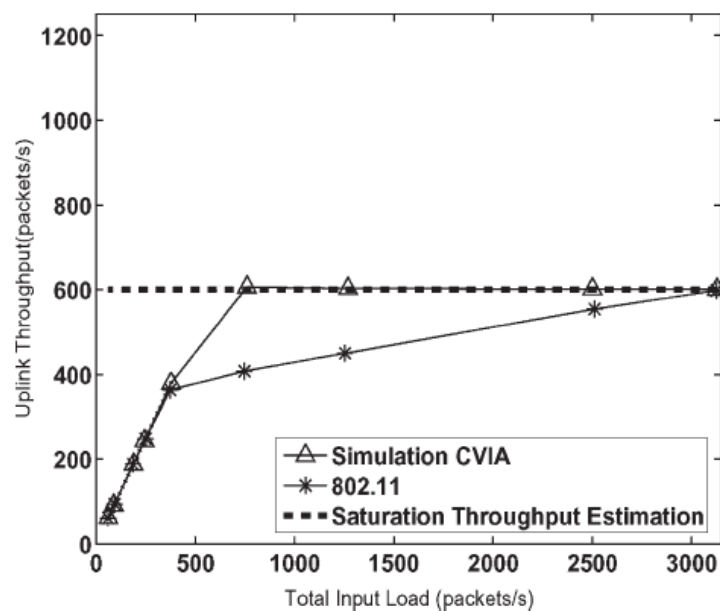
Scenario	payload	VTR
I	2304 bytes	$4R$ ($N=4$)
II	500 bytes	$4R$ ($N=4$)
III	2304 bytes	$8R$ ($N=8$)



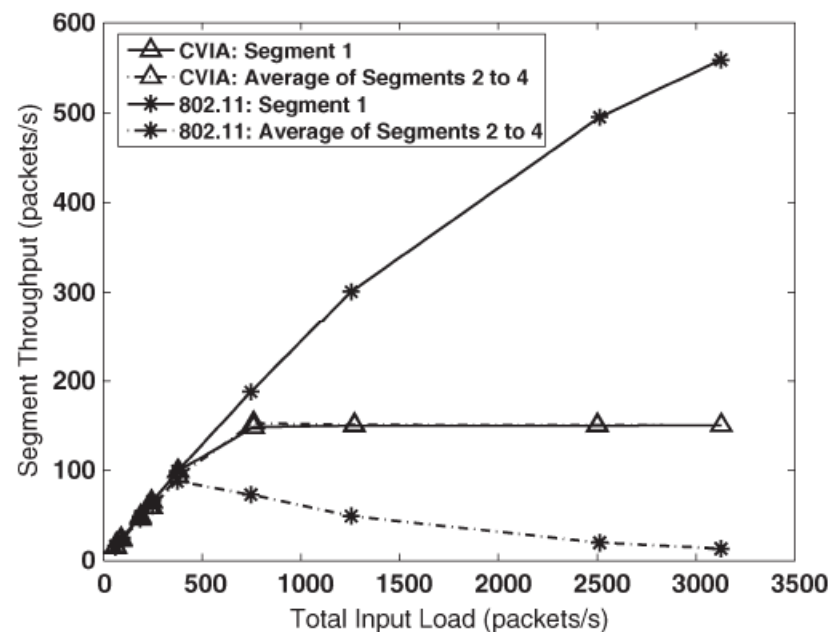
Performance Evaluation

Scenario I

Throughput



Detailed segment throughputs

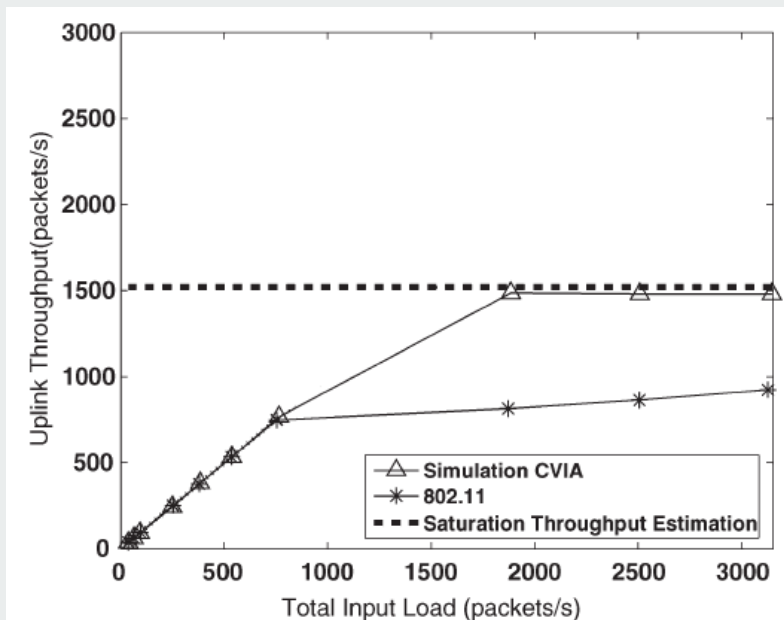




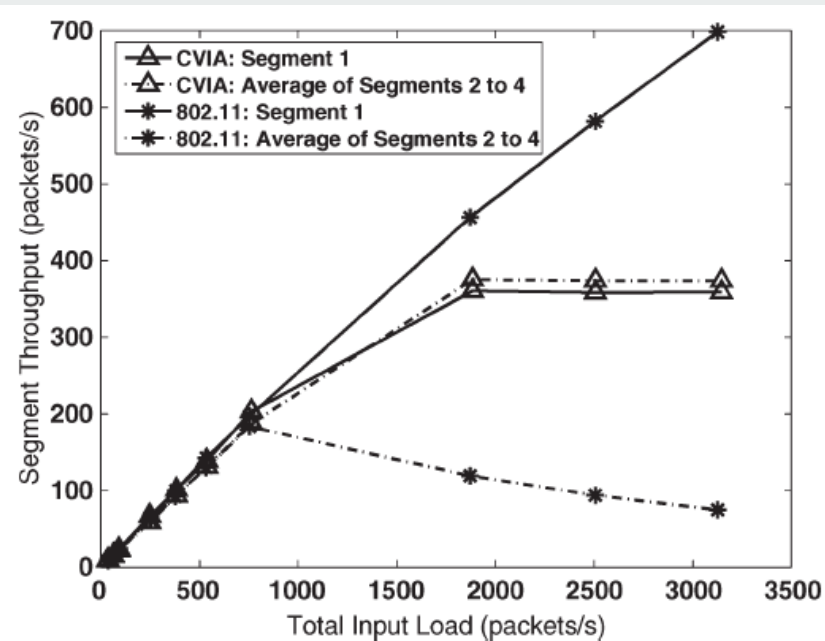
Performance Evaluation

Scenario II

Throughput



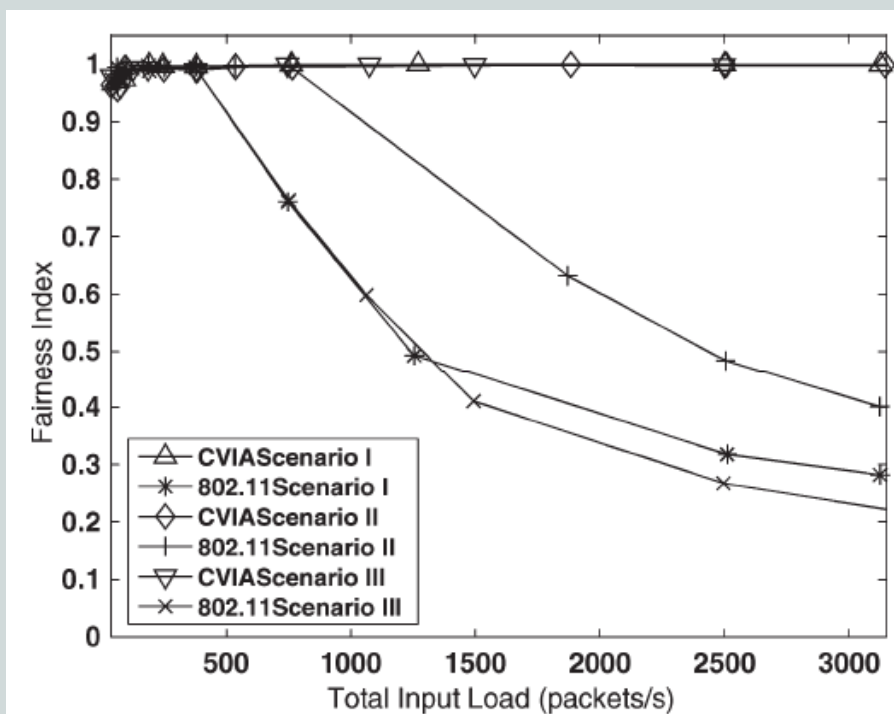
Detailed segment throughputs





Performance Evaluation

Fairness



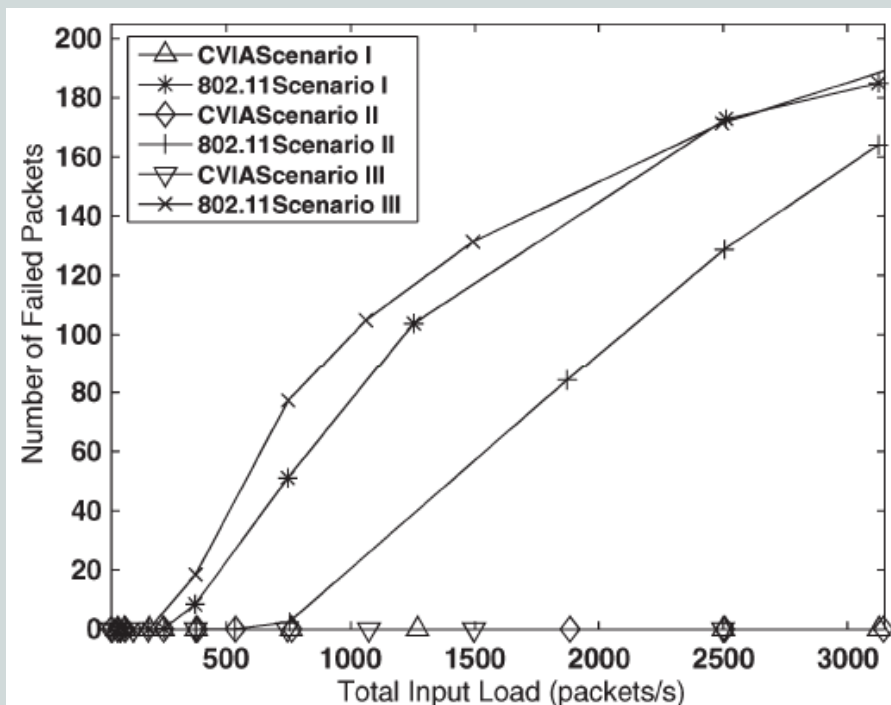
$$\text{Fairness Index (FI)} = \frac{(\sum_{i=1}^N \text{Thr}_i)^2}{N \cdot \sum_{i=1}^N \text{Thr}_i^2}$$

Thr_i : the throughput of segment i



Performance Evaluation

Failed packets



In LPG phase, the contention-based channel access is used.

Probability of packet collision ≈ 0.3

Probability of packet dropping $\approx 0.3^{10}$
 $\approx 5 \times 10^{-6}$



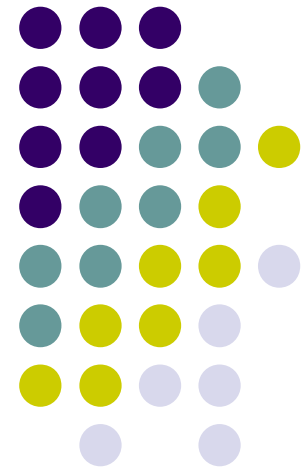
Conclusion

- | Mitigates the hidden node problem.
- | Avoids contention while moving packets among road segments.
- | Provides fairness among segments by controlling the contents of the packet trains.
- | Unnecessary RTS/CTS overhead while moving packets among road segments.
- | The throughput gain obtained using the proposed CVIA protocol is higher for short payload scenario.

Self-Configuring TDMA Protocols for Enhancing Vehicle Safety With DSRC Based Vehicle-to-Vehicle Communications

Fan Yu and Subir Biswas

*Department of electrical and computer engineering,
Michigan State University*



*IEEE Journal on Selected Areas in
Communications, JSAC 2007*

Outline

- | Introduction
- | *VeSOMAC* MAC protocol
- | Simulation
- | Conclusion





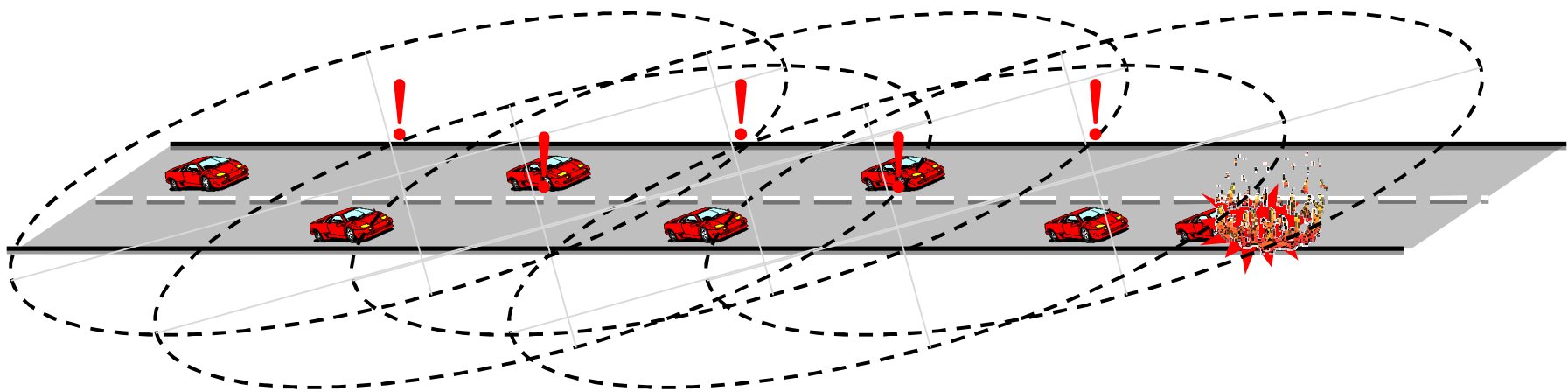
Introduction

- | The ITS architecture is currently being developed for enhancing vehicle safety using V2V and V2R communications.
- | DSRC is a communication standard for ITS.
- | IEEE 802.11a is recommended as the MAC for DSRC.
- | It suffers from unbounded delivery latency at higher loads because of the underlying random access.
- | This implies that 802.11a may not be able to provide the required message delivery latency for a delay-sensitive ITS safety application



Goal

- ┆ Propose a MAC protocol for distributed TDMA allocation in V2V wireless networks.
 - ┆ Minimizing the multi-hop delivery delay.
 - ┆ Fast reconfiguration during frequent topology changes.
 - ┆ Can work without network time synchronization.

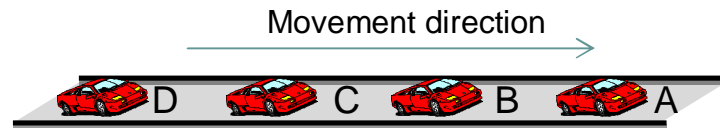


Network Assumption

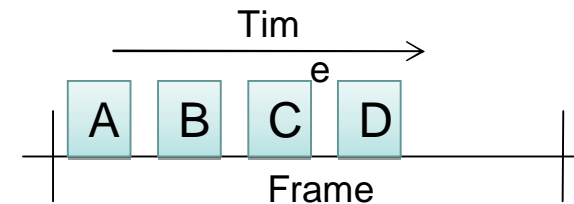


- Each vehicle is equipped with GPS devices used for obtaining vehicle positions and speed.

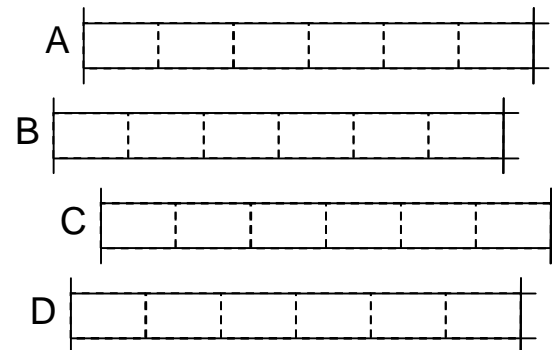
- Uni-direction



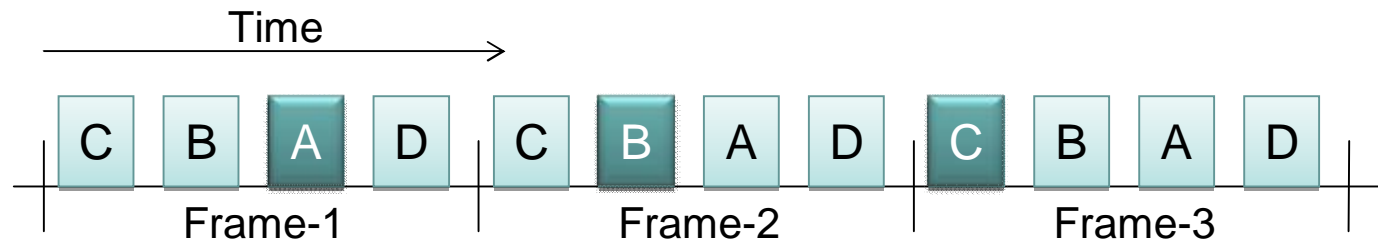
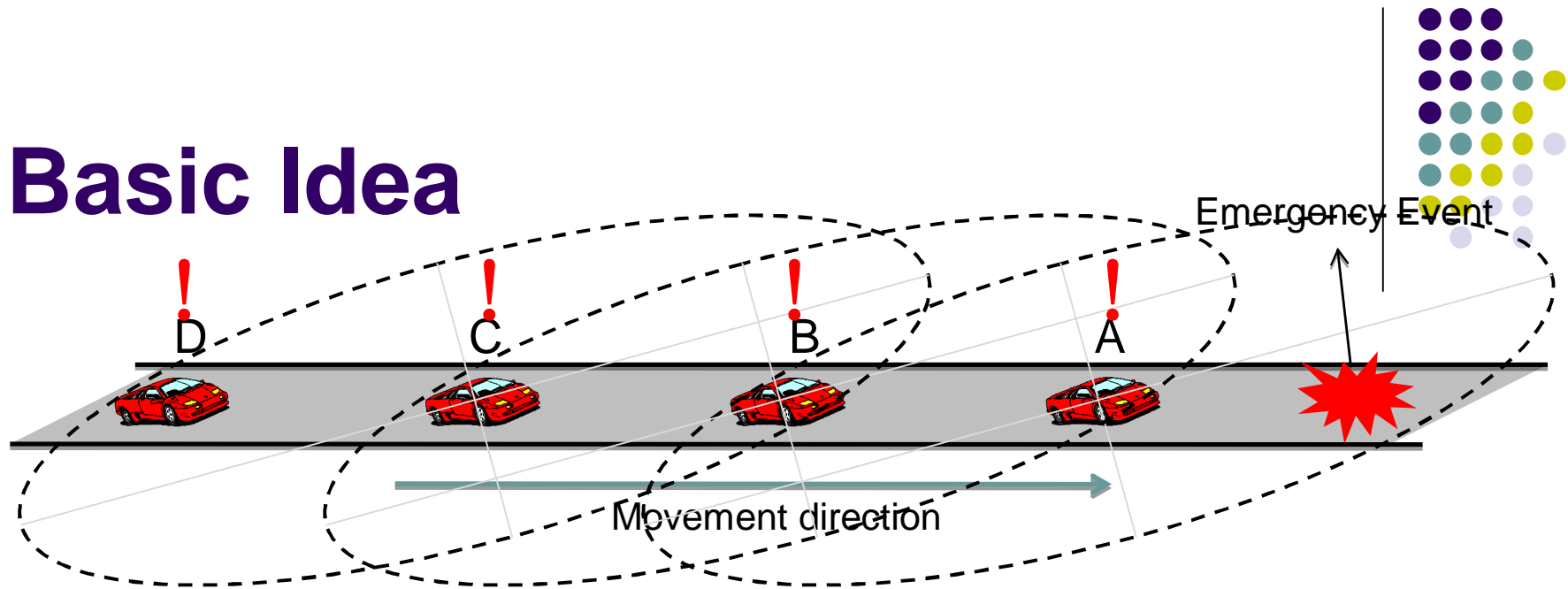
- Each vehicle allocated one slot in one frame.



- Asynchronous

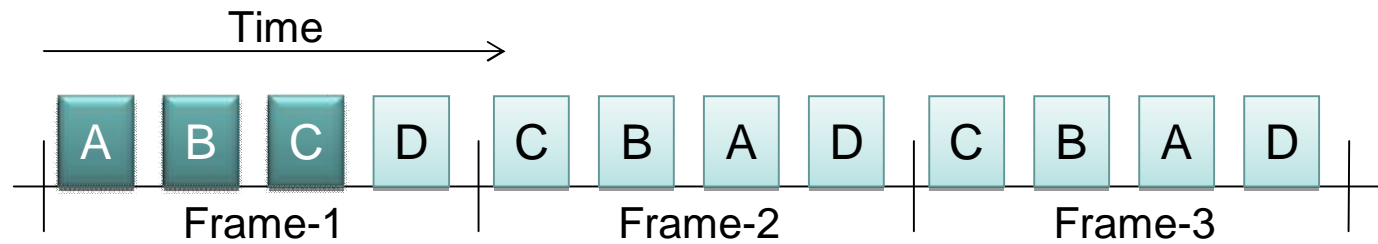
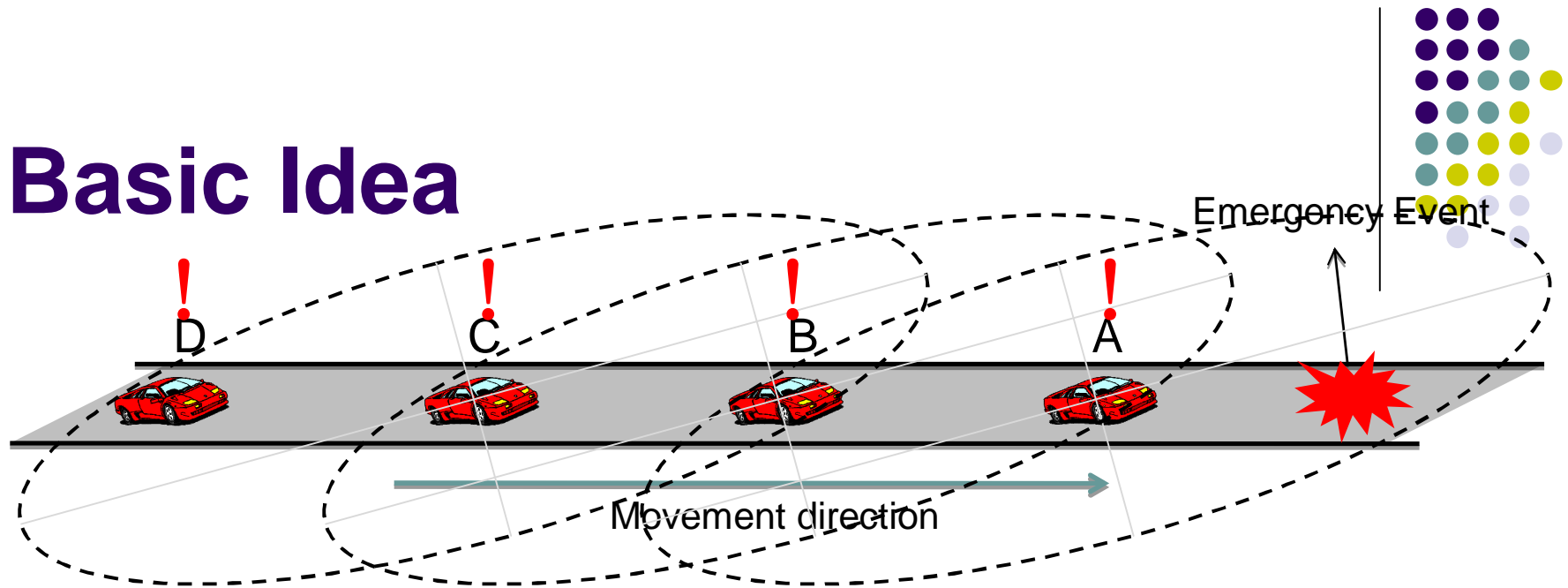


Basic Idea



Need 3 Frames

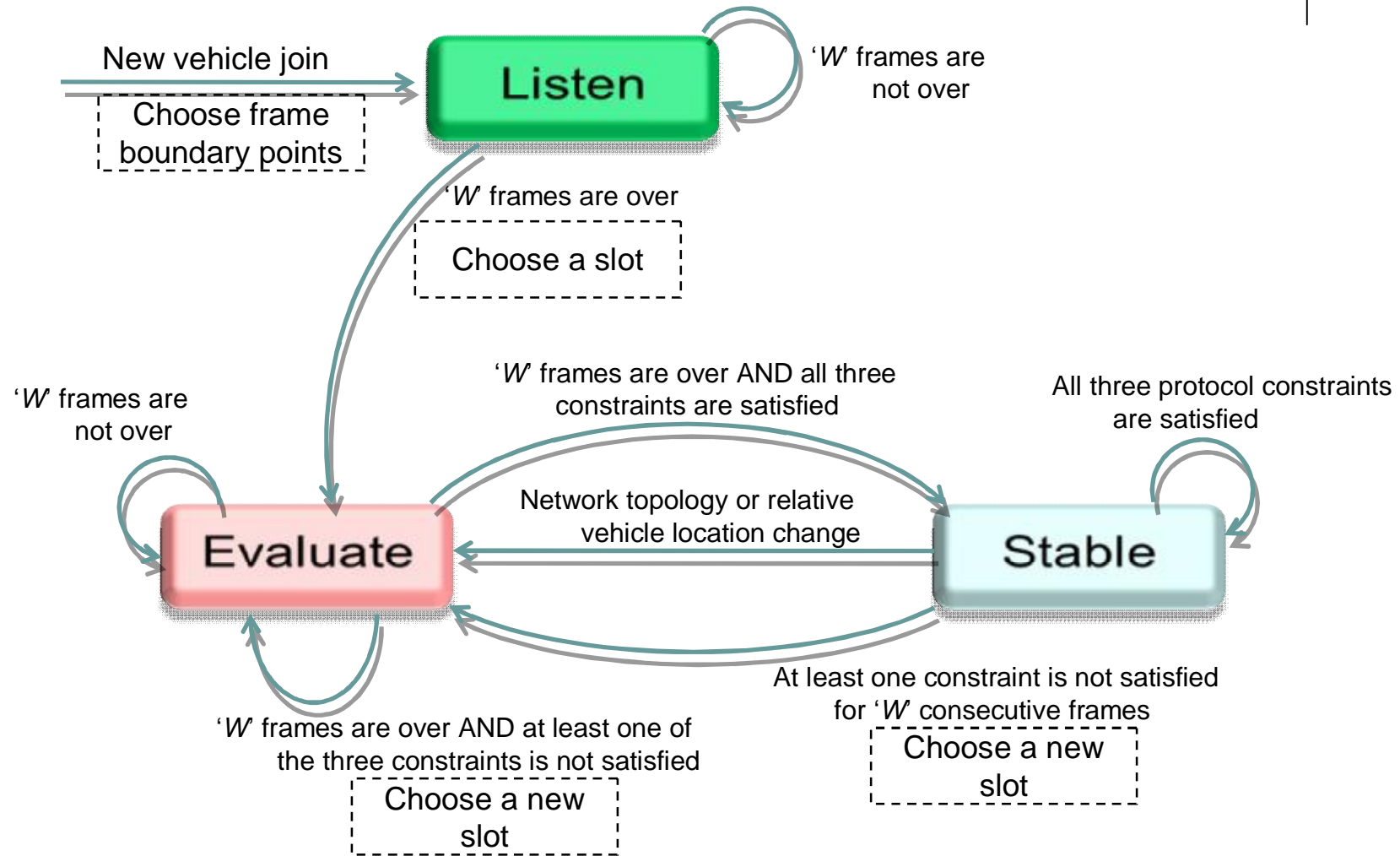
Basic Idea



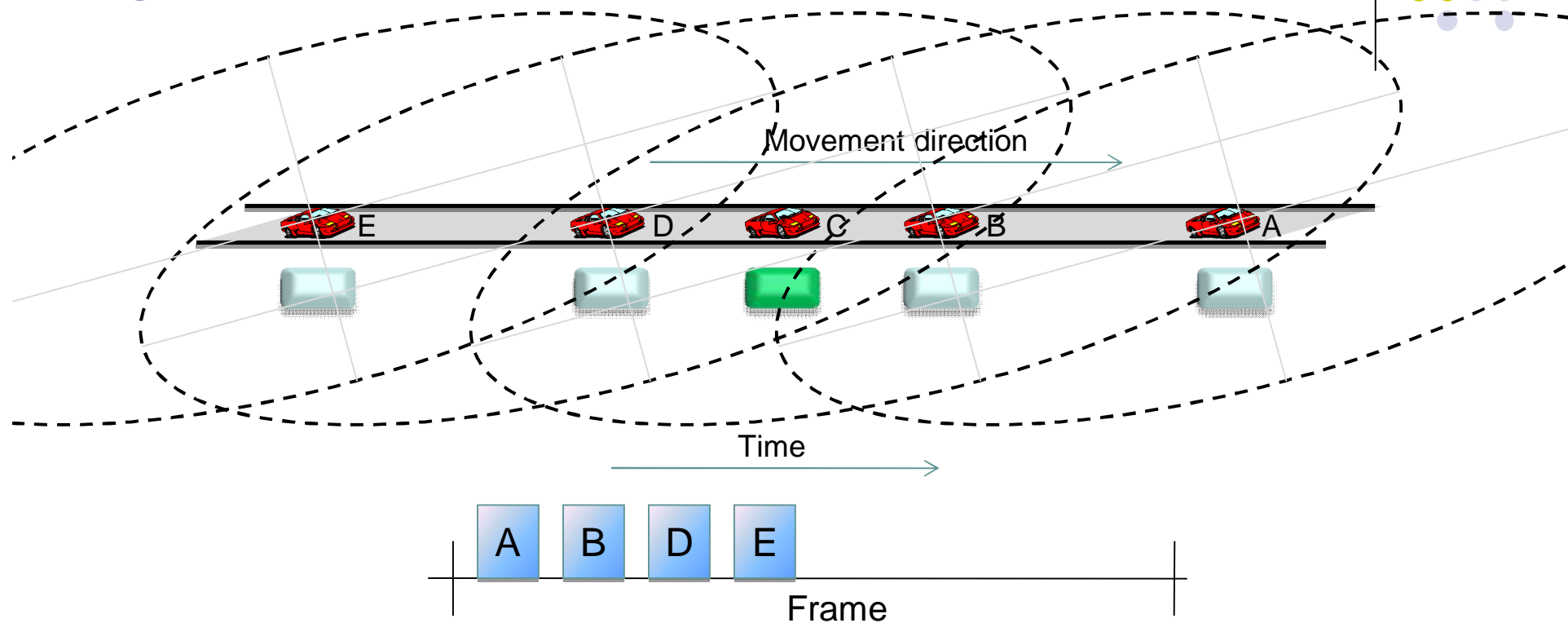
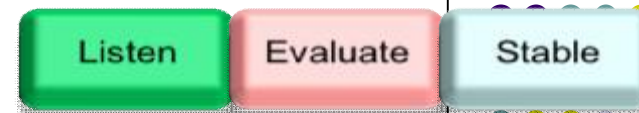
Need 1 Frame



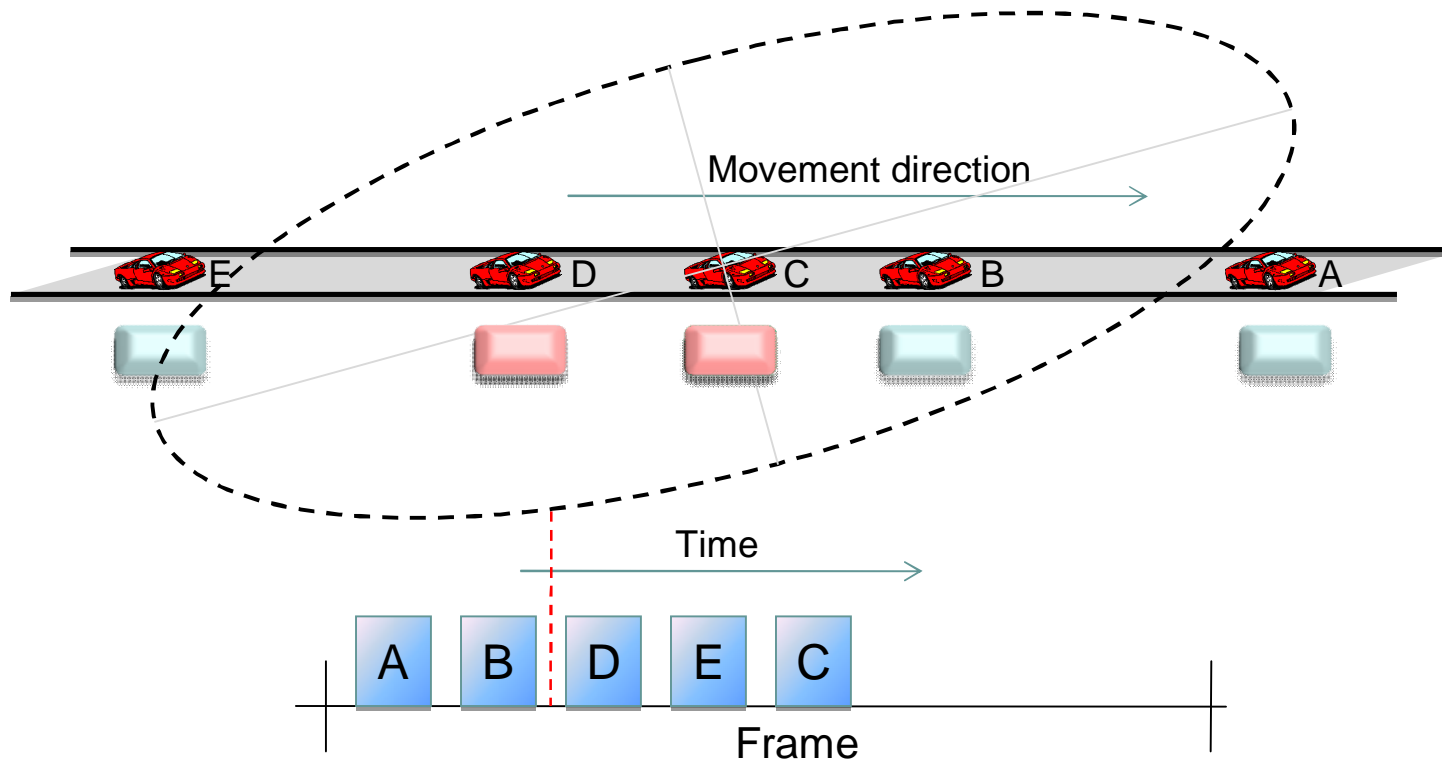
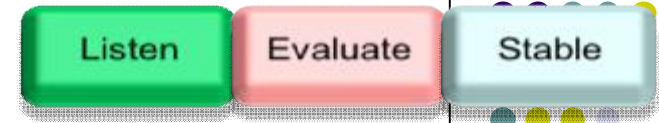
Protocol State Machine



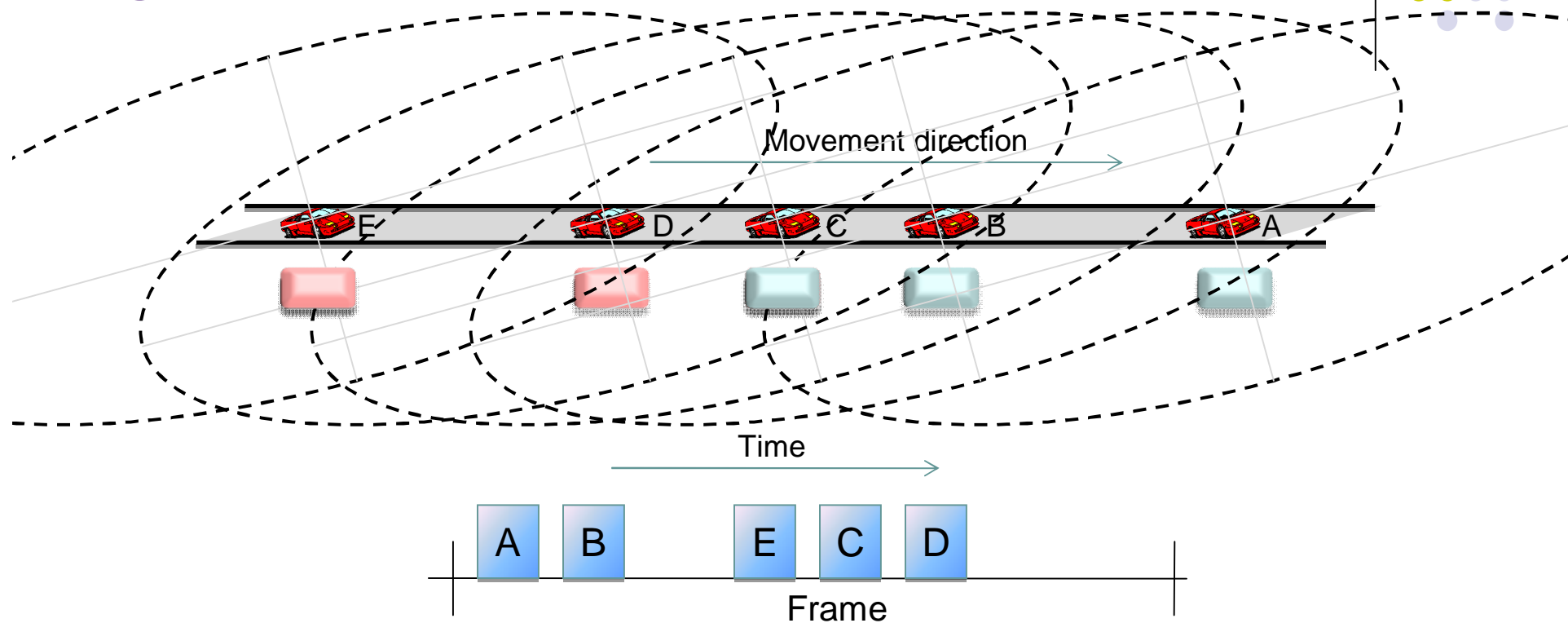
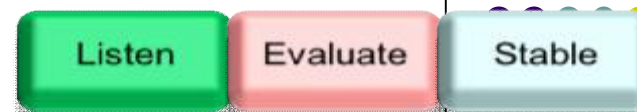
Overview



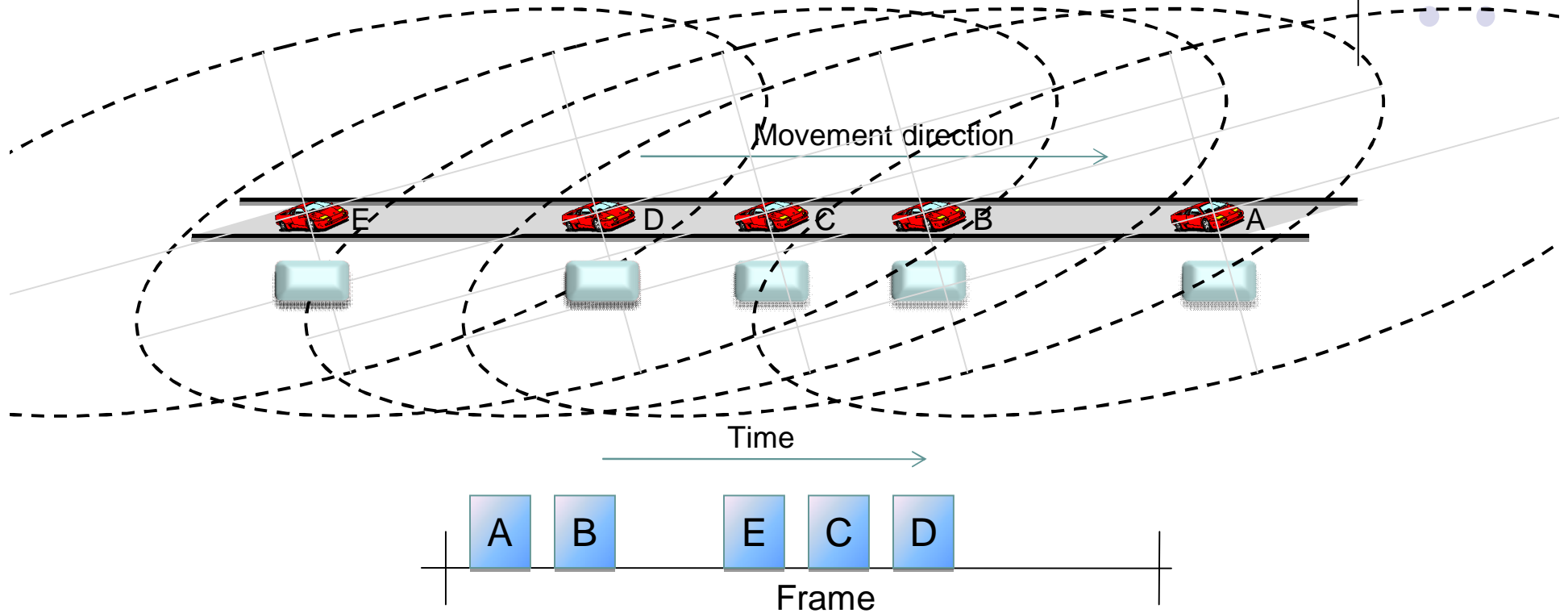
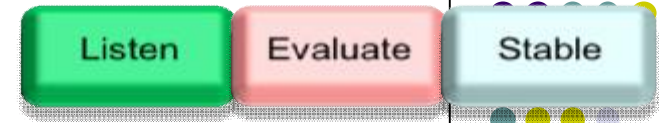
Overview



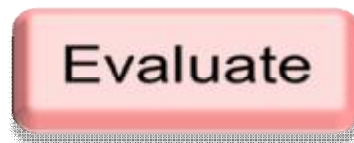
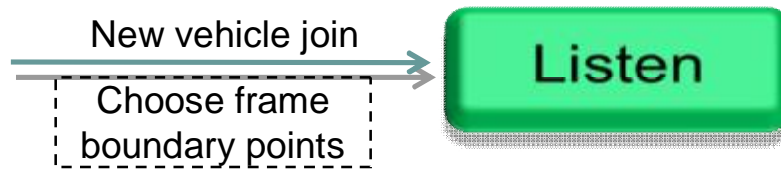
Overview



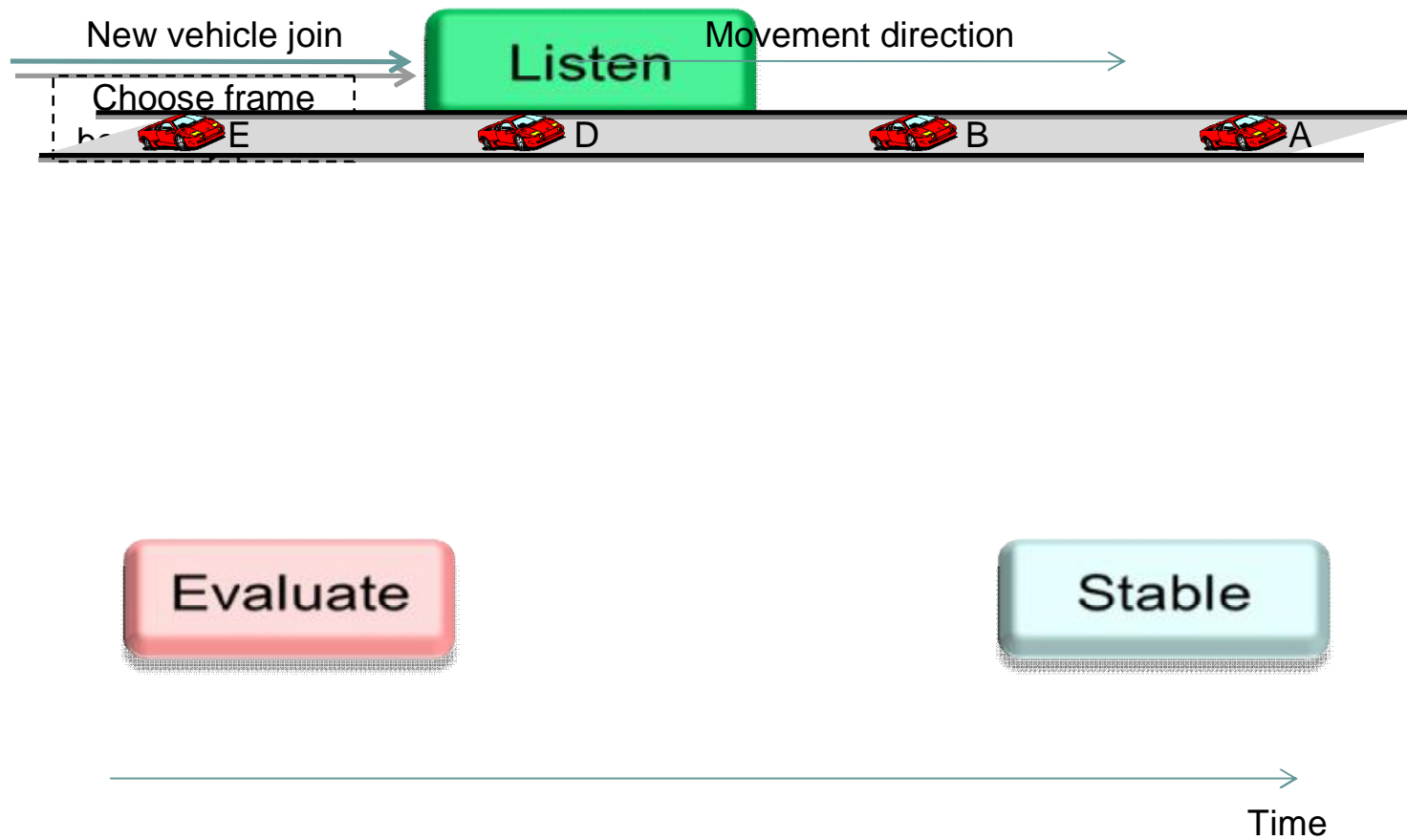
Overview

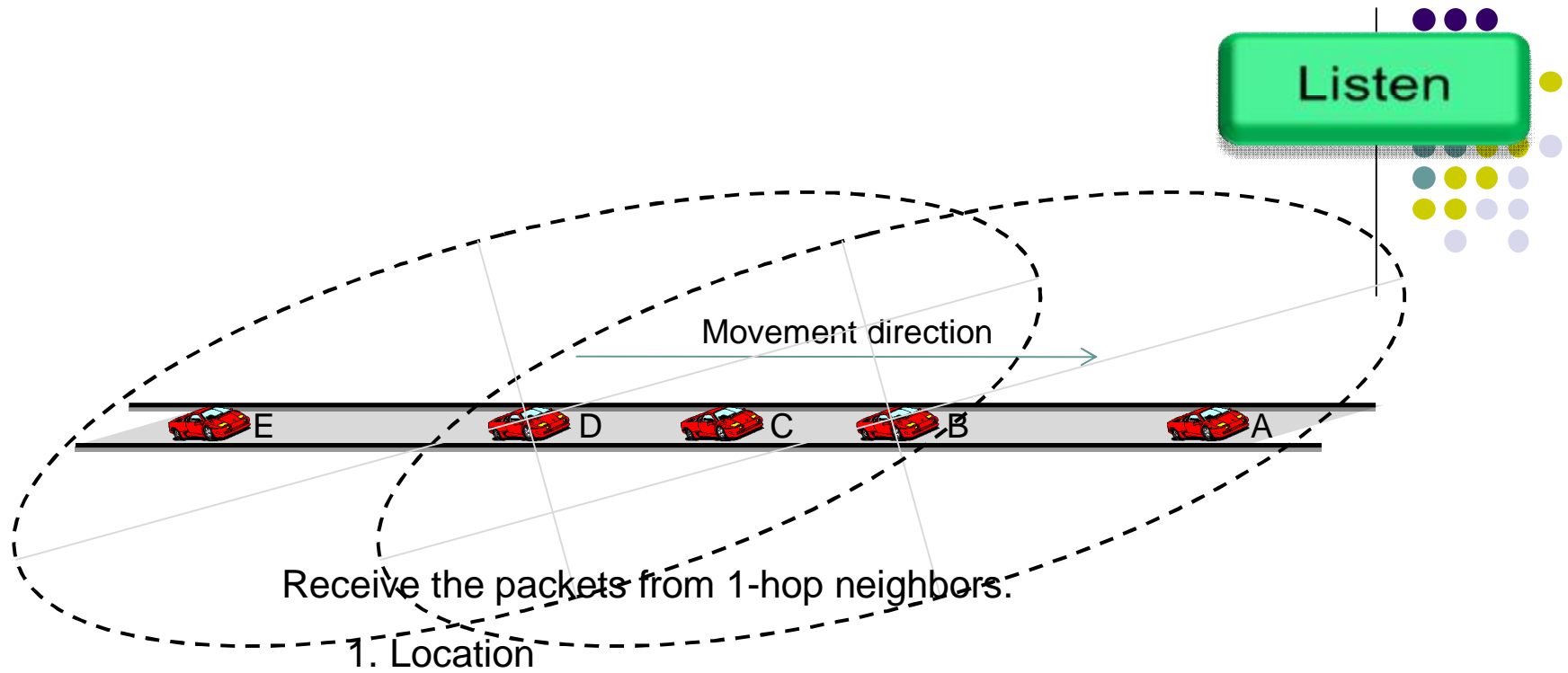


Protocol State Machine



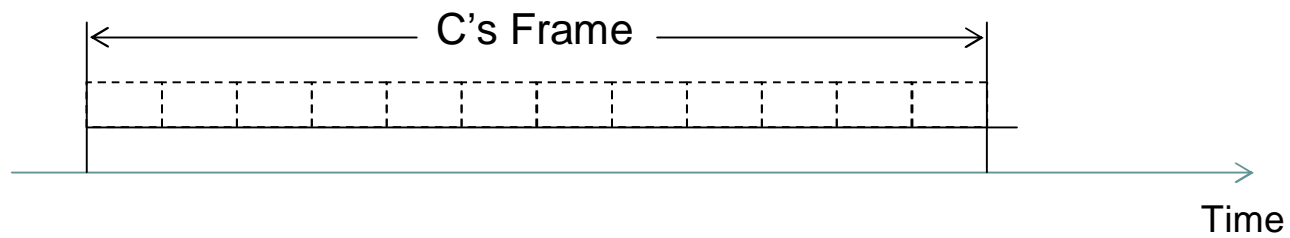
Protocol State Machine





2. Allocated slot

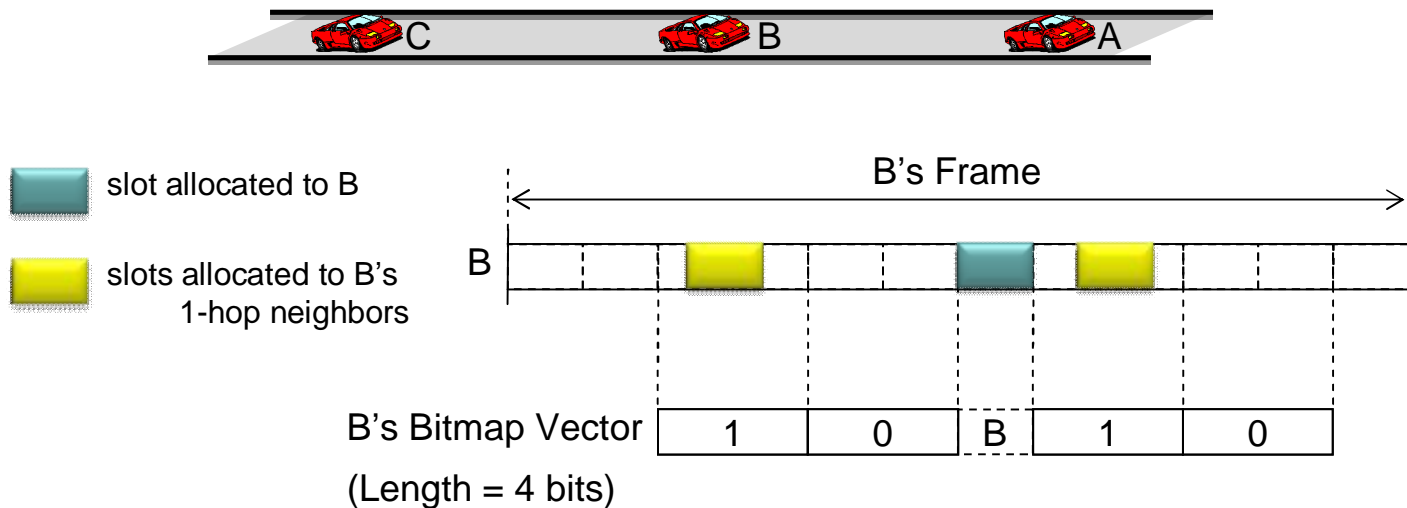
3. Bitmap



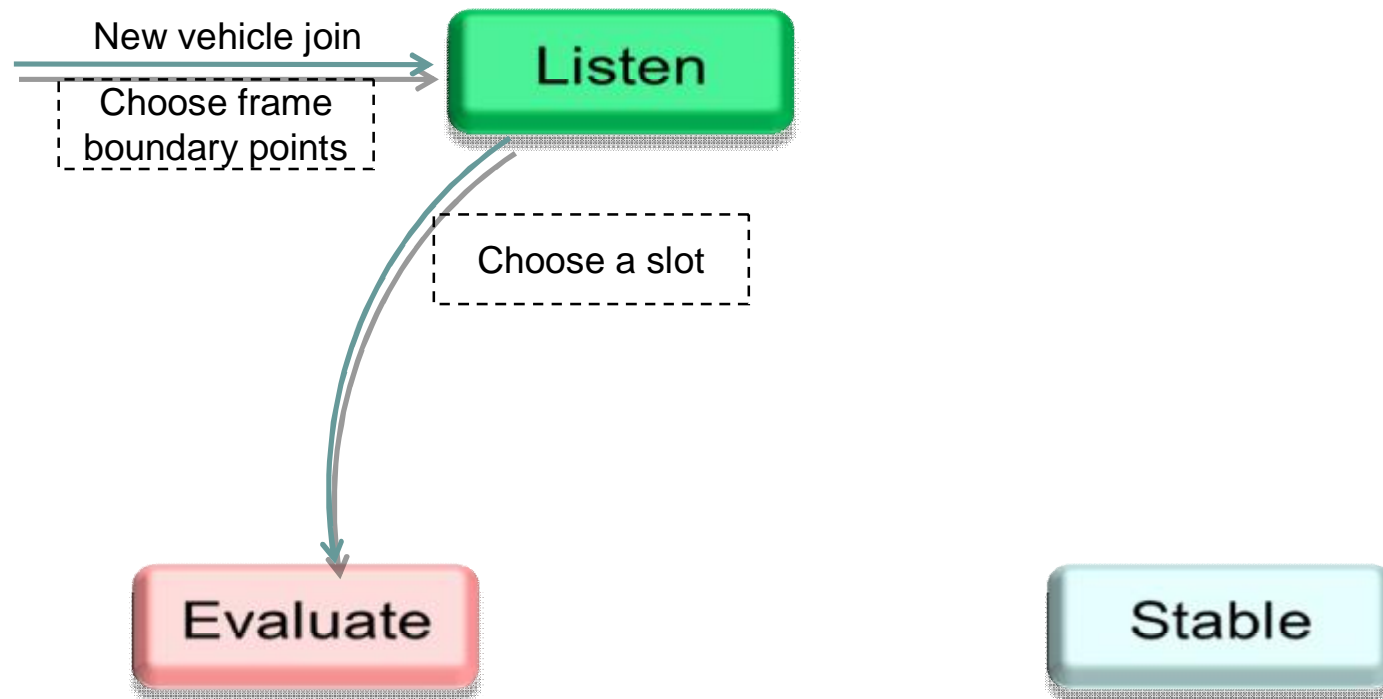
Bitmap



- I Informs vehicle's 1-hop neighbors about the slots occupied by its 1-hop neighbors.
- I By listening to the bitmaps in all received packets, a vehicle can detect the slot locations of its 1-hop and 2-hop neighbors.



Protocol State Machine



Constraint



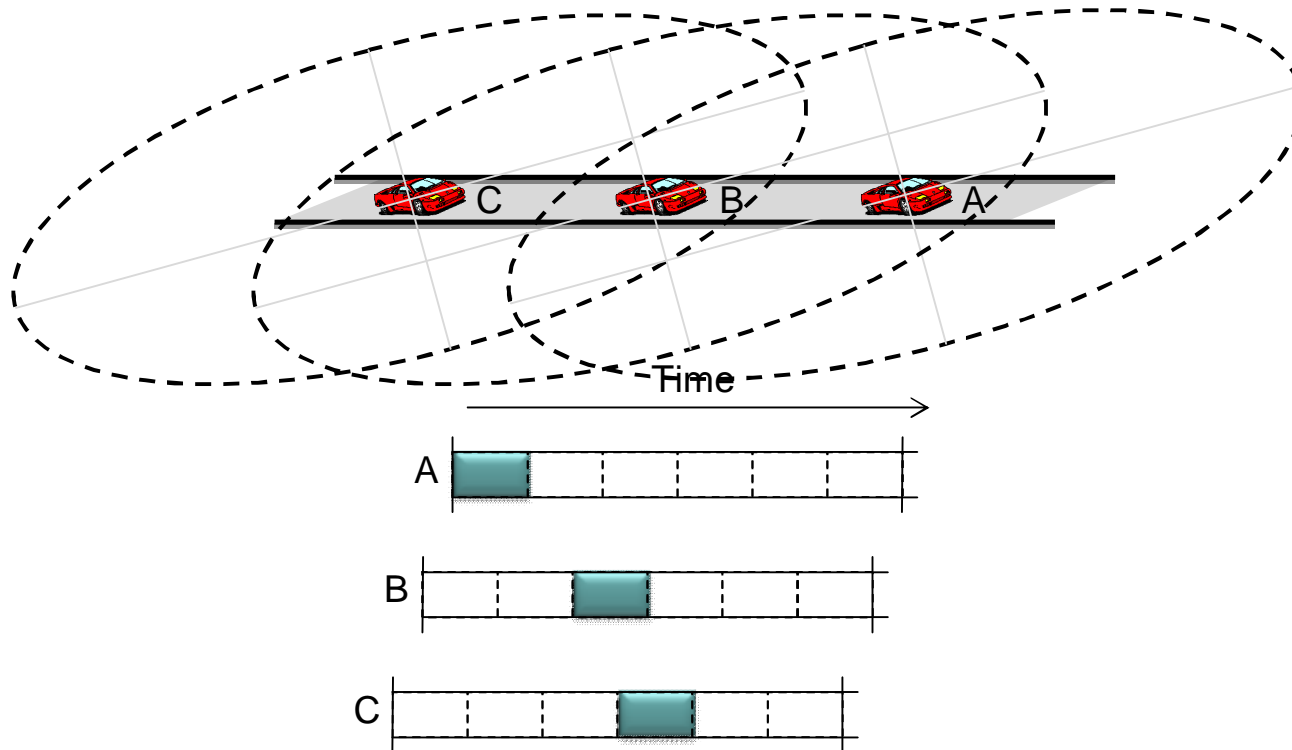
- | Slot allocation in VeSOMAC needs to satisfy the three constraint
 - | Timing Constraint
 - | Bitmap Constraint
 - | Ordering Constraint

Constraint



- | Slot allocation in VeSOMAC needs to satisfy the three constraint
 - | Timing Constraint
 - | No two one-hop or two-hop neighbors' slots can overlap.
 - | Bitmap Constraint
 - | Ordering Constraint

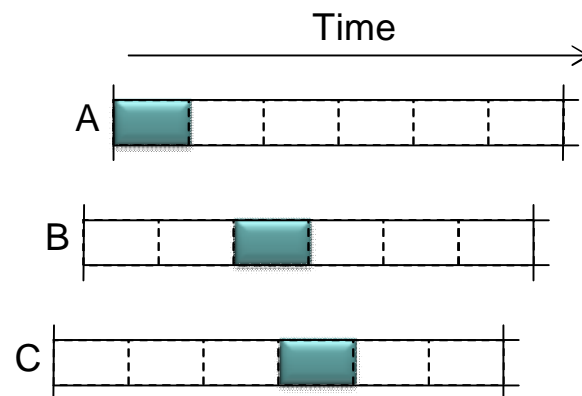
Constraint - Timing Constraint



Constraint - Timing Constraint



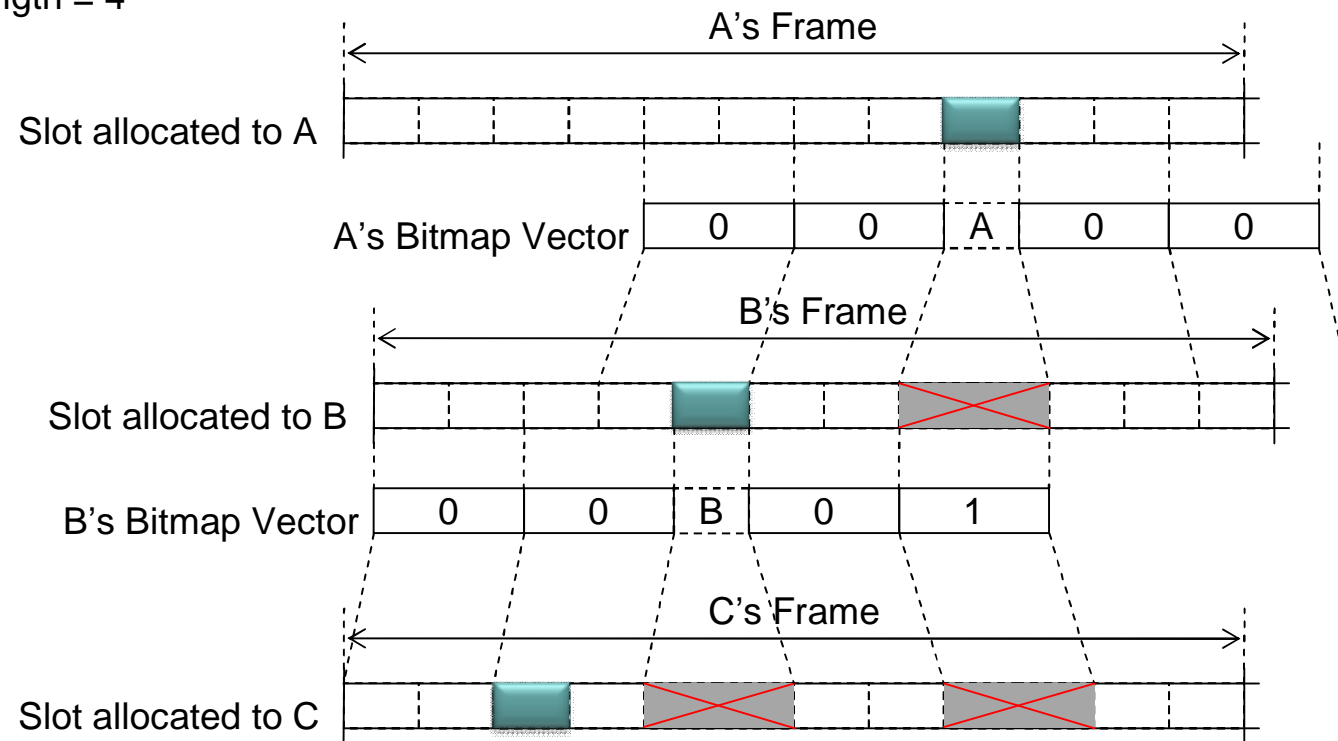
Bitmap Length = 4



Constraint - Timing Constraint



Bitmap Length = 4



Constraint



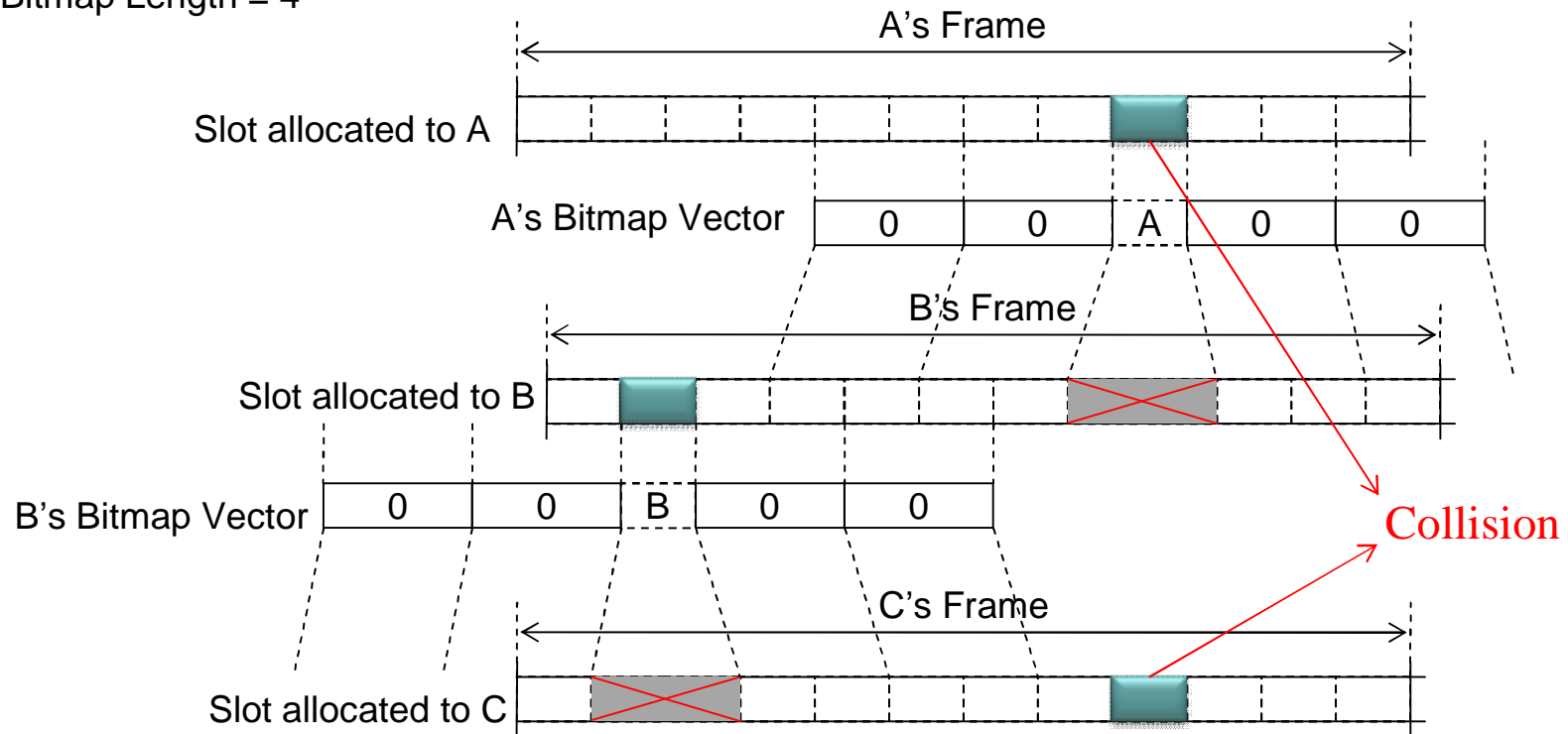
- | Slot allocation in VeSOMAC needs to satisfy the three constraint
 - | Timing Constraint
 - | Bitmap Constraint
 - | For 1-hop neighbors i and j , i 's chosen slot should be able to be represented within the bitmap vector of j . The same is applicable for vehicle j 's slot.
 - | Ordering Constraint

Constraint - Bitmap Constraint



Why Bitmap Constraint?

Bitmap Length = 4

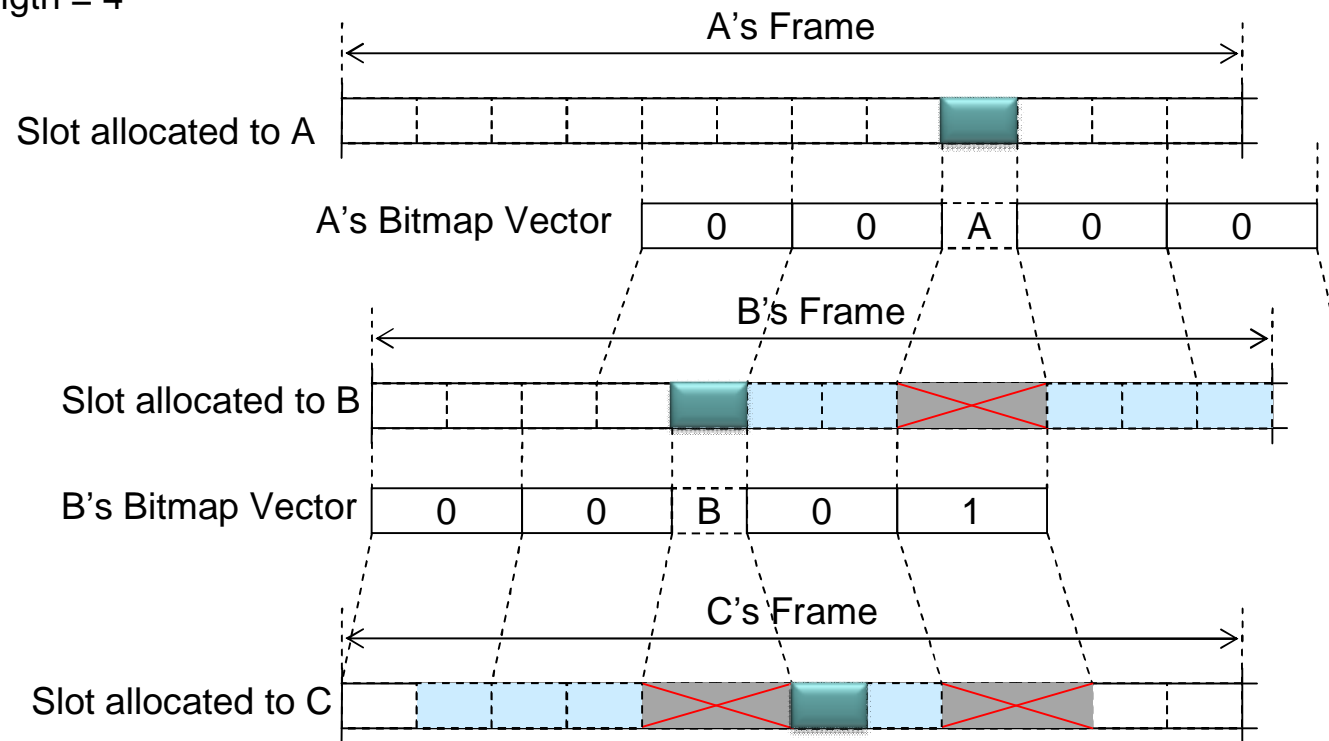


Constraint - Bitmap Constraint



A and B are 1-hop neighbors
A and C are 2-hop neighbors

Bitmap Length = 4

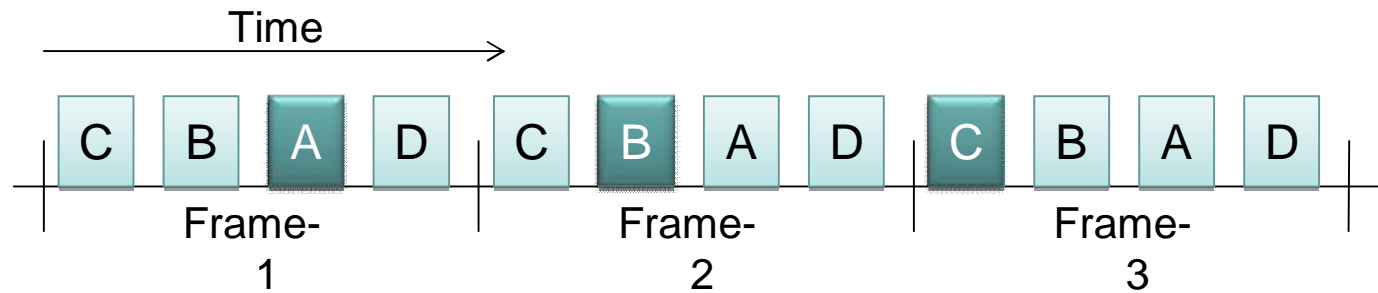
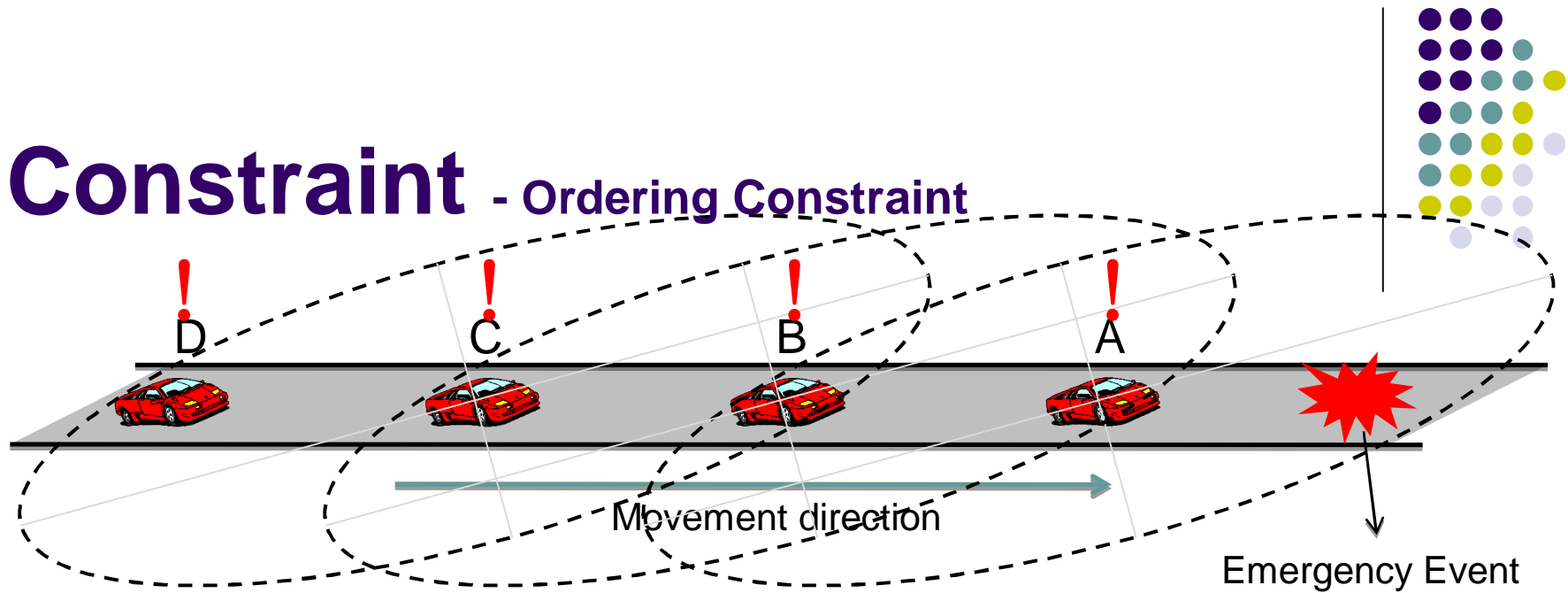


Constraint



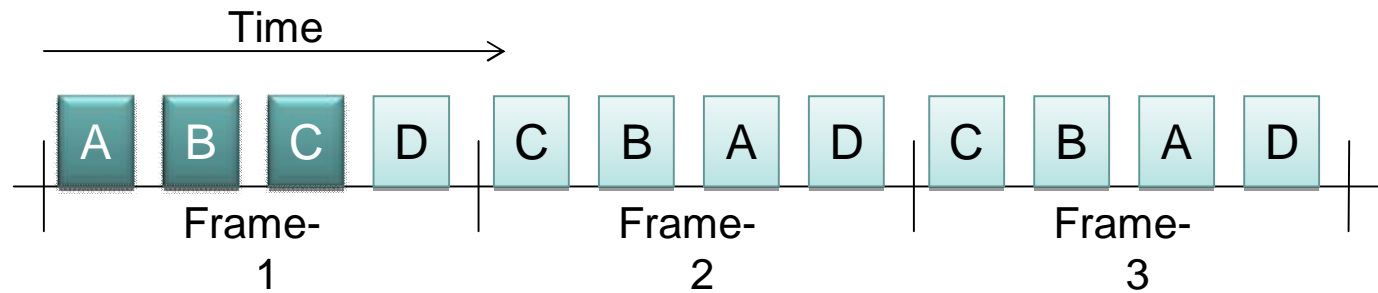
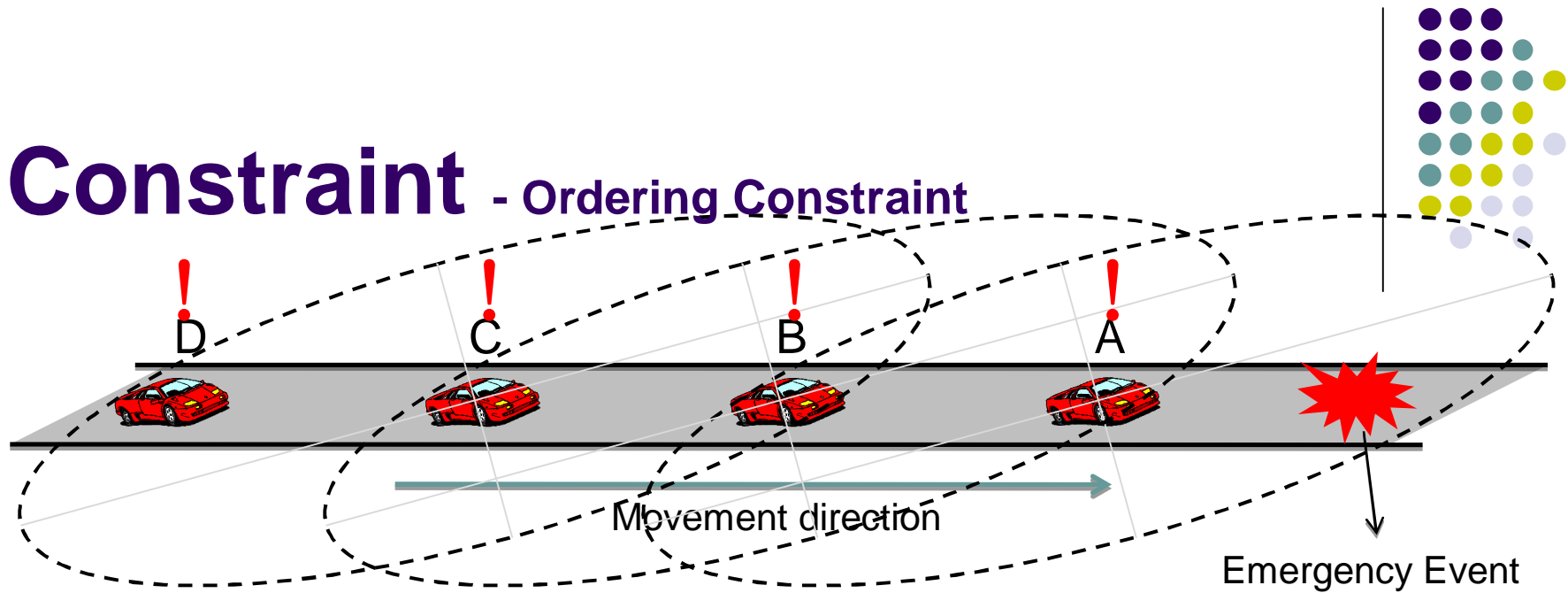
- | Slot allocation in VeSOMAC needs to satisfy the three constraint
 - | Timing Constraint
 - | Bitmap Constraint
 - | Ordering Constraint
 - | If two vehicles i and j are geographical neighbors and i 's location is ahead of j in the platoon, then i 's chosen slot should be earlier than j 's slot in the time domain.

Constraint - Ordering Constraint



Need 3
Frames

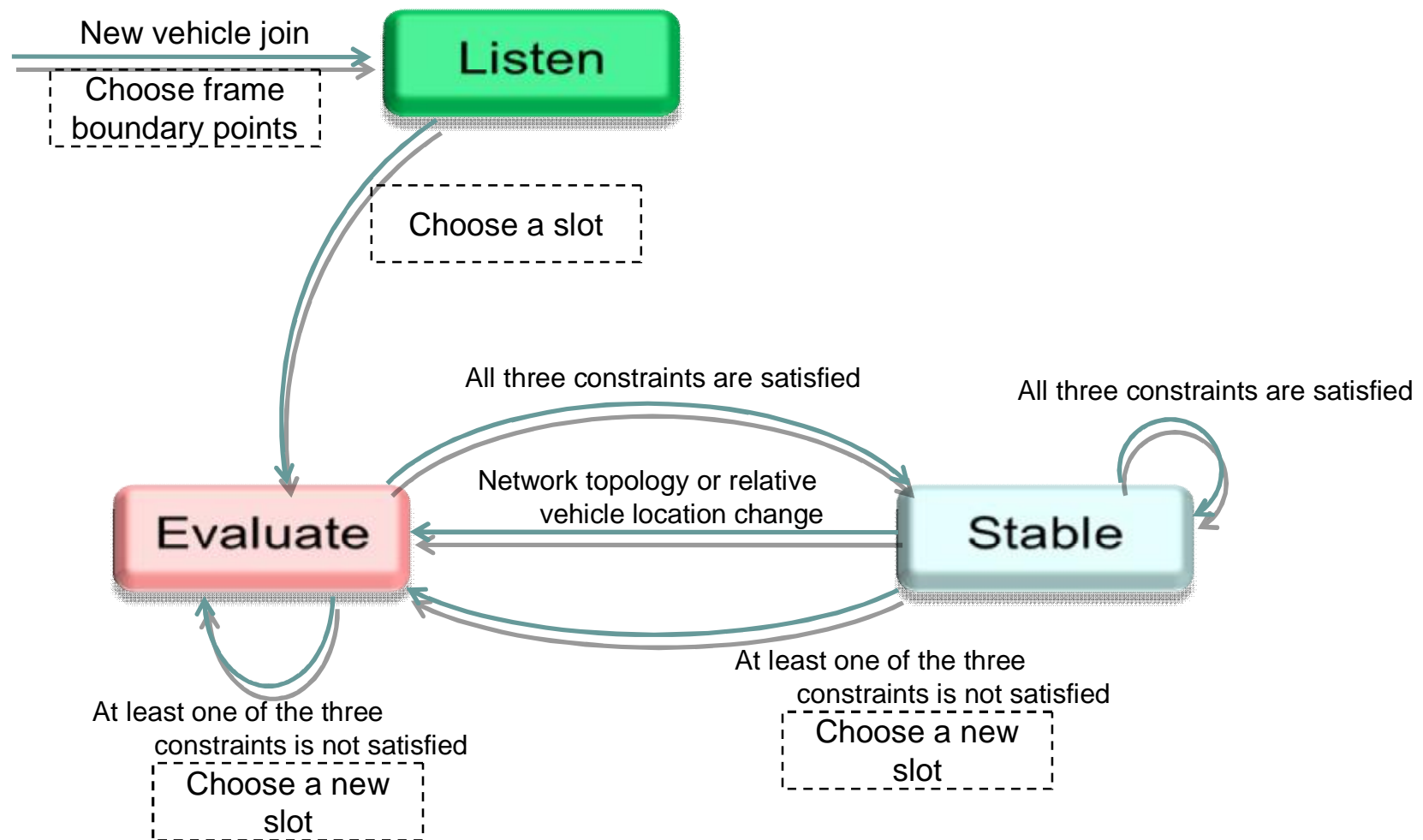
Constraint - Ordering Constraint



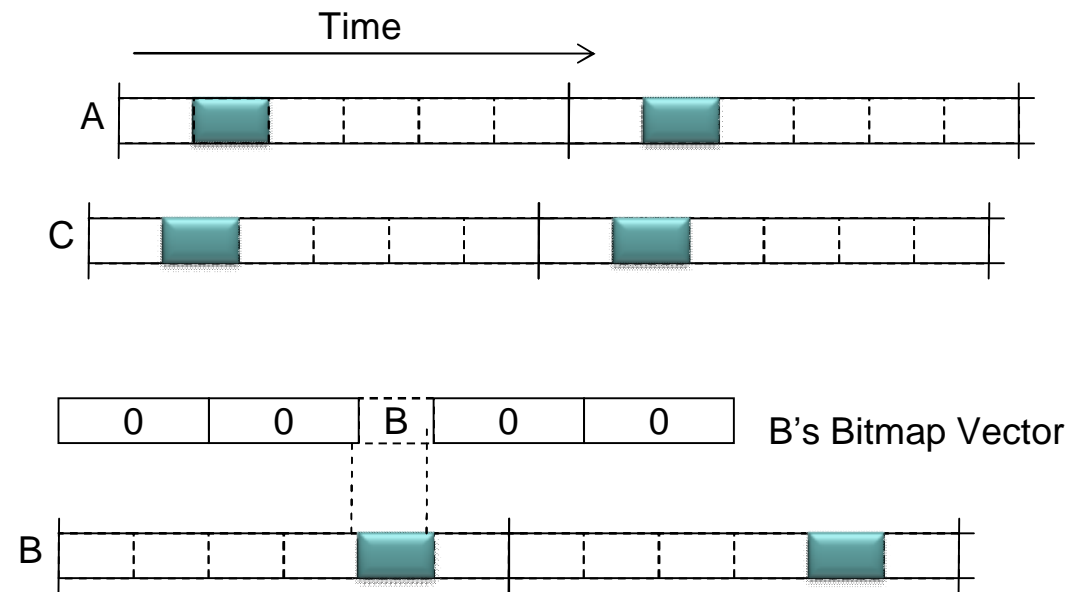
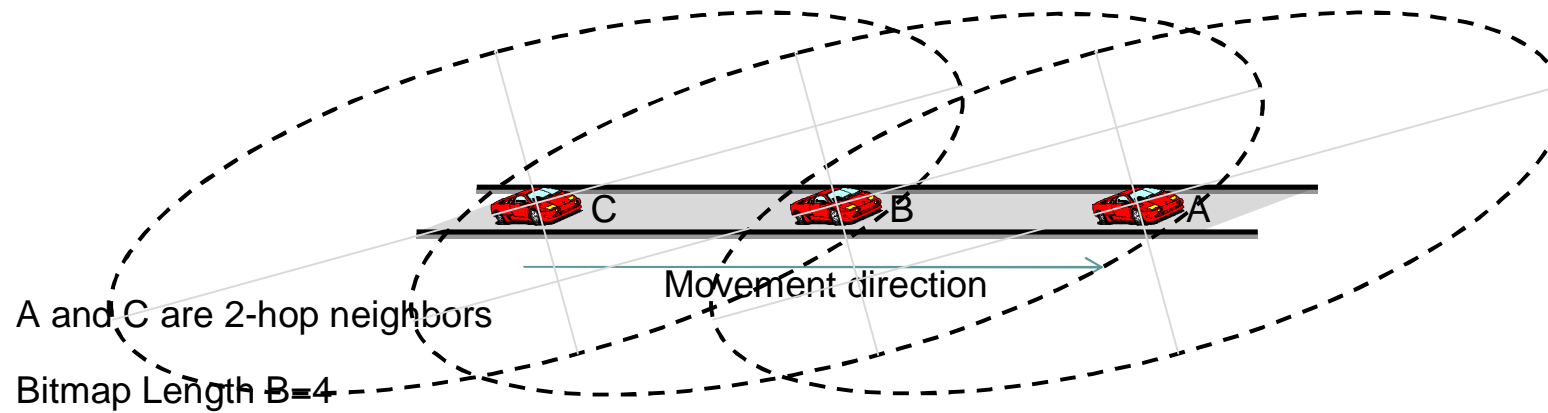
Need 1 Frame



Protocol State Machine



Collision Detection and Resolution

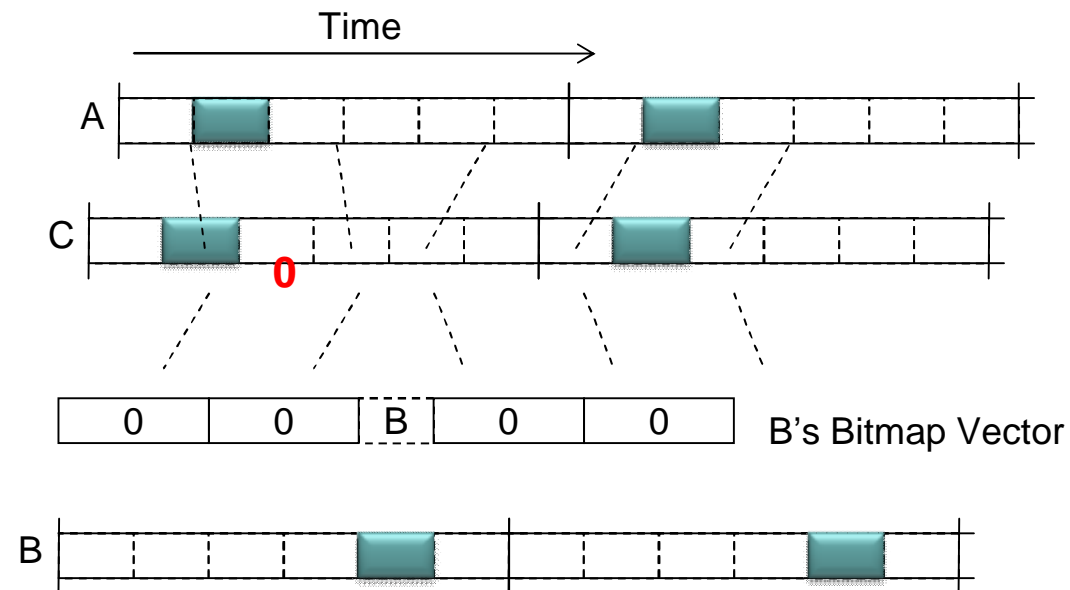


Collision Detection and Resolution



A and C are 2-hop neighbors

Bitmap Length B=4





Performance Evaluation

Simulation	
Simulator	ns-2
ITS application	CCA (Cooperative Collision Avoidance)

Simulation parameters (Vehicle Related)	
Platoon Size	50 vehicles
Vehicle Speed	68mph (= 108Km/hr = 30m/sec)
Inter-vehicle Spacing	25m to 45m (= 0.8sec to 1.5sec)
Vehicle Length	4m
Emergency Deceleration	8m/s ² (56.25m)
Regular Deceleration	4m/s ² (115.25m)
Drivers' Reaction Time	0.75sec to 1.5sec



Performance Evaluation

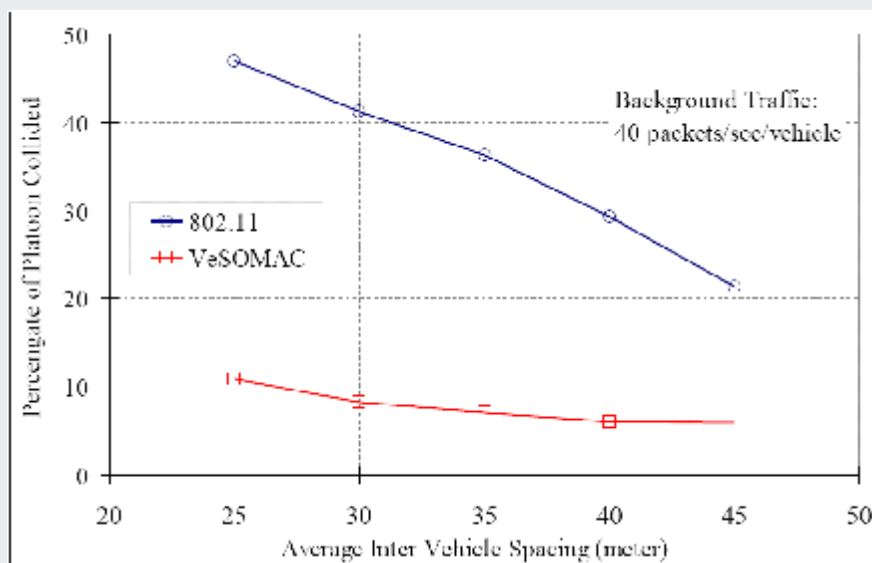
Simulation parameters (Network Related)	
Channel	DSRC 5.9GHz band, 24Mbps
Radio Range	300m
MAC Protocols	IEEE 802.11 and Worst case <i>VeSOMAC</i>
WCW Packet Size	300bytes (0.1ms)
WCW Message Period	100ms
<i>VeSOMAC</i> Frame Size	100 packets (10ms)
<i>VeSOMAC</i> Bitmap Size	96 Bits (very weak bitmap constraint)
<i>VeSOMAC</i> Evaluation Time	$W = 3$ frames



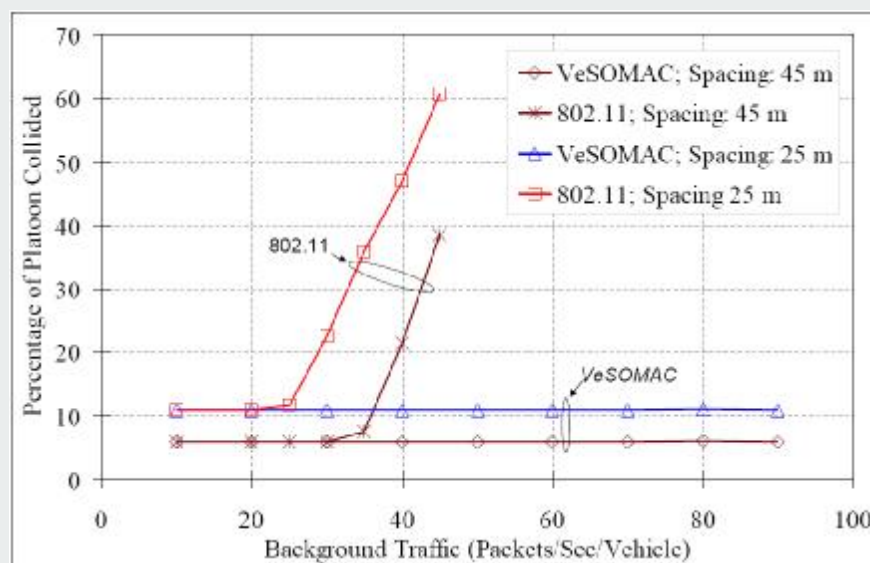
Performance Evaluation

Vehicle Crash Performance

Effects of Vehicle Spacing



Effects of Background Traffic

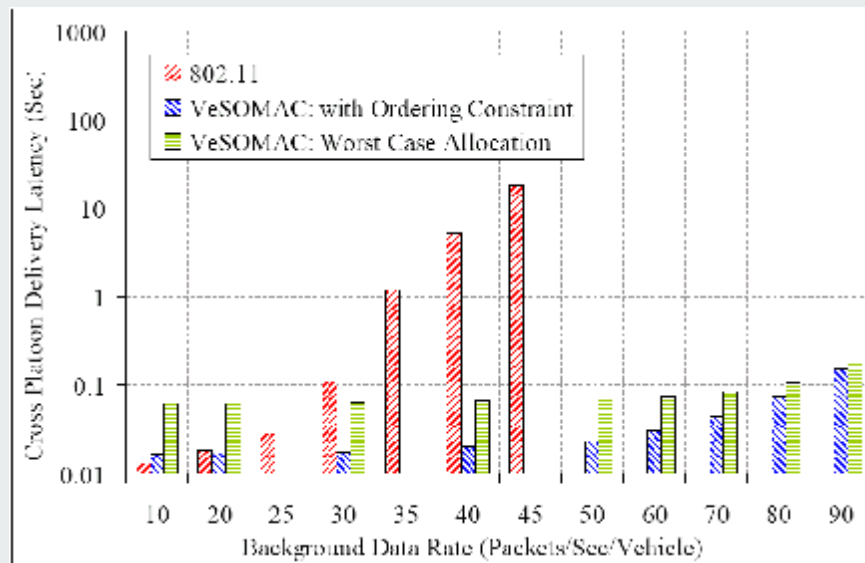


Performance Evaluation

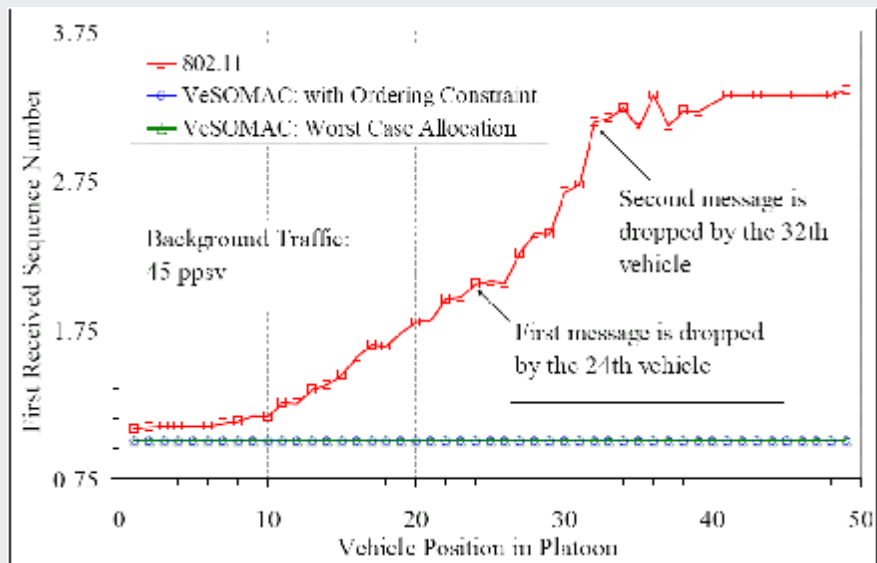


Data Plane Network Performance

Cross-Platoon Message Delivery



Packet Drops

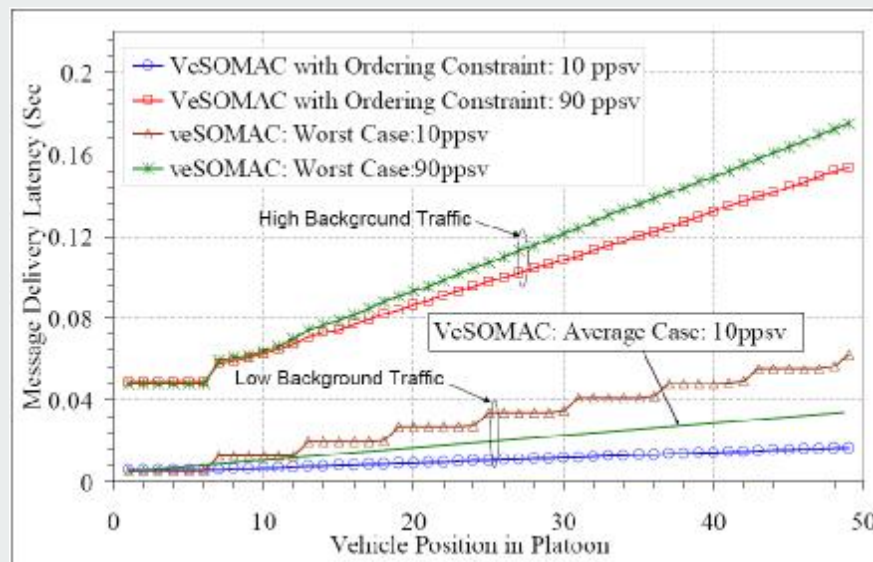




Performance Evaluation

Data Plane Network Performance

Cross-Platoon Message Delivery



Conclusion



- | Distributed TDMA allocation.
- | A bitmap vector is used for exchanging relative slot timing information across the 1-hop and 2-hop neighbor vehicles.
- | Can work without network time synchronization.
- | Minimizing the multi-hop delivery delay of ITS safety messages.
- | The protocol convergence during topology changes is fast.



Thanks~~~