

Design and Implementation of A SIP-Based Mobile and Vehicular Wireless Network with Push Mechanism

Yu-Chee Tseng, *Senior Member, IEEE*, Jen-Jee Chen, *Student Member, IEEE*, and Yu-Li Cheng

Abstract—As portable devices are gaining more popularity, maintaining Internet connectivity anytime and anywhere becomes critical, especially for mobile and vehicular networks. Network mobility (NEMO) and IP mobility are gaining more and more importance. In this work, we develop a SIP-based mobile network architecture to support network mobility for vehicular applications. We propose to form a mobile ad-hoc network (MANET) by the mobile hosts (MHs) inside a vehicle or a cluster of vehicles. The MANET is connected to the outside world via a SIP-based Mobile Network Gateway (SIP-MNG), which is equipped with one or multiple external wireless interfaces and some internal IEEE 802.11 interfaces. The external interfaces of the SIP-MNG are to support Internet connectivity by aggregating users' traffics to and from the Internet. In addition, exploiting session information carried by SIP signaling, the SIP-MNG supports resource management (RM) and call admission control (CAC) for the MHs. Wireless access, however, incurs charges, power consumption, and overheads of mobility management. So, it is desirable to allow the SIP-MNG to disconnect its external interfaces when necessary. To guarantee that users inside the mobile network will not lose any incoming request, we propose a *push mechanism* through short message services (SMS) to wake up these wireless interfaces in an on-demand manner. We show the detail signaling to support such mechanism. The proposed system is fully compatible with existing SIP standards. Our real prototyping experience and some experimental results are also reported.

Index Terms—MANET (Mobile Ad-Hoc Network), mobile computing, NEMO (Network Mobility), push mechanism, SIP (Session Initiation Protocol), wireless network.

I. INTRODUCTION

As mobile devices are gaining more popularity, users expect to connect to the Internet at anytime and from anywhere. Extensive research has focused on how to maintain the global reachability of a device without interruption even when it is moving around [1], [2], [3]. In [3], it has proposed methods for mobile devices to continue their sessions when address changes. However, these *host mobility* management schemes

manage mobility and connectivity of mobile devices in an individual manner. Group mobility is not addressed. For example, in a transportation carriage, there may exist tens/hundreds of users roaming together. If managing users in a one-by-one manner, not only the core network has to track each user individually, but also all users have to consume computation and radio transmission power to maintain their connectivity to the Internet. Clearly, supporting host mobility when users exhibit group mobility causes significant costs. Instead of tracking individual users' mobility, a concept called *Network Mobility (NEMO)* [4], [5] is proposed recently. NEMO packs all users in a vehicle as a network unit and conducts mobility management through a gateway/router in the mobile network. Supporting vehicular networking services by NEMO provides following advantages: 1) there is less power consumption for a mobile device to connect to the local gateway/router than to a base stations (BS) outside the vehicle; 2) the complexity of mobility management is lower or even transparent for a mobile when connecting to a local gateway/router; and 3) there are fewer handoffs since the MAC layer handoff only occurs on the central gateway/router rather than on all mobile devices and the mobility management of a single gateway/router can ensure the reachability of the whole group of users inside the gateway/router [6].

To support NEMO, the Internet Engineering Task Force (IETF) has created a working group, where a *Mobile IPv6-based NEMO (MIPv6-NEMO)* protocol is proposed [7], [8], [9]. In a mobile network, there is a central node called *Mobile Router (MR)*, which can connect to the Internet and provide networking services for the attached mobile devices. The MR has a home network and is assigned with one or more address blocks belonging to the home network, called *mobile network prefixes*, which can be assigned to the attached devices. Whenever the mobile network changes location, the MR will update with its Home Agent (HA), so packets addressed to the mobile network prefixes will still be routed to the mobile network from the home network. In addition to the MIPv6-NEMO, the *Session Initiation Protocol-based NEMO (SIP-NEMO)* is proposed in [10]. In SIP-NEMO, central to the mobile network is a *SIP Network Mobility Server (SIP-NMS)*, which is responsible for the mobility management of the mobile network. Each SIP-NMS has a *SIP Home Server (SIP-HS)*, which records the current location of the SIP-NMS and forwards requests to the SIP-NMS. SIP clients can communicate with each other directly without tunneling. So, SIP-NEMO also provides an additional way to achieve route

Y.-C. Tseng holds joint appointments at the Department of Computer Science, National Chiao-Tung University, Taiwan and the Department of Information and Computer Engineering, Chung-Yuan Christian University, Taiwan (email: yctseng@cs.nctu.edu.tw). His research is co-sponsored by Taiwan MoE ATU Program, by NSC grants 93-2752-E-007-001-PAE, 96-2623-7-009-002-ET, 95-2221-E-009-058-MY3, 95-2221-E-009-060-MY3, 95-2219-E-009-007, 95-2218-E-009-209, and 94-2219-E-007-009, by Realtek Semiconductor Corp., by MOEA under grant number 94-EC-17-A-04-S1-044, by ITRI, Taiwan, by Microsoft Corp., and by Intel Corp.

J.-J. Chen and Y.-L. Cheng are with the Department of Computer Science, National Chiao-Tung University, Hsinchu, Taiwan (email: chencz@csie.nctu.edu.tw; yulicheng.eic94g@nctu.edu.tw).

optimization and tunnel header reduction for network mobility. However, both MIPv6-NEMO and SIP-NEMO schemes have following shortcomings: 1) Wireless resource is limited and valuable, but both architectures do not consider how to manage wireless resources to guarantee QoS of sessions and how to dynamically adjust its bandwidth to the Internet. 2) Both methods use an always-on approach, which means that the central gateway/router has to continuously collect network advertisements and deal with handoffs even when the network has no sessions for a long period of time. This will incur unnecessary charges and energy consumption for the external wireless interfaces. 3) For SIP-NEMO, in addition to existing SIP servers, extra servers, such as SIP-HS and SIP-FS (SIP Foreign Server), are required. This will increase the implementation cost for public transportation operators. 4) Wide spread of SIP-FS deployment is needed in all foreign network domains that users may roam into. This causes significant barrier for such services to be widely accepted.

Also based on the NEMO concept, this paper proposes to combine the innate mobility and scalability of MANET with a SIP-based mobile network architecture to support vehicular networking services for users on the roads. Our system is designed over IPv4. Compared to SIP-NEMO [10], our system only needs an additional *SIP-based Mobile Network Gateway* (SIP-MNG), which follows the SIP standard and is compatible with the existing SIP framework, in each mobile network. No extra servers are needed in foreign networks. In our implementation, the SIP-MNG connects to the intra MANET by some IEEE 802.11 interfaces configured at the ad hoc mode and attaches to the Internet through one or more external wireless interfaces (such as GSM, GPRS, PHS, 3G, WiMAX, and IEEE 802.11 interfaces). These external interfaces allow the SIP-MNG to provide dynamic bandwidths to the Internet according to users' demand and roaming characteristics. For example, by exploiting the session parameters carried by SIP signaling (such as session type, codec, bandwidth requirement, ...), our SIP-MNG can support resource management (RM) and call admission control (CAC) to guarantee the QoS of sessions. Therefore, wireless resources are utilized efficiently.

Although the aforementioned system is quite attractive, dialing to the Internet (via wireless interfaces) incurs charges, power consumption, and maintenance overheads. When there is no Internet activity, it is desirable to disconnect the SIP-MNG's wireless interfaces to save energy and cost. However, this will result in users inside our system unreachable by other users. To solve this problem, we exploit the SIP session control feature and propose to establish a *push server* in the Internet as a representative of our system when being disconnected. When a SIP request arrives, we develop a *push mechanism* by dropping a short message to the SIP-MNG to "wake up" the connection by redialing to the Internet. Once the SIP-MNG resumes its connection, it informs the push server its contact address of the destined MH. Via the contact, the push server can transfer the suspended session to the MH by the SIP session transfer. Therefore, the proposed push mechanism can guarantee reachability of users in the mobile network, while saving energy, cost, and wireless resources.

The push mechanism has also been used in other applica-

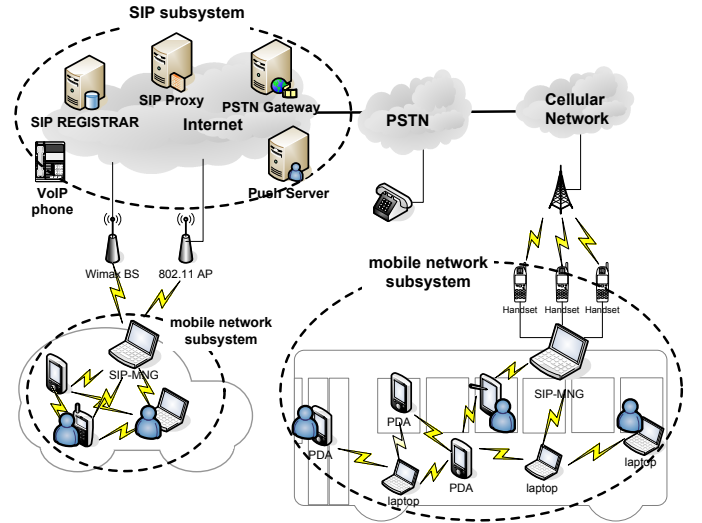


Fig. 1. System architecture.

tions. In GPRS networks, a MH must activate a PDP context before the corresponding service can be used. However, maintaining a PDP context without actually using it will consume network resources. So, short messages are used in [11] to activate a PDP context on-the-fly. For a dual-mode handset (with a cellular and a WLAN interfaces, for instance), to reduce energy consumption, it is desirable to disable the handset's WLAN module when it is not used. However, this will prevent the handset from receiving incoming VoIP calls from the WLAN interface. To solve this problem, reference [12] uses a paging mechanism via the cellular interface to inform the handset to activate its WLAN module. Then the VoIP call can be connected via the relatively inexpensive WLAN. However, both [11], [12] involve some modifications on SIP components and end devices to support such functions, which is undesirable. In addition, because of the overhead to go through the push procedure, both approaches may suffer from the timeout problem if the callee's response time exceeds a predefined threshold. In our architecture, the standard protocols in SIP servers and SIP clients are unchanged. Also, with an external server and SIP call control feature, our approach can effectively relieve the timeout problem.

This paper is organized as follows. Section II introduces our system architecture to support vehicular networking services and the design motivation. Detailed system operations are presented in Section III. The push mechanism is discussed in Section IV. Section V presents our prototyping experience and performance measurement results. Conclusions are drawn in Section VI.

II. SYSTEM ARCHITECTURE AND MOTIVATION

Fig. 1 shows our SIP-based mobile network architecture, which contains a *mobile network subsystem* and a *SIP subsystem*. The former is a SIP-based mobile network connecting to the Internet. The latter includes some servers to support SIP-based networking services.

Central to the mobile network subsystem is a SIP-based mobile network gateway (SIP-MNG). It is equipped with one

TABLE I
CHARGE PLANS OF DIFFERENT TYPES OF WIRELESS NETWORKS.

Type of Interfaces	Charge plans		
	By time	Flat rate	By packet
GPRS			✓
3G		✓	✓
PHS	✓	✓	
IEEE 802.11	✓	✓	
WiMAX	✓	✓	

or multiple wireless interfaces (such as GSM, GPRS, PHS, 3G, WLAN, and WiMAX interfaces) that can dial up to the Internet and some IEEE 802.11 interfaces to connect to the internal MANET. In real applications, cellular interfaces may be applied in freeways, country areas, and trains, while WLAN and WiMAX interfaces may be applied in hot-spot and metropolitan areas. The MANET consists of a set of mobile hosts (MHs), each equipped with an 802.11 interface configured at the ad hoc mode. Because real-time services are stringent in responsiveness, routing in the MANET is supported by a proactive protocol, such as DSDV [13] and CGSR [14], which will attempt to maintain consistent, up-to-date routing information at each host.

The SIP subsystem has four components: SIP registrar, SIP proxy, PSTN (Public Switched Telephone Network) gateway, and push server. The SIP registrar is a database containing users' subscription and status information. Users need to send their SIP registrar SIP REGISTER request messages to update their current locations periodically or when IP changes. The SIP proxy is responsible for routing SIP messages. The SIP registrar and SIP proxy are logical entities and are commonly implemented in a single host, called a *SIP server*. The PSTN gateway interconnects the Internet and the PSTN. So, with the PSTN gateway, SIP servers can set up/receive calls to/from the PSTN, respectively. The push server is to support our push mechanism and can "wake up" the SIP-MNG when its wireless interfaces are disconnected from the Internet. This can prevent incoming requests to the mobile network subsystem from loss.

Before presenting the detail system operations of our system, we first introduce our design motivations.

- **Saving charges of Internet access:** Table I shows the charge plans of different types of wireless interfaces. There are 3 types: 1) charge by time, 2) flat rate, and 3) charge by packet. Clearly, for charge plans 1) and 2), accessing the Internet for a group of users in a vehicle through a few wireless interfaces can save a great deal of charges for individuals. For plan 3), operators can save wireless resources.
- **QoS guarantee:** For realtime and multimedia sessions, QoS has to be guaranteed. Exploiting session information carried by SIP signaling, our SIP-MNG is designed with CAC and RM mechanisms to guarantee the QoS of realtime or multimedia applications.
- **Push mechanism to save charges, power consumption, and wireless resources:** Considering that the current battery technology for cell phones can operate 2~5 hours in the active mode, and 5~14 days in the standby mode, it is desirable to disconnect the cellular interfaces of

SIP-MNG (by putting them to the standby mode) to save power consumption when no Internet activity exists. Also, this can save charges for plans 1) and 3), too. Moreover, the SIP-MNG does not need to collect network advertisements to maintain global reachability. Via SIP session control and SMS, we design a push mechanism to allow the wireless interfaces to stay off-line when there is no Internet connection and to be "woken up" when necessary.

- **An added service for public transportation operators:** The public transportation operators can use our system as an added service to attract customers. In addition, adding management functions (such as accounting) becomes easy via SIP-MNG.
- **Backward compatibility:** Our goal is to serve existing SIP clients without modification. So we design our system by adding some new servers that can work transparent to existing SIP clients.
- **Reducing handoffs:** When a vehicle moves from one BS to another, BSs have to handle a large number of handoff events if host mobility is used. With our architecture, BSs only have to handle the mobility of the SIP-MNG, thus significantly reducing BSs' load.
- **Saving the power consumption of MHs:** Since a MH connects to the Internet via the local SIP-MNG, its power consumption is significantly reduced.
- **Decreasing the complexity of MHs:** Since handoff events are handled by the SIP-MNG, MHs do not have to do movement detection and location update. The design of MHs is simplified. This is also known as *movement transparency*.

III. BASIC OPERATIONS OF THE SIP-BASED MOBILE NETWORK

In this section, we discuss the basic network operations in our system. Since the MANET is considered a private network, to provide Internet access, the SIP-MNG serves as a NAT (Network Address Translation) server for MHs and is responsible for the translation of SIP messages. To achieve the NAT traversal for SIP, techniques such as ALG (Application Layer Gateway) [15], STUN (Simple Traversal of UDP Through NATs) [16], and ICE (Interactive Connectivity Establishment) [17] have been proposed. Here we adopt the ALG scheme in our SIP-MNG. Below, we will discuss the entrance and session establishment procedures and CAC, RM, and handoff mechanisms.

A. MH Joining the Mobile Network

When a mobile device moves into a vehicle with the proposed SIP-based mobile networking services, its IEEE 802.11 interface will detect the existence of the network. After attaching to the mobile network, the MH will get a new IP address from the SIP-MNG. With this address, the MH can actively send a SIP REGISTER message to its SIP server to update its contact information. The REGISTER message will trigger the SIP-MNG to serve as a representative of the MH. Fig. 2 shows the detail procedure of the SIP registration. In

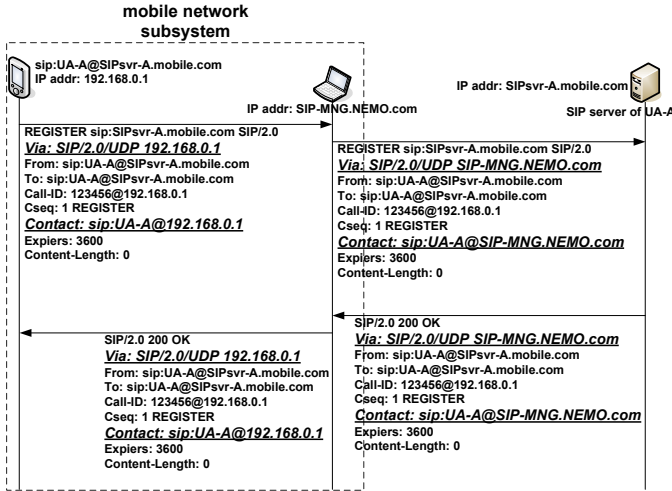


Fig. 2. The SIP registration procedure in our SIP-based mobile network.

this scenario, we assume the MH UA-A gets an IP address 192.168.0.1 from the SIP-MNG, and its SIP URL is sip:UA-A@SIPsvr-A.mobile.com. Addresses of the SIP-MNG and UA-A's SIP server are SIP-MNG.NEMO.com and SIPsvr-A.mobile.com, respectively. In the registration procedure, UA-A first sends a SIP REGISTER message to its SIP server with Via and Contact fields equal to 192.168.0.1 and sip:UA-A@192.168.0.1, respectively. Since the SIP-MNG monitors and translates all SIP messages to/from the Internet, it will capture the REGISTER message. On intercepting the message, the SIP-MNG will record UA-A's affiliation by adding UA-A's information into a *SIP_client_table* and relay the message by translating the Via and Contact fields into SIP-MNG.NEMO.com and sip:UA-A@SIP-MNG.NEMO.com, respectively. On receiving the SIP REGISTER message, the UA-A's SIP server will update UA-A's information and then reply a SIP 200 OK message. Since the Via field in the REGISTER request is translated into the SIP-MNG's address, the 200 OK message will be forwarded to the SIP-MNG. Also, with the SIP-MNG's address as the contact address, UA-A's SIP server can later forward SIP request messages destined for UA-A to the SIP-MNG, which will then relay them to UA-A. Upon the receipt of the 200 OK message from UA-A's SIP server, the SIP-MNG will translate the Via and Contact fields back to UA-A's address and relay the SIP response to UA-A. Then, the SIP registration procedure is completed.

B. Session Setup Procedure and CAC and RM Mechanisms

In this subsection, we present how a session is established between a MH in the MANET and an external CN (Corresponding Node). We also discuss how our SIP-MNG supports CAC and RM. Recall that a SIP-MNG has one or more external interfaces. An interface is *active* if it has been dialed up to the Internet; otherwise, it is *idle*. RM is responsible for evaluating the required bandwidth of a requested session and assigning a serving external interface for the session. The required bandwidth can be estimated by the SDP (Session Description Protocol) [18] provided in the SIP signal. If any active wireless interface has sufficient spare resource, RM will

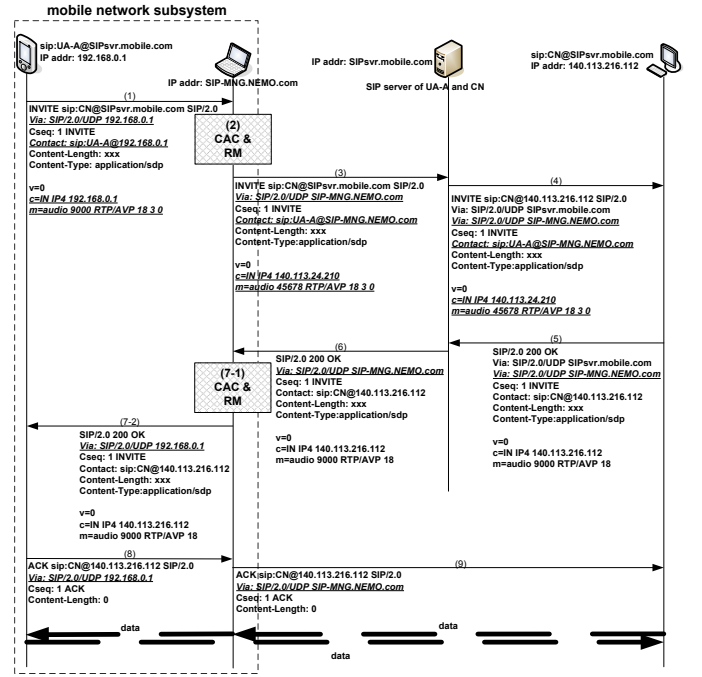


Fig. 3. The message flow to set up a session.

assign it to the session. Otherwise, RM will activate a new wireless interface to serve the session. If all wireless interfaces are full, CAC will drop the SIP signal and reject the request.

Here we use a voice call request as an example to describe the detail message flow of session setup, CAC, and RM. Fig. 3 depicts the message flow, with some SIP headers omitted for simplicity. We assume that UA-A's IP address and SIP URL are 192.168.0.1 and sip:UA-A@SIPsvr-A.mobile.com, respectively. CN's IP address and SIP URL are 140.113.216.112 and sip:CN@SIPsvr-A.mobile.com, respectively. For simplicity, UA-A and CN use the same SIP server, whose address is SIPsvr-A.mobile.com. The SIP-MNG's address is SIP-MNG.NEMO.com. The corresponding call setup steps when UA-A calls CN are as follows.

- 1) UA-A sends a SIP INVITE to the SIP server to invite CN. Fields *c* (connection information) and *m* (media description) in the SDP are set to the connection address (192.168.0.1) and port number (9000), respectively (by which UA-A can receive data from CN). Field *m* also describes that the requested session is audio, so RM can tell that this is a realtime/multimedia session.
- 2) When intercepting the SIP INVITE message, the SIP-MNG will initiate the CAC and RM procedures. CAC will accept the requested session if RM returns ok, and reject the session otherwise. RM includes 3 steps: 1) evaluate the required bandwidth; 2) allocate a wireless interface for the session if there is enough bandwidth; and 3) return the status ok/failure to CAC. Since the codec information can be derived from field *m* (18 = G.729, 3 = GSM, and 0 = G.711 mu-law) and the default packetization interval (PI) is 20 ms, the required bandwidth can be predicted. If there is no resource, the SIP-MNG will drop the SIP INVITE message and

respond a SIP "480 Temporarily not available" message to UA-A. Otherwise, the session will be recorded into a *session table* and the SIP INVITE will be relayed to the SIP server after translation. Since the SDP may provide multiple candidate codecs for the session, RM will reserve the maximum required bandwidth for the session. For non-audio/video sessions (which will be described in field m), the SIP-MNG may reserve one or few wireless interfaces dedicated to these types of sessions. For such best-effort sessions, RM will skip the resource evaluation step and directly allocate an interface for it.

- 3) Then fields Via, Contact, c, and m of the SIP INVITE message are translated. In this example, the (connection address):(port) is translated from 192.168.0.1:9000 to 140.113.24.210:45678.
- 4) On receiving the SIP INVITE message, the SIP server forwards it to the CN, which will register the caller's address as 140.113.216.112.
- 5) The CN then replies with a SIP 200 OK message, where it chooses G.729 as the codec to communicate with UA-A. According to the Via header in the received SIP INVITE message, this 200 OK message will be sent back to the SIP server at SIPsvr.mobile.com.
- 6) Upon the receipt of the 200 OK message, the SIP server forwards it to address SIP-MNG.NEMO.com.
- 7) When the SIP-MNG receives the 200 OK message, RM will update the session table according to the final media information. Then, the SIP-MNG translates the 200 OK message and relays it to UA-A.
- 8) UA-A feedbacks SIP ACK message to the CN directly according to the Contact information in the received 200 OK message.
- 9) The SIP-MNG relays the SIP ACK message to CN after translating the Via header. Then the call is established.

C. Handoff Procedure

As the vehicle moves, an external interface may change its network domain and the SIP-MNG may change its active interface. In both cases, a handoff happens. The SIP-MNG must recover sessions on these handoff interfaces by sending re-INVITE messages and maintain the global reachability of all MHs in the MANET by sending re-REGISTER messages. Since recovering on-going sessions is more urgent, it will be done first. Below, we outline the detail steps (refer to Fig. 4).

- 1) According to the SIP_client_table and the session table, the SIP-MNG retrieves the endpoint MHs and CNs of the on-going sessions on each handoff interface.
- 2) The SIP-MNG then sends each CN a SIP INVITE message to re-invite it.
- 3) A CN, on receiving the SIP INVITE message, will register the new contact address and port of the MH. Then the CN replies a SIP 200 OK message.
- 4) On receiving the 200 OK messages, the SIP-MNG will update the session table and reply the corresponding CN a SIP ACK message. Then, the session between the CN and the MH can be continued.

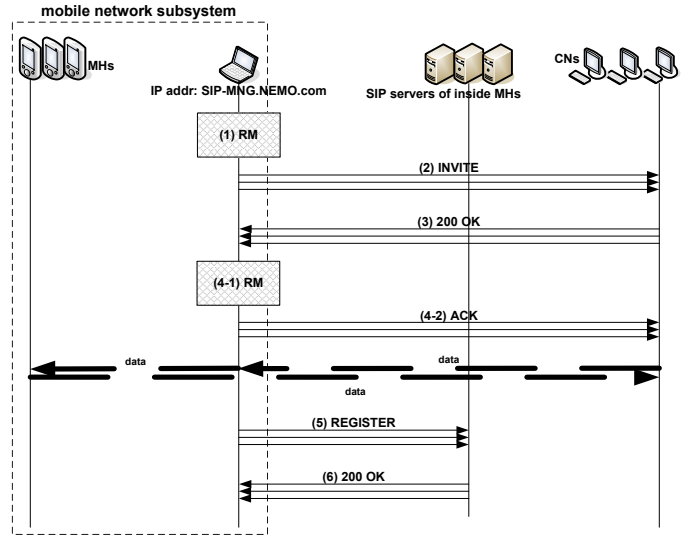


Fig. 4. The message flow of handoff.

- 5) After recovering all handoff sessions, the SIP-MNG will send a SIP REGISTER message for each MH that has changed to a new external interface to update its contact address.
- 6) When a SIP server receives the above SIP REGISTER message, it will update the corresponding MH's data and reply a SIP 200 OK message. After this, all MHs are guaranteed to be reachable from outside.

D. MH Leaving the Mobile Network

When a MH leaves the mobile network, it may detect other network and update its contact information by sending a SIP REGISTER message. If there is an on-going session, it can be resumed by SIP re-INVITE. However, since the MH does not deregister with the SIP-MNG, the SIP-MNG will keep its information in the SIP_client_table. If the MH has an on-going session before it leaves, the SIP-MNG will still reserve resource for the session. Fortunately, if a SIP request arrives, because the MH does not exist in the network, a new session will not be set up. However, the allocated resource will never be released. To solve this problem, we suggest to set a timer for each session and integrate the SIP-MNG with the underlying routing protocol in the MANET. If the SIP-MNG finds that the MH does not exist after the timer times out, the corresponding resource will be recycled. Noted that the SIP-MNG can determine whether a specific MH exists or not by searching the corresponding entry in the routing table. Alternatively, the SIP OPTIONS message can be used to query the capability of a SIP client or server. The message should be responded with a SIP 200 OK response with the supported capabilities. This can be used to determine a MH's existence.

IV. THE PROPOSED PUSH MECHANISM

Our system allows the external interfaces of SIP-MNGs to be disconnected when there is no Internet connection. When any interface is connected, the mobile network becomes a part of the Internet, so all sessions can be handled as usual. When

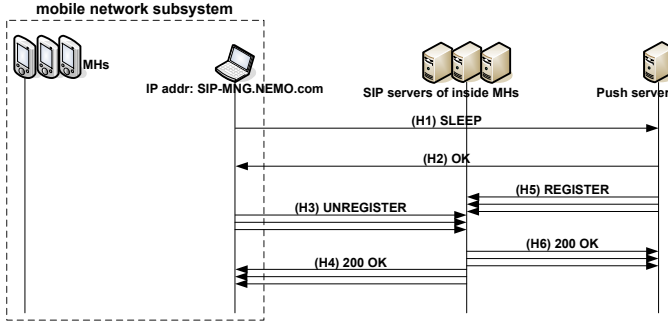


Fig. 5. Sleep procedure.

all interfaces are disconnected, if an outgoing SIP request is sent by a user in the mobile network, the SIP-MNG can buffer the invitation, dial up to the Internet, register this user with the SIP registrar, and then send the invitation to the callee. However, when the mobile network is out of reach from the Internet, an incoming SIP request can not be completed. To solve this problem, we propose a *push mechanism*.

In our push mechanism, the SIP-MNG will carry out a *sleep procedure* when it decides to disconnect from the Internet. Recall that there is a push server in the SIP subsystem (refer to Fig. 1). The SIP-MNG will solicit the push server as its agent during the disconnection period. When an incoming SIP request arrives, the push server will be notified first and will trigger a *wake-up procedure*, which includes a *wake-up process* and a *session transfer process*. In the wake-up process, the push server will activate the SIP-MNG via SMS (Short Message Service) and establish a connection with the caller to hold the session. After the SIP-MNG reconnects to the Internet, the transfer process will help build the link between the caller and the internal callee. In this way, the SIP-MNG can stay off-line but remain reachable from the Internet.

A. Sleep Procedure

To avoid modifying the existing SIP standard, the push server works as an agent for users inside the mobile network when the SIP-MNG is disconnected from the Internet. The sleep procedure is to inform the SIP server to redirect all SIP messages destined to mobile network users to the push server.

Fig. 5 shows the message flow of the sleep procedure. When there is no Internet connection, the SIP-MNG may initiate the sleep procedure by sending a sleep request to the push server (H1). The sleep request includes the SIP URIs of MHs inside the mobile network and the MSISDN (Mobile Station International ISDN Number) of one of the cellular interfaces of the SIP-MNG. The MSISDN will later be used by the push server to notify the SIP-MNG of new incoming SIP requests via short messages.

The push server maintains a *gateway table* and a *SIP client table*. Each entry in the gateway table is to track the status of one SIP-MNG and includes four fields: gateway id, status, MSISDN, and IP address. Each entry in the SIP client table is to track one MH and includes three fields: SIP URI, SIP-MNG id, and registration expiration time. On receiving the sleep request from the SIP-MNG, the push server updates

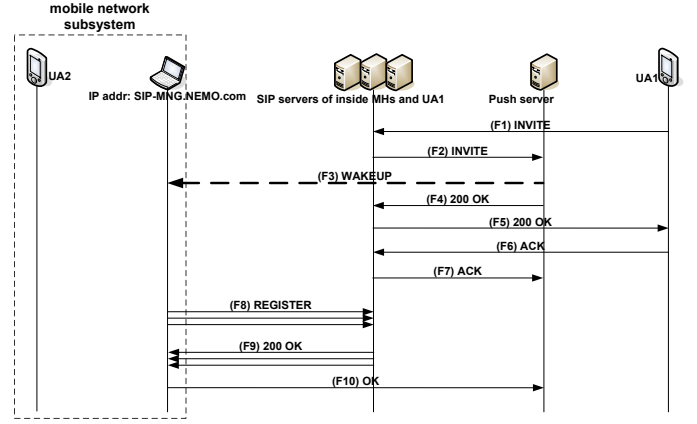


Fig. 6. Wake-up process.

these two tables and marks the status of the SIP-MNG as off-line. The push server then replies an OK to the SIP-MNG (H2). Upon receipt of the OK message, the SIP-MNG will generate REGISTER messages for all internal SIP clients to their corresponding SIP servers (H3) with an EXPIRE value of 0 (which means unregister). In return, the SIP server will reply SIP 200 OK messages (H4). On the other hand, as soon as the push server replies an OK message to the SIP-MNG (H2), it also sends REGISTER messages for all MHs served by the sleeping SIP-MNG to their SIP servers with a non-zero EXPIRE value (H5). These REGISTER requests should contain the push server's IP address in the CONTACT field, so that all future SIP INVITE requests to these MHs will be forwarded to the push server. Upon receipt of a REGISTER message, a SIP server will update its record and reply a SIP 200 OK message to the push server (H6). After step H4, the SIP-MNG can disconnect all its wireless interfaces, and after step H6 the push server will become the agent of the mobile network subsystem and send periodic REGISTER messages for it.

B. Wake-Up Procedure

Consider a SIP request from a SIP client UA1 in the Internet to a client UA2 in the mobile network. To complete this session, there are two parts in the wake-up procedure: wake-up process and session transfer process.

1) *Wake-Up Process*: Fig. 6 illustrates the message flow of the wake-up process. To set up a session to UA2, UA1 first sends a SIP INVITE message to UA2's SIP server containing the session information, connection address, and port number of UA1 in the SDP (F1). The SIP server will identify that the contact of UA2 is the push server and forward the INVITE to the push server (F2). The push server then checks its SIP client table and gateway table and retrieves UA2's SIP-MNG information. Because the status of the SIP-MNG is off-line, the push server will send a short message to the MSISDN registered by the SIP-MNG (F3). This short message carries the event type and IP address of the push server to inform the SIP-MNG to reconnect to the Internet. In the meantime, the push server will temporarily set up the session with UA1 to keep the session alive (F4–F7). This can prevent the SIP

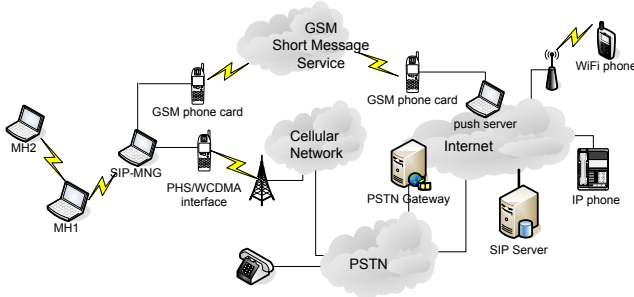


Fig. 8. A testing environment of our SIP-based mobile network system.

support multi-hop ad hoc routing, MNs in the MANET all run the Destination-Sequenced Distance-Vector (DSDV) protocol.

The testing environment is shown in Fig. 8. MHs are IBM X23 notebooks equipped with ASUS WL-167G USB2.0 WLAN adapters configured at the ad hoc mode. They run the Microsoft Windows XP with Windows messenger 5.1 as their SIP client software. The SIP-MNG is an IBM T42 notebook running the Linux Fedora Core Release IV; a Nokia card phone 2.0 is equipped to access GSM SMS. To access the Internet, the SIP-MNG is equipped with an embedded 802.11 WLAN chipset and a cellular interface (Huawei E612 WCDMA PCMCIA card/PHS J88 cellular phone). The PHS phone interface is driven by the SIP-MNG via a P-Card MC-6550 PCMCIA adaptor. The push server is an ASUS centrino notebook running the Microsoft Windows XP. It is also equipped with a Nokia card phone 2.0 to transmit short messages. An iptel SIP server [23] is used to serve as a server. In the Internet side, we have several SIP client terminals including D-Link IP phone and WiFi phone. All SIP clients inside the MANET or in the Internet use the SIP server as their home SIP server.

B. Call Setup Time and Maximum Number of Supported Calls

Based on the above environment, we have measured the call setup time and the maximum number of supported calls for different wireless interfaces. Some testing results are shown in Table II. Each experiment here and in the following subsections was repeated 10 times. It shows that, to set up a call via a PHS J88 phone interface, it takes 1.58 seconds in average for MH2 to receive the ringing tone of a call from an IP phone in the Internet. If via the WCDMA PCMCIA card/embedded IEEE 802.11b interface, it requires 1.44/1.11 seconds in average, respectively. We do not provide results of the GPRS interface because in our experience a GPRS interface cannot provide enough bandwidth to support even one single voice call (GPRS downlink bandwidth is only 28.8kbps, uplink bandwidth is even less). The *call setup time* by cellular interfaces is longer than that by 802.11 interfaces because connecting MH2 and the IP phone via a cellular interface has to go through four networks: the Internet, PSTN, cellular network, and our MANET. However, connecting them via a wireless interface only goes through two networks: the Internet and our MANET. The proposed system allows multiple calls to be supported by a single interface. With a PHS interface, the SIP-MNG can support two concurrent

TABLE II
MEASUREMENT OF CALL SETUP TIME AND CAPACITIES OF DIFFERENT INTERFACES.

Item	external interface			
	PHS	WCDMA	GPRS	802.11b
call setup time	1.58 sec	1.44 sec	—	1.11 sec
maximum number of calls (with G.729)	2	5	0	12

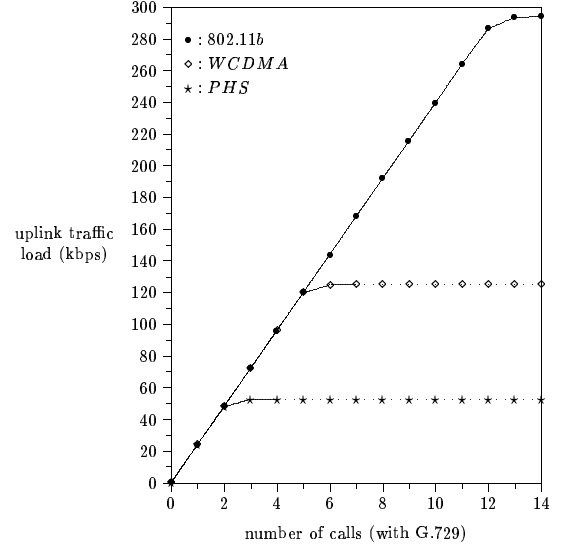


Fig. 9. Uplink traffic load against the number of concurrent calls using one PHS/WCDMA/802.11b interface.

voice calls with acceptable quality ($< 1\%$ packet dropping rate), while with one WCDMA interface, it can support up to five concurrent voice calls. Interestingly, with an IEEE 802.11b wireless interface, the SIP-MNG can support up to 12 concurrent voice calls with G.729 as the codec.

We have also measured the uplink traffic load on a PHS/WCDMA/802.11b interface (we did not measure downlink load because the uplink bandwidth is lower than the downlink bandwidth in all existing cellular interfaces). Fig. 9 plots the uplink traffic load against the number of concurrent calls on a PHS/WCDMA/802.11b interface. As shown in Fig. 9, in the beginning, the uplink traffic load increases linearly as the number of calls increases. As the number of calls keeps increasing, the PHS's curve will saturate after there are more than two calls, the WCDMA's curve will saturate after there are more than five calls, and the IEEE 802.11b's curve will saturate after there are more than 12 calls. This gives the maximum number of calls that can be supported in Table II. Our measured capacity for the 802.11b wireless interface is close to the analysis in [24], which claims that 11.4 G.729 calls can be supported in average.

C. Handoff Delay

Fig. 10 illustrates the handoff procedure when an IEEE 802.11 interface is used. The handoff delay can be divided into five parts: T1 (the delay of layer 2 handoff for a MH switching from one AP to another), T2 (the delay of layer 3 handoff for a MH getting a new IP in a new subnet), T3

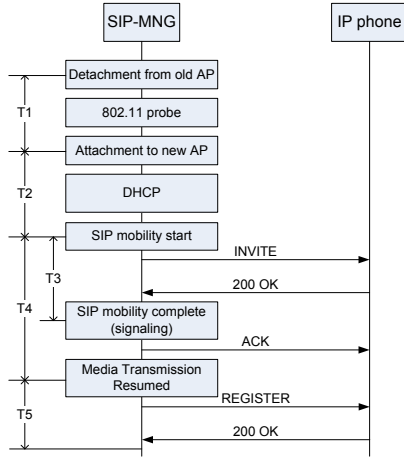


Fig. 10. The handoff procedure when an IEEE 802.11 interface is used.

TABLE III
MEASUREMENT OF HANDOFF LATENCIES.

Item	external interface	
	Cellular (PHS)	Wireless (802.11b)
T3	742.87 ms	268.26 ms
T4	967.37 ms	279.54 ms
T5	267.69 ms	42.34 ms

(the delay of SIP mobility), T4 (the delay from the SIP-MNG activating the SIP mobility procedure to the media session being resumed), and T5 (the re-registration latency for the SIP-MNG to update MHs' contact information).

T1 and T2 are not the focus in our experiments because handoff rarely happens for cellular interfaces. In addition, lots of previous work [25], [26] have already provided empirical results of WLAN layer 2 and layer 3 handoff latencies. Therefore, we only evaluate T3, T4, and T5 for cellular and 802.11 interfaces. Table III shows the values of T3, T4, and T5 for different interfaces. For cellular interfaces, T4 consumes 967.37 ms in average, which is somewhat long. Fortunately, for cellular interfaces, handoff seldom happens because an operator's network normally covers quite a large area. For cellular interfaces, T5 is 267.69 ms in average. This is why we do re-invitation before re-registration when a handoff occurs, or the SIP mobility latency will be further increased by 270 ms. For 802.11 interfaces, T3 and T4 are 268.26 ms and 279.54 ms, respectively, in average. This is much shorter than the cellular interface case. The re-registration delay T5 for such interfaces is only 42.34 ms in average.

D. Performance of the Push Mechanism

By the proposed push mechanism, our system allows the external interfaces of the SIP-MNG to be disconnected when there is no Internet connection. In this subsection, we evaluate the call setup latency when the SIP-MNG is disconnected from the Internet with our push mechanism. Table IV shows our testing results with different interfaces. If the interface is a PHS J88 phone (in the countryside areas), with our push mechanism, it costs 23.03 seconds in average for MH2 to receive the ringing tone from the IP phone. This latency consists of two major components: short message transmission time and

TABLE IV
LATENCIES OF THE PUSH MECHANISM.

Item	external interface		
	PHS	WCDMA	802.11b
call setup time (with push mechanism)	23.03 sec	18.77 sec	23.11 sec
short message transmission time	5.31 sec	5.31 sec	5.31 sec
wireless interface reconnection time	13.05 sec	8.47 sec	14.05 sec

wireless interface reconnection time. We found that the short message transmission time is 5.31 seconds in average and the wireless interface reconnection time takes 13.05 seconds in average. Replacing the PHS interface with a Huawei E612 WCDMA PCMCIA card, it shows that the call setup time can be reduced to 18.77 seconds in the disconnected case. To further divide the cost, Table IV shows that the short message transmission time is also 5.31 seconds and the wireless interface reconnection time is 8.47 seconds in average. If the wireless interface is an IEEE 802.11 interface, with our push mechanism, it shows that the average call setup time is 23.11 seconds, within which the wireless interface reconnection time requires 14.05 seconds in average. The call setup time in the 802.11 case is shorter than the cellular interfaces because SIP signaling only goes through two networks. However, no matter which type of interfaces is used, the call setup time is not short. This is why we design our push server to temporarily answer an incoming call to keep the session alive, or the caller may hang up the call before the call is established.

E. Comparison of Signaling Cost

Next, we compare the signaling cost of our proposed approach and the MIPv6-NEMO scheme with routing optimization. We consider two cases. One is the offline case, where SIP-MNG is disconnected from the Internet; the other is the online case, where SIP-MNG is connected to the Internet. In the offline case, for SIP-MNG, there is no SIP signaling cost when it moves from one network domain to another. The external push server would be SIP-MNG's representative, which answers and transfers incoming sessions for the mobile network. On the other hand, since the MIPv6-NEMO scheme does not support the disconnected feature, the MR still has to track network signaling and update with its HA when handoff. Suppose that the HA binding update cost is $cost_{HABU}$. In the online case, the SIP signaling cost of SIP-MNG when handoff will be $N \times cost_{SIP-reregistration} + S \times cost_{SIP-reINVITE}$, where N is the number of MHs in the mobile network, S is the number of sessions between the mobile network and the Internet. On the other hand, the MIPv6-NEMO scheme with routing optimization will require a cost of $cost_{HABU} + M \times cost_{BU}$ when handoff, here we assume that the routing optimization approach based on binding update for network prefixes is used, and M is the number of CNs that are communicating with the mobile network. To summarize, in the online case, the SIP has its disadvantage in terms of signaling cost as compared to the MIPv6-NEMO scheme with routing optimization. However, in the offline case, our approach incurs lower cost than the

MIPv6-NEMO scheme, which may compensate some of the signaling cost in the on-line case.

VI. CONCLUSIONS

This paper proposes a SIP-based mobile network architecture to support networking services on the roads. With multiple wireless interfaces, a SIP-MNG can provide dynamic bandwidth to internal users based on their bandwidth requirements. Also, by allowing multiple sessions to share one interface, our system can help users or public transportation operators to save Internet access fees. Moreover, through our system, vehicles can provide Internet access to passengers with support of group mobility. By interpreting SIP signaling, our RM and CAC mechanisms inside the SIP-MNG can guarantee QoS for users. In addition, through session control and SMS, we propose a push mechanism to allow the SIP-MNG to stay off-line when there is no calling activity and to be “woken up” when necessary. Our push approach can save energy and call charges, while maintaining global reachability of users. In our architecture, we do not modify the current SIP client-server architecture and protocol.

A prototype is developed and some experimental results are presented. For IEEE 802.11, WCDMA, and PHS networks, we demonstrated that it is feasible to allow multiple stations to share one interface. It is also shown that, by cellular interfaces, the call setup time and handoff delay are longer than that by 802.11 interfaces, because connecting to the Internet via a cellular interface has to go through more networks. For our push mechanism, based on current technologies, the call setup time is in the range of 20 seconds, which is some what long. So we have designed the push server to temporarily pick up the session and apply the REFER scheme to transfer the session to the user inside SIP-MNG. The wireless interface reconnection time takes the largest part of the time. Although this is not in the scope of this work, we believe that a lot of research results can help reduce the reconnection time [27], [28].

REFERENCES

- [1] C. Perkins. IP Mobility Support for IPv4. *IETF RFC 3344*, Aug. 2002.
- [2] D. Johnson, C. Perkins, and J. Arkko. Mobility Support in IPv6. *IETF RFC 3775*, Jun. 2004.
- [3] J. Rosenberg, H. Schulzrinne, G. Camarillo, A. Johnston, J. Peterson, R. Sparks, M. Handley, and E. Schooler. SIP: Session Initiation Protocol. *IETF RFC 3261*, June 2002.
- [4] T. Ernst, K. Uehara, and K. Mitsuya. Network Mobility from the InternetCAR Perspective. In *Proc. of the 17th IEEE Int'l Conf. on Advanced Inf. Networking and Appl.*, pages 19–25, Mar. 2003.
- [5] H.-Y. Lach, C. Janneteau, and A. Petrescu. Network Mobility in Beyond-3G Systems. *IEEE Communications Magazine*, 41(7):52–57, Jul. 2003.
- [6] E. Perera, V. Sivaraman, and A. Seneviratne. Survey on Network Mobility Support. *ACM SIGMOBILE Mobile Computing and Communications Review*, 8(2):7–19, Apr. 2004.
- [7] V. Devarapalli, R. Wakikawa, A. Petrescu, and P. Thubert. Network Mobility Basic Support Protocol. *IETF RFC 3963*, Jan. 2005.
- [8] T. Ernst. Network Mobility Support Goals and Requirements. *IETF DRAFT: draft-ietf-nemo-requirements-06*, Nov. 2006.
- [9] T. Ernst and H.-Y. Lach. Network Mobility Support Terminology. *IETF DRAFT: draft-ietf-nemo-terminology-06*, Nov. 2006.
- [10] C.-M. Huang, C.-H. Lee, and J.-R. Zheng. A Novel SIP-Based Route Optimization for Network Mobility. *IEEE Journal on Selected Areas in Communications*, 24(9):1682–1691, Sep. 2006.
- [11] Y.-B. Lin, Y.-C. Lo, and C.-H. Rao. A Push Mechanism for GPRS Supporting Private IP Addresses. *IEEE Communications Letters*, 7(1):24–26, Jan. 2003.

- [12] Y.-B. Lin, W.-E. Chen, and C.-H. Gan. Effective VoIP Call Routing in WLAN and Cellular Integration. *IEEE Communications Letters*, 9(10):874–876, Oct. 2005.
- [13] C. E. Perkins and P. Bhagwat. Highly Dynamic Destination-Sequenced Distance-Vector Routing (DSDV) for Mobile Computers. *Proc. of the ACM SIGCOMM*, 234–244, 1994.
- [14] C.-C. Chiang, H.-K. Wu, W. Liu, and M. Gerla. Routing in Clustered Multihop, Mobile Wireless Networks with Fading Channel. *Proc. IEEE Singapore Int. Conf. on Networks*, pages 197–211, Apr. 1997.
- [15] P. Srisuresh and K. Egevang. Traditional IP Network Address Translator (Traditional NAT). *IETF RFC 3022*, Jan. 2001.
- [16] J. Rosenberg, J. Weinberger, C. Huitema, and R. Mahy. STUN - Simple Traversal of User Datagram Protocol (UDP) Through Network Address Translators (NATs). *IETF RFC 3489*, Mar. 2003.
- [17] J. Rosenberg. Interactive Connectivity Establishment (ICE): A Methodology for Network Address Translator (NAT) Traversal for Offer/Answer Protocols. *IETF DRAFT: draft-ietf-mmusic-ice-15*, Mar. 2007.
- [18] M. Handley, V. Jacobson, and C. Perkins. SDP: Session Description Protocol. *IETF RFC 4566*, July 2006.
- [19] R. Sparks. The Session Initiation Protocol (SIP) Refer Method. *IETF RFC 3515*, Apr. 2003.
- [20] J. Rosengerg, J. Peterson, H. Schulzrinne, and G. Camarillo. Best Current Practices for Third Party Call Control (3pcc) in the Session Initiation Protocol (SIP). *IETF RFC 3725*, Apr. 2004.
- [21] netfilter. <http://www.netfilter.org/>.
- [22] C.-H. Rao, D.-F. Chang, and Y.-B. Lin. iSMS: An Integration Platform for Short Message Service and IP Networks. *IEEE Networks*, 15:48–55, Mar./Apr. 2001.
- [23] SIP Express Router (ser). <http://www.iptel.org/ser/>.
- [24] W. Wang, S.-C. Liew, and V. O. K. Li. Solutions to Performance Problems in VoIP Over a 802.11 Wireless LAN. *IEEE Transactions on Vehicular Technology*, 54(1):366–384, Jan. 2005.
- [25] A. Mishra, M. Shin, and W. Arbaugh. An Empirical Analysis of the IEEE 802.11 MAC Layer Handoff Process. *ACM SIGCOMM Computer Comm. Rev.*, 33(2):93–102, April 2003.
- [26] B. Aboba. Fast Handoff Issues. IEEE-03-155r0-1. *IEEE 802.11 Working Group*, March 2003.
- [27] H. Velayos and G. Karlsson. Techniques to Reduce the IEEE 802.11b Handoff Time. In *IEEE ICC*, volume 7, pages 3844–3848, June 2004.
- [28] S. Park, P. Kim, M. Lee, and Y. Kim. *Parallel and Distributed Computing: Applications and Technologies*, chapter Fast Address Configuration for WLAN, pages 396–400. Springer Berlin/Heidelberg, 2004.

PLACE
PHOTO
HERE

Yu-Chee Tseng Yu-Chee Tseng received the B.S. degree from National Taiwan University, Taipei, Taiwan, in 1985 and the M.S. degree from National Tsing-Hua University, Hsinchu, Taiwan, in 1986, both in computer science, and the Ph.D. degree in computer and information science from Ohio State University, Columbus, in 1994.

He was an Associate Professor with Chung-Hua University, Hsinchu, from 1994 to 1996 and with National Central University, Taoyuan, Taiwan, from 1996 to 1999, and a Professor with National Central University from 1999 to 2000. Since 2000, he has been a Professor with the Department of Computer Science, National Chiao-Tung University, Hsinchu, where he is also currently the Department Chairman. His research interests include mobile computing, wireless communication, network security, and parallel and distributed computing.

PLACE
PHOTO
HERE

Jen-Jee Chen Jen-Jee Chen received the B.S. and M.S. degrees in Computer Science and Information Engineering from National Chiao Tung University, Hsinchu, Taiwan, in 2001 and 2003, respectively. He is currently pursuing his PhD degree in Computer Science at National Chiao Tung University, Hsinchu, Taiwan. His research interests are primarily in mobile computing, wireless networks, network performance modeling, and internet communication service.

PLACE
PHOTO
HERE

Yu-Li Cheng Yu-Li Cheng received the B.S. degree in Engineering Science and Ocean Engineering from the National Taiwan University, Taipei, Taiwan, in 1995 and the M.S. degree in Computer Science from the National Chiao-Tung University, Hsinchu, Taiwan, in 2007. His research interests include wireless communication and mobile computing. He is working at Alpha Networks Inc. in Hsin-Chu Science Park, Taiwan as a software engineer in the Wireless Broadband product division since 2001.