

Downlink Packet Scheduling in LTE Cellular Network-Key Design Issues and a Survey

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Outline

- Abstract
- Introduction
- Overview on LTE Network
- Scheduling in LTE System
- Scheduling Strategies for LTE Downlink
- New Direction and Future Challenges
- Conclusions and Lesson Learned



Abstract

- LTE systems represent an important milestone towards the so called 4G cellular networks.
- This paper provides an overview on the key issues that arise in the design of a resource allocation algorithm for LTE networks.



Introduction

- The Radio Resource Management (RRM) block exploits a mix of advanced MAC and Physical functions.
- The design of effective resource allocation strategies becomes crucial :
 - 1. meet the system performance targets.
 - 2. satisfy user needs according to specific Quality of Service (QoS) requirements.
- The downlink shared channel is the key facets of LTE scheduling.



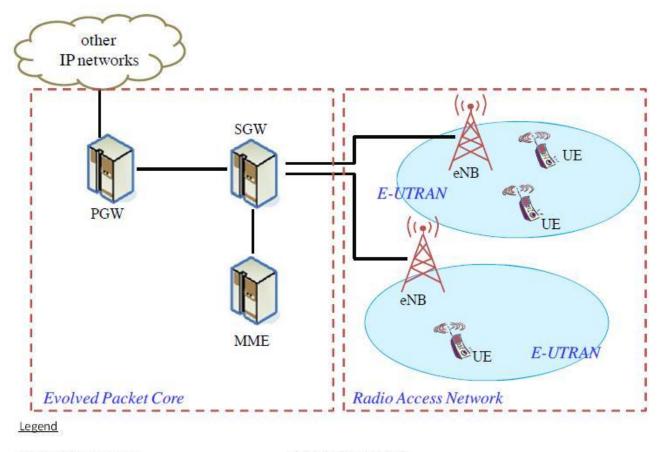
MAIN LTE PERFORMANCE TARGETS

Peak Data Rate	- Downlink: 100 Mbps					
	- Uplink: 50 Mbps					
Spectral Efficiency	2 - 4 times better than 3G systems					
Cell-Edge Bit-Rate	Increased whilst maintaining same site locations as deployed today					
User Plane Latency	Below 5 ms for 5 MHz bandwidth or higher					
Mobility	- Optimized for low mobility up to 15 km/h					
	- High performance for speed up to 120 km/h					
	- Maintaining connection up to 350 km/h					
Scalable Bandwidth	From 1.4 to 20 MHz					
RRM	- Enhanced support for end-to-end QoS					
	- Efficient transmission and operation of higher layer protocols					
Service Support	- Efficient support of several services (e.g., web-browsing, FTP, video-streaming, VoIP)					
	- VoIP should be supported with at least a good quality as voice traffic over the UMTS network					



A. System Architecture and Radio Access Network

• The LTE system is based on a flat architecture.



PGW: Packet Gateway MME: Mobility Management Entity SGW: Service gateway E-UTRAN: Evolved Universal Terrestrial Radio Access Network



B. Radio Bearer Management

- A radio bearer is in charge of managing QoS provision on the E-UTRAN Interface.
- When an UE joins the network, a default bearer is created for basic connectivity and exchange of control messages.
- Dedicated bearers are set up every time a new specific service is issued.

QCI	Resource Type	Priority	Packet Delay Budget [ms]	Packet Loss Rate	Example services					
1	GBR	2	100	10^{-2}	Conversational voice					
2	GBR	4	150	10^{-3}	Conversational video (live streaming)					
3	GBR	5	300	300 10 ⁻⁶ Non-Conversational video (buffered						
4	GBR	3	50	Real time gaming						
5	non-GBR	1	100	IMS signaling						
6	non-GBR	7	100	10^{-3}	Voice, video (live streaming), interactive gaming					
7	non-GBR	6	300	10^{-6}	Video (buffered streaming)					
8	non-GBR	8	300	10^{-6}	TCP based (e.g., WWW, e-mail), chat, FTP, P2P file sharing					
9	non-GBR	9	300	10^{-6}						

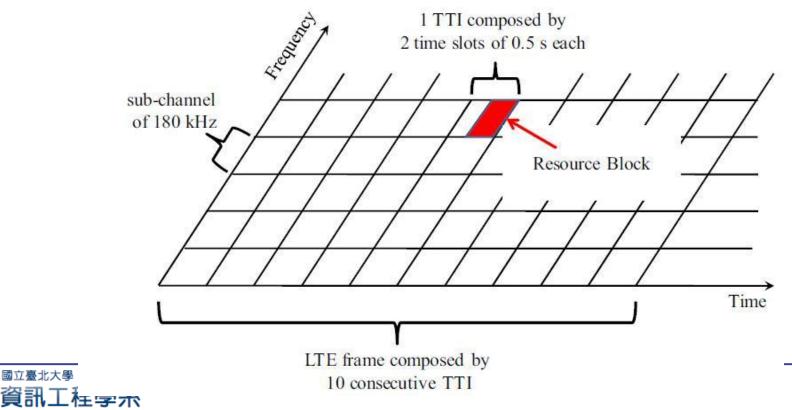
STANDARDIZED QOS CLASS IDENTIFIERS FOR LTE

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C. Physical layer

- Radio spectrum access is based on the Orthogonal Freq. Division Multiplexing (OFDM) scheme.
 - Single Carrier Freq.Multiple Access (SC-FDMA) and OFDMA are used in uplink and downlink directions.



- LTE radio interface supports two types of frame structure
 - Frequency Division Duplex(FDD):allowing simultaneous downlink and uplink data transmissions
 - Time Division Duplex (TDD): divided into two consecutive half-frames
- The selection of the TDD frame configuration is performed by the RRM module.

	sub-frame number									
	1^{st} half frame			2^{nd} half frame						
configuration number	0	1	2	3	4	5	6	7	8	9
0	D	S	U	U	U	D	S	U	U	U
1	D	S	U	U	D	D	S	U	U	D
2	D	S	U	D	D	D	S	U	D	D
3	D	S	U	U	U	D	D	D	D	D
4	D	S	U	U	D	D	D	D	D	D
5	D	S	U	D	D	D	D	D	D	D
6	D	S	U	U	U	D	S	U	U	D

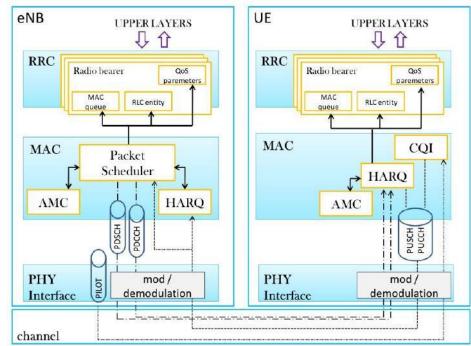
TDD FRAME CONFIGURATIONS

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D. Radio Resource Management

• 1. CQI reporting

- CQI enables the estimation of the quality of the downlink channel at the eNB.
- Each CQI is calculated as a quantized and scaled measure of the experienced Signal to Interference plus Noise Ratio (SINR).



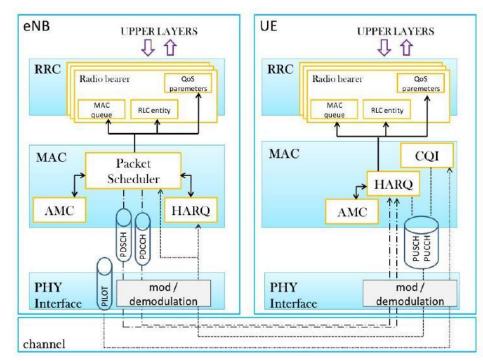
Legend

AMC: Adaptive Modulation and Coding RLC: Radio Link Control PDCCH: Physical Downlink Control Channel PUSCH: Physical Uplink Schared Channel CQI: Channel Quality Indicator

HARQ: Hybrid Automatic Repeat Request RRC: Radio Resource Control PDSCH: Physical Downlink Shared Channel PUCCH: Physical Uplink Control Channel



- 2. AMC and Power Control:
 - AMC module selects the proper Modulation and Coding Scheme (MCS) trying to maximize the supported throughput with CQI and a given target Block Error Rate (BLER).



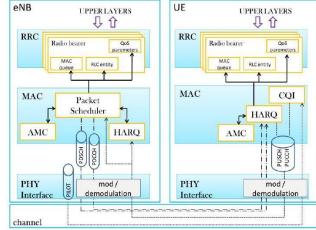
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- 3. Physical Channels:
 - Downlink data are transmitted by the eNB over the Physical Downlink Shared Channel(PDSCH).
 - Physical Downlink Control Channel (PDCCH) carries assignments for downlink resources and uplink grants, including the used MCS.
 - The Physical Uplink Control Channel (PUCCH) and the Physical Uplink Shared Channel (PUSCH) are defined in the uplink direction.
 - Due to single carrier limitations, simultaneous transmission on both channels is not allowed.
 Image: Competitive competitintex competitive competitive competitive competitive competiti



HARQ: Hybrid Automatic Repeat Request RRC: Radio Resource Control

PDSCH: Physical Downlink Shared Channel

PUCCH: Physical Uplink Control Channel

Legend

AMC: Adaptive Modulation and Coding

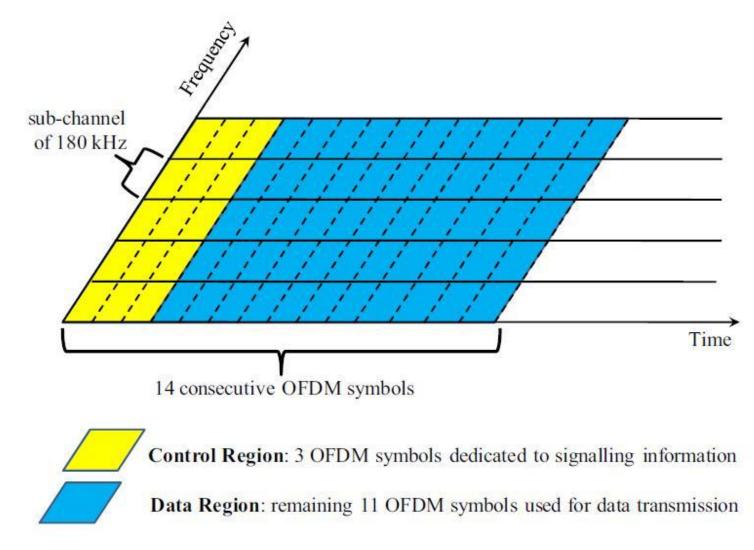
PDCCH: Physical Downlink Control Channel

PUSCH: Physical Uplink Schared Channel

RLC: Radio Link Control

CQI: Channel Quality Indicator







- 4. HARQ(Hybrid Automatic Repeat request)
 - It is the retransmission procedure at MAC layer, based on stopand-wait algorithm.
 - A NACK is sent over the PUCCH when a packet transmitted by the eNB is unsuccessfully decoded at the UE.
 - the eNB will perform a retransmission, sending the same copy of the lost packet.
 - the UE will try to decode the packet combining the retransmission with the original version, and will send an ACK message to the eNB upon a successfully decoding.



Scheduling in LTE System

- Multi-user scheduling is one of the main feature in LTE systems.
- Resource allocation for each UE is usually based on the comparison of per-RB metrics : transmission priority of each user on a specific RB

 $m_{j,k} = \max_i \{m_{i,k}\}$



- Status of transmission queues
- Channel Quality
- Resource Allocation History
- Buffer State
- Quality of Service Requirements



A. Key design aspects

- Complexity and Scalability
- Spectral efficiency
- Fairness
- QoS Provisioning



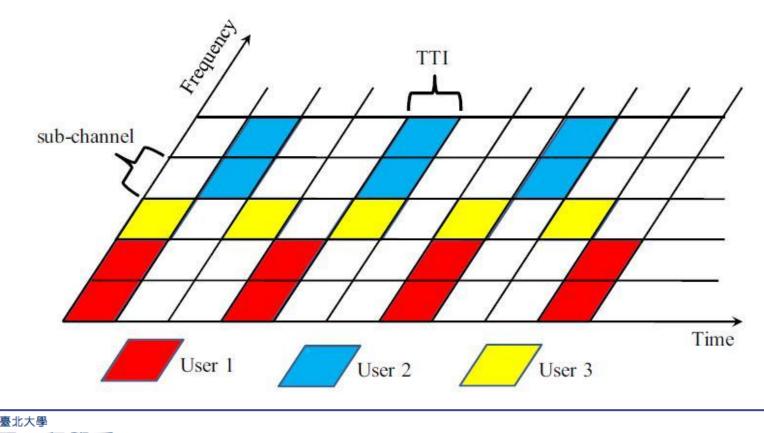
B. Practical limitations in real LTE systems

- Uplink Limitations
- Control overhead
- Limitation on multi-user diversity gain
- Energy Consumption: Discontinuous Reception (DRX) methods



C. Persistent and semi-persistent scheduling

• Dynamic frequency domain strategies have the main benefit of exploiting multi-user diversity gain, but this comes at the cost of increased control overhead.



Scheduling Strategies for LTE Downlink

- A. Channel-unaware Strategies
- 1. First In First Out
 - t : current time
 - T_i : the time instant when the request was issued by the i-th user

$$m_{i,k}^{FIFO} = t - T_i$$

• very simple, but both inefficient and unfair.



- 2. Round Robin
 - t : current time
 - T_i : last time when the user was served

$$m_{i,k}^{FIFO} = t - T_i$$

 The concept of fairness is related to the amount of time in which the channel is occupied by users but not in terms of user throughput.



• 3. Blind equal throughput

$$m_{i,k}^{BET} = 1/\overline{R^i}(t-1)$$

with

$$\overline{R^i}(t) = \beta \overline{R^i}(t-1) + (1-\beta)r^i(t)$$

where $0 \le \beta \le 1$.

NOTATIONS USED FOR SCHEDULING METRICS

Expression	Meaning					
$m_{i,k}$	Generic metric of the i -th user on the k -th RB					
$r^i(t)$	Data-rate achieved by the i -th user at time t					
$\overline{R^i}(t)$	Past average throughput achieved by the i -th user until time t					
$\overline{R^i_{sch}}(t)$	Average throughput achieved by data flow of the i -th user when scheduled					
$D_{HOL,i}$	Head of Line Delay, i.e., delay of the first packet to be transmitted by the i -th user					
$ au_i$	Delay Threshold for the <i>i</i> -th user					
δ_i	Acceptable packet loss rate for the <i>i</i> -th user					
$d^i(t)$	Wideband Expected data-rate for the i -th user at time t					
$d_k^i(t)$	Expected data-rate for the i -th user at time t on he k -th RB					
Γ^i_{k}	Spectral efficiency for the <i>i</i> -th user over the <i>k</i> -th RB					



- 4. Resource preemption
 - This approach can be fruitfully exploited to handle the differentiation among QoS (high priority) and non-QoS flows



- 5. Weighted Fair Queuing
 - $m_{i,k}^{RR}$ is the RR specific metric for the i-th user

$$m_{i,k}^{WFQ} = w_i \cdot m_{i,k}^{RR}$$

• No starvation is possible because the RR metric would control that the waiting time of a given user does not indefinitely grow.



- 6. Guaranteed Delay
 - require that each packet has to be received within a certain deadline to avoid packet drops.

$$m_{i,k}^{EDF} = \frac{1}{(\tau_i - D_{HOL,i})} \ .$$

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- 7. General considerations on channel-unaware strategies
 - We referred to them as channel-unaware schedulers, because their working rationales do not account for channel quality variations, making them unsuitable in cellular networks.



• 1. Maximum Throughput

- maximizing the overall throughput by assigning each RB to the user that can achieve the maximum throughput (indeed) in the current TTI.
- dⁱ_t(t) : the achievable throughput expected for the i-th user at the t-th TTI over the k-th RB

$$m_{i,k}^{MT} = d_k^i(t).$$

• MT is obviously able to maximize cell throughput, but it performs unfair resource sharing since users with poor channel conditions.



- 2. Proportional Fair Scheduler
 - MT : Maximum Throughput
 - BET : Blind equal throughput

$$m_{i,k}^{PF} = m_{i,k}^{MT} \cdot m_{i,k}^{BET} = d_k^i(t) / \overline{R^i}(t-1) \ . \label{eq:minimum}$$

It find a trade-off between requirements on fairness and spectral efficiency.



- 3. Throughput to Average
 - $d_t^i(t)$: the achievable throughput expected for the i-th user at the t-th TTI over the k-th RB
 - $d^i(t)$: the achievable throughput expected for the i-th user at the t-th TTI over all the bandwidth.

$$m_{i,k}^{TTA} = \frac{d_k^i(t)}{d^i(t)}.$$

- It quantifies the advantage of allocating a specific RB, guaranteeing that the best RBs are allocated to each user.
- It exploits channel awareness for guaranteeing a minimum level of service to every user.



- 4. Joint Time and Frequency domain schedulers
 - i. at first, a Time Domain Packet Scheduler (TDPS) selects a subset of active users in the current TTI among those connected to the eNB.
 - ii. Then, RBs are physically allocated to each user by a FDPS.
 - The computational complexity at the FDPS is reduced, due to the number of candidate users for resource allocation decreases.
 - For each phase a different policy can be selected.

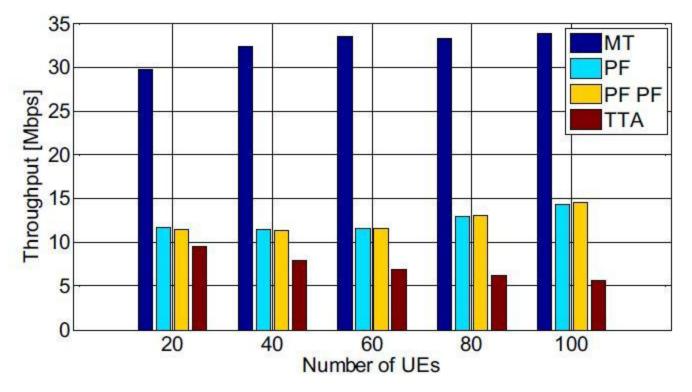


• 5. Buffer-aware Schedulers

- Buffer-Aware Traffic-Dependent (BATD) deal with the packet dropping probability due to a receiver buffer overflow.
- BATD makes use of buffer status information reported by the user to the eNB and of traffic statistics for setting dynamic priorities associated to each MAC queue.

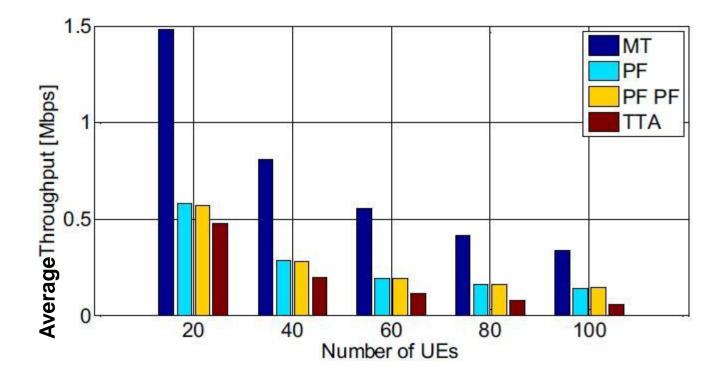


 6. Performance evaluation of most relevant channelaware/QoS-unaware strategies



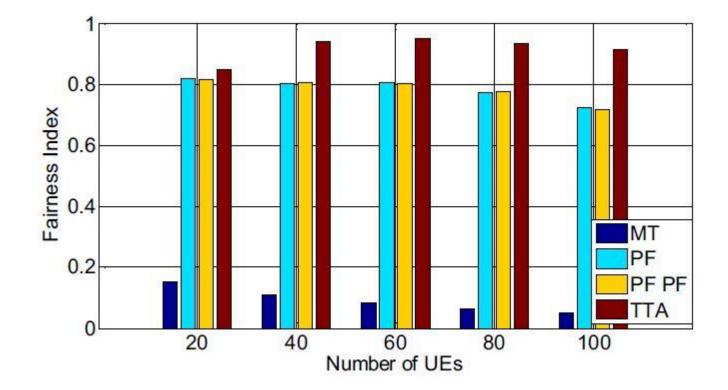
Aggregate cell throughput for different number of users





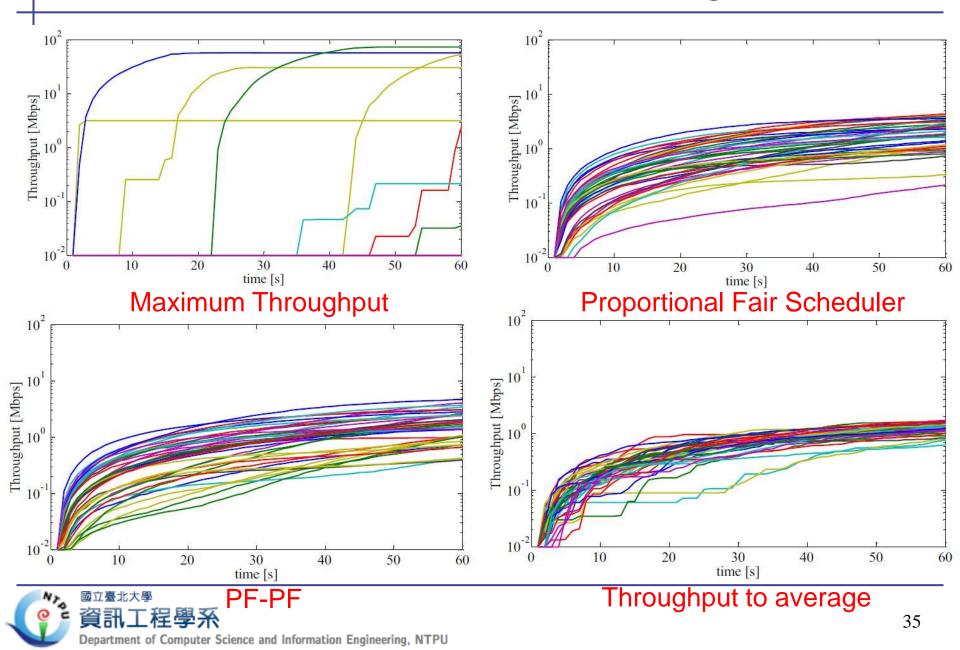
Average user throughput for different number of users





Fairness index for different number of users





- 7. General considerations on channel-aware/QoS-unaware strategies
 - Channel-awareness is a fundamental concept for achieving high performance in a wireless environment.



- 1. Schedulers for Guaranteed Data-Rate
 - It works in both time and frequency domains
 - For the time domain, users with flows below their target bitrate form a high priority set, the rest of the users forms a lower priority set.
 - Users belonging to first and second sets are managed by using Blind equal throughput(BET) and Proportional Fair Schedule (PF) algorithms
 - The approach use ordered lists to prioritize the most delayed flows and to meet the FDPS is QoS-unaware.



- 2. Schedulers for Guaranteed Delay Requirements
- non real-time : Modified Largest Weight Delay First(M-LWDF)
 - $D_{HOL,i}$: Head of Line Delay

$$m_{i,k}^{M-LDWF} = \alpha_i D_{HOL,i} \cdot m_{i,k}^{PF} = \alpha_i D_{HOL,i} \cdot \frac{d_k^i(t)}{\overline{R^i(t-1)}}$$

 M-LWDF uses information about the accumulated delay for shaping the behavior of PF, assuring a good balance among spectral efficiency, fairness, and QoS provisioning.



- Real-time : Exponential/PF (EXP/PF)
 - N_{rt} : the number of active downlink real-time flows
 - EXP/PF distinguishes between real-time and best effort flows

$$m_{i,k}^{EXP/PF} = \exp\left(\frac{\alpha_i D_{HOL,i} - \chi}{1 + \sqrt{\chi}}\right) \cdot \frac{d_k^i(t)}{\overline{R^i}(t-1)}$$

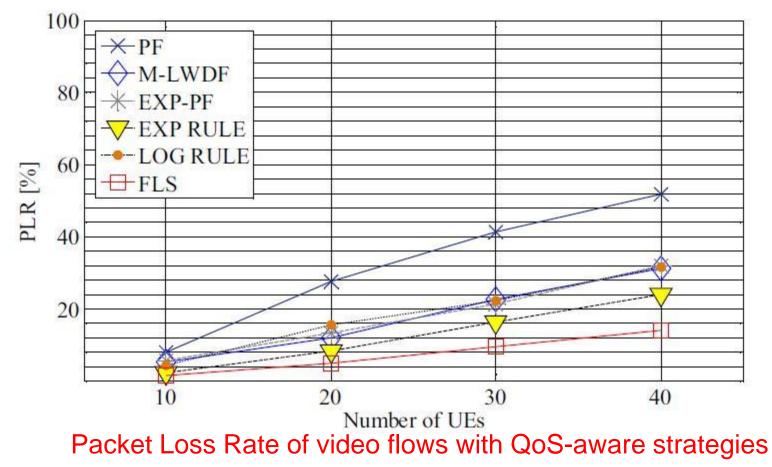
where

$$\chi = \frac{1}{N_{rt}} \sum_{i=1}^{N_{rt}} \alpha_i D_{HOL,i}$$

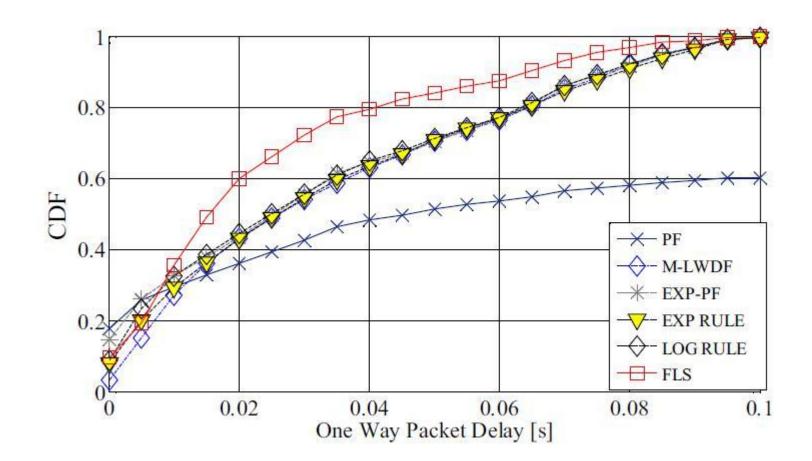
 They both try to guarantee good throughput and an acceptable level of fairness.



 3. Performance evaluation of most relevant channelaware/Qos-aware strategies









- 4. General considerations on channel-aware/Qos-aware strategies
 - support of multimedia flows is very important in LTE.
 - In fact, the working rationales of many of them are based on parameter settings whose optimality might depend on the specific scenario.
 - Further, more evolved approaches risk to be too complex and to waste resources.



D. Semi-persistent scheduling for VoIP support

- D. Semi-persistent scheduling for VoIP support
 - the radio resources are divided in several groups of RBs, and each block is pre-configurated and associated to VoIP.
 - It might be convenient to interleave persistent approaches, specific for VoIP, with dynamic ones that should perform allocation decisions for other flows.



E. Energy-aware strategies

• E. Energy-aware strategies

- The modification of resource allocation policies on a TTI basis may not have strong impact on energy performance of a cellular network, unless under very low traffic load.
- In this case, the best choice should be the maximization of the spectral efficiency.



New Direction and Future Challenge

- Its natural evolution, the Long Term Evolution-Advanced (LTE-A) solution, has been introduced by the 3GPP with the Release 10 for fulfilling and even surpassing IMT-A targets.
- Innovative technological solutions are standardized or foreseen for LTE-A, such as carrier aggregation, enhanced multi-antenna support, Coordinated Multi-Point (CoMP) transmission techniques, relaying, multi-userMultiple input multiple output (MIMO) communications, and Heterogeneous Networks (HetNets) deployment



Conclusion and Lesson Learned

- In this paper we provided an extensive survey on downlink packet allocation strategies recently proposed for LTE networks
- A main issue we found while studying literature on this topic, is the lack of a common reference scenario that can be used to compare different solutions.
- For this reason, we tried to identify a reference scenario and we used it to compare performance of the most representative solutions.

